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#### VALUATIONS OF POLYNOMIAL RINGS

by

#### Ronald Alvin Macauley

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#### Abstract

is a field on which all (non-archimedean) valua-If R tions are known, then all valuations on R(x), where x is transcendental over R, are also known. Ostrowski described such valuations of R[x] by means of pseudo-convergent sequences in the algebraic completion of A of R. MacLane later showed that if all valuations of R are discrete, then any valuation V of R[x] can be represented by certain "key" polynomials in R [x]. The present paper exhibits the connection between these two treatments. This is achieved by first determining keys for the evaluation which a pseudoconvergent sequence defines on A[x], and then relating these keys to those for V .

# <u>A c k n o w l e d g m e n t</u>

The writer wishes to express his thanks to Dr. B.N. Moyls of the Department of Mathematics at the University of British Columbia for his advice and guidance. His numerous criticisms and suggestions proved invaluable in the preparation of this thesis. 1. <u>Introduction</u>. A non-archimedean valuation V, hereafter simply called a valuation, of an integral domain R is a single-valued mapping of the elements of R into the real numbers and  $+\infty$  such that:

1) Va is a unique finite real number for  $a \neq 0$ ,

- 2)  $V_0 = +\infty$
- 3) V(ab) = Va + Vb for all a,  $b \in \mathbb{R}$ ,
- 4)  $V(a + b) \ge \min \{Va, Vb\}$  for all  $a, b \in \mathbb{R}$ .

An extremely important property of these valuations is that, if  $Va \neq Vb$ , then  $V(a + b) = min \{Va, Vb\}$ ; hence, if V(a + b) > Va, Va = Vb.

Ostrowski, and later MacLane, attacked the problem of finding all extensions of valuations on an integral domain R to the ring of polynomials R[x], where x is transcendental over R. MacLane's results are based on the assumption that all valuations of R are discrete; that is, the real numbers used as values form an isolated point set. It is the purpose of this paper to provide a connection between the valuations of Ostrowski and MacLane on R[x], where R is a field with only discrete valuations.

<u>Definition 1.1</u>: Let R be a field with a valuation V. The sequence  $\{\alpha_i\}$ , where  $\alpha_i \in \mathbb{R}$ , is a pseudo-convergent sequence with respect to V if  $V(\alpha_i - \alpha_{i-1}) < V(\alpha_{i+1} - \alpha_i)$  for all i > N, some fixed positive integer.

If  $\{a_i\}$  is a pseudo-convergent sequence with respect to

V, then the sequence  $\{Va_i\}$  is eventually strictly monotone increasing or eventually attains a constant value; as long as 0 is not a limit of  $\{\alpha_i\}, \{\forall \alpha_i\}$  converges to a finite limit. This is important, for it is essentially this property that Ostrowski uses to extend V on R to W on R[x], where W is a valuation of R[x]. He shows ([1], III, page 371) that if  $f(x) \in R[x]$ , then  $\{f(\alpha_i)\}$  is also a pseudo-convergent sequence. This implies that  $\{Vf(\alpha_i)\}$  is convergent to a finite limit except when  $\{\alpha_i\}$  converges to a root of f(x). Hence, if  $\{\alpha_i\}$  is a pseudo-convergent sequence possessing no limit in the algebraic completion A of R, the function W on. R[x] defined by  $Wf(x) = \lim_{x \to \infty} Vf(\alpha_i)$  is a valuation ([1], section 65, page 374) of R[x]. Further, ([1], IX, page 378) every valuation of R[x] may be obtained by means of some pseudo-convergent sequence in A. The pseudo-convergent sequences in A are valued by an extension of V on R to A; this extension always exists([1], II, page 300). This last reference implies that any valuation of R(x) may be extended A(x). Hence, if all valuations of A[x] are found, all to valuations of R[x] are automatically found. This result is of prime importance to the development of the theory in this paper.

<u>Definition 1.2</u>: Let K be an integral domain with a valuation V. Two elements a, b  $\varepsilon$  K are equivalent with respect to V, written a ~ b (V), if V(a - b) > Va.

<u>Definition 1.3</u>: For a, b  $\varepsilon$  K, a equivalence divides b in V if there exists c  $\varepsilon$  K such that b ~ ca (V); notation:

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a|b(V).

is any valuation of R[x] which reduces to a If V discrete valuation  $V_0$  of R, MacLane ([2]) represents V by the following inductive method: a value  $V_1 x = V x = \mu_1$ is assigned to x. Then for any polynomial  $f(x) \in \mathbb{R}[x]$ ,  $f(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_0$ , a function  $V_1$  on R[x] is defined by  $V_{l}f(x) = \min_{i} \{ V_{0}a_{i} + i\mu_{l} \}$ . This function  $V_{l}$  may be shown to be a valuation of R [x] such that  $V_1 \leq V$ ; that is,  $V_1g(x) \leq Vg(x)$  for all  $g(x) \in R[x]$ . The value  $V_1$  is called a first stage value and is symbolized by  $V_1 = [V_0, V_1x = \mu_1]$ . Either  $V_1 = V$ , that is,  $V_1g(x) = Vg(x)$  for all  $g(x) \in \mathbb{R}[x]$ , or there exists an  $f(x) \in \mathbb{R}[x]$  such that  $V_1f(x) < Vf(x)$ . If that latter is the case,  $\varphi_2(x) \in R[x]$  is chosen such that  $\varphi_2$ is a monic polynomial of the smallest degree satisfying  $V_1 \varphi_2 < V \phi_2$ . This polynomial satisfies, over  $V_1$ , MacLane's conditions for a key polynomial.

<u>Definition 1.4</u>: Let W be any valuation of R[x]. A polynomial  $\varphi \in R[x]$  is a key polynomial over the value W if:

(i)  $\varphi$  is equivalence irreducible -  $\varphi|a(x)b(x)$  (W) implies either  $\varphi|a(x)$  (W) or  $\varphi|b(x)$  (W),

(ii)  $\varphi$  is minimal -  $\varphi|a(x)$  (W) implies deg  $a(x) \ge \deg \varphi$ , (iii)  $\varphi$  is monic.

It is shown ([2], theorem 4.2) that if a key polynomial  $\varphi$  over W is assigned a value  $\mu = W'\varphi > W\varphi$ , then the function W' on R[x] defined by

 $W'f(x) = \min_{i} \{Wf_{i}(x) + i\mu\},\$ 

where  $f(x) = f_n(x) \phi^n + f_{n-1}(x) \phi^{n-1} + \dots + f_{\Theta}(x)$ , deg  $f_i(x) < \deg \varphi$ , is a valuation of R [x]. Further,  $W \le W'$ and Wf(x) < W'f(x) if and only if  $\varphi f(x)$  (W); in particular, Wf(x) = W'f(x) if deg  $f(x) < \deg \varphi([2], \text{ theorem 5.1})$ . In the original valuation V, the polynomial  $\varphi_2$  chosen above will define a valuation  $V_2$  on R[x] if assigned the value  $\mu_2$  = V  $_2$  > V  $_1 \phi_2$  . The value V  $_2$  satisfies V  $_2 \leq$  V and and  $V_2f(x) = V_1f(x) = Vf(x)$  for all  $f(x) \in R[x]$  such that deg  $f(x) < \deg \varphi_2$ . The second-stage value V<sub>2</sub> is symbolized by  $V_2 = [V_0, V_1 x = \mu_1, V_2 \varphi_2 = \mu_2]$ . As before, either  $V_2 = V$ or there exists a monic polynomial  $\varphi_3$  of minimum degree satisfying  $V_2 \varphi_3 < V \varphi_3$ . Again, if  $\varphi_3$  exists it is a key polynomial over  $V_2$  and may be used to define a valuation  $V_3$  such that  $V_3 \leq V$  and  $V_3f(x) = V_2f(x) = Vf(x)$  for all  $f(x) \in R[x]$ with deg  $f(x) < \deg \varphi_3$ . The third-stage value  $V_3$  is symbolized by  $V_3 = [V_0, V_{1x} = \mu_1, V_2 \varphi_2 = \mu_2, V_3 \varphi_3 = \mu_3]$ . MacLane shows ([2], theorem 8.1) that if this procedure is continued, equality will occur after either a finite or countable number of steps. In the first case V will have a representation

 $V = V_k = [V_0, V_1 x = \mu_1, V_2 \varphi_2 = \mu_2, \dots, V_k \varphi_k = \mu_k],$ and is called an inductive value. In the latter case

 $V = V_{\infty} = [V_0, V_1 x = \mu_1, V_2 \varphi_2 = \mu_2, \dots, V_k \varphi_k = \mu_k, \dots],$ where  $V_{\infty}f(x) = \lim_{k \to \infty} V_k f(x)$ , and V is called a limit value. Hence each valuation of R[x] may be represented by one of these two cases if every valuation of R is discrete.

The key polynomials defining the above inductive and limit values satisfy:

(iv) 
$$\varphi_i \sim \varphi_{i-1} (V_{i-1})$$
 is false for all  $i \ge 2$ ,

(v) deg  $\varphi_i \ge \deg \varphi_{i-1}$  for all  $i \ge 2$ . Now since every valuation of  $\mathbb{R}[x]$  is either an inductive value or the limit of a sequence of inductive values it is necessary only to consider key polynomials which also satisfy (iv) and (v). For this reason it will be assumed in this paper that a key polynomial is a polynomial satisfying (i), (ii), (iii), (iv) and (v).

The representation of a valuation on R[x] is not necessarily unique, but if one additional restriction is placed on the key polynomials the representation becomes unique when V<sub>o</sub> on R is discrete. Let  $V_{k} = [V_{o}, V_{1}x = \mu_{1}, \dots, V_{k}\phi_{k} = \mu_{k}]$ be an inductive valuation of R[x]. The  $V_{k}$  value of  $f(x) \in R[x]$  is found from the expansion

$$f(x) = f_n(x) \phi_k^n + f_{n-1}(x) \phi_k^{n-1} + \dots + f_o(x)$$

where deg  $f_i(x) < \deg \varphi_k$ . By expanding each  $f_i(x)$  in powers of  $\varphi_{k-1}$  and the coefficients of this expansion again in powers of  $\varphi_{k-2}$  and continuing these expansions, finally the expression

$$f(\mathbf{x}) = \sum_{j} a_{j} \mathbf{x}^{m_{j}} \phi_{2}^{m_{2}} \cdots \phi_{k}^{m_{k}}$$

where a; E R and each

$$n_{ij} < \frac{\deg \varphi_{i+1}}{\deg \varphi_i}$$
 ,

is obtained. Furthermore,  $V_k f(x) = \min_j \left\{ V_k \begin{bmatrix} a_j x^{m_1} j \phi_2^{m_2} j & \cdots & \phi_k \end{bmatrix} \right\}$ . Now, the elements of R may be partitioned into classes of equi-

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valent elements with respect to  $V_0$  and a representative may chosen from each class. In particular, the element 1 is to be chosen as a representative. These representatives are called the  $V_0$ -representatives. If in the above expansion of f(x)each aj is a  $V_0$ -representative and all terms have the minimum value  $V_k f(x)$ , then f(x) is called homogeneous in  $V_k$ . Every polynomial  $f(x) \in R[x]$  is equivalent in  $V_k$  to one and only one homogeneous polynomial  $h(x) \in R[x]$  ([2], lemma 16.2); h(x) is called the homogeneous part of f(x). An inductive or limit value is called a homogeneous value if each  $\varphi_1$ ,  $i \ge 2$ , is homogeneous in  $V_{i-1}$ . MacLane has shown ([2], theorem 16,3 and 16.4) that any inductive or limit value constructed from a discrete  $V_0$  may be represented by one and only one homogeneous inductive or limit value.

The inductive and limit values of R[x] will always be considered to be homogeneous values.

2. The relation between the valuations of Ostrowski and MacLane will first be established on A [x], where A is an algebraically complete field.

It will be found convenient, in this section and future sections, to remove the condition that a MacLane value has first key x. It is necessary only that the first key be linear and monic. The properties of MacLane values will be preserved.

Every valuation V of A[x] may be defined by some pseudoconvergent sequence  $\{\alpha_i\}$ , with respect to V<sub>o</sub> on A, which does

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not possess a limit in A ; Vf(x) is defined as  $Vf(x) = \lim_{i \to \infty} V_0f(\alpha_i)$ .

These pseudo-convergent sequences may be divided into two types. To obtain the desired classification, a pseudo-limit is defined.

<u>Definition 2.1</u>: An element  $\alpha \in A$  is pseudo-limit of the pseudo-convergent sequence  $\{\alpha_i\}$ , where  $\alpha_i \in A$ , with respect to the valuation  $V_0$  if  $V_0(\alpha - \alpha_i) = \delta_i$ , where  $\delta_i < \delta_{i+1}$ . for i >some integer N.

<u>Note</u>: "Pseudo-limit" as defined here is not the same as that defined by Ostrowski.

Now  $V_0(\alpha_i - \alpha_{i+1}) = \mathcal{X}_i$ , where  $\mathcal{X}_i < \mathcal{X}_{i+1}$  for i > some integerN'. Since  $\mathcal{X}_i = V_0(\alpha_i - \alpha_{i+1}) = V_0[(\alpha_i - \alpha) + (\alpha - \alpha_{i+1})]$  it follows that, for i > N,  $\mathcal{X}_i = \delta_i$ .

The pseudo-convergent sequences are now divided into two classes:

(1)  $\{\alpha_i\}$  possesses a pseudo-limit in A,

(2)  $\{\alpha_i\}$  does not possess a pseudo-limit in A,

<u>Theorem 2.2</u>: If the pseudo-convergent sequence  $\{\alpha_i\}$ , with respect to  $V_0$ , has a pseudo-limit  $\alpha \in A$ , then the Ostrowski valuation V of A [x] defined by  $\{\alpha_i\}$  is the same as the first-stage valuation V, defined by  $V_i = [V_0, V_i (x - \alpha) = Y]$ , where  $Y = \lim_{i \to \infty} Y_i = \lim_{i \to \infty} V_0 (\alpha_i - \alpha_{i+1})$ .

<u>Proof</u>: It is sufficient to consider a monic linear polynomial  $x - \beta$  in A[x]. Since  $\alpha_i - \beta = (\alpha_i - \alpha) + (\alpha - \beta)$  and  $V_0(\alpha_i - \alpha) = \mathcal{V}_i$ , either  $V_0(\alpha_i - \beta) = \mathcal{V}_i$  or  $V_0(\alpha_i - \beta)$ =  $V_0(\alpha - \beta)$ , for i sufficiently large. Hence

$$\mathbb{V}(\mathbf{x} - \beta) = \lim_{i \to \infty} \mathbb{V}_0(\alpha_i - \beta) = \min\left\{\mathcal{X}, \mathbb{V}_0(\alpha - \beta)\right\} = \mathbb{V}_1(\mathbf{x} - \beta)$$

<u>Theorem</u> 2.3: Given a finite inductive value  $V = [V_0, V(x - \alpha) = \delta]$ on A [x], a pseudo-convergent sequence  $\{\alpha_i\}$  with pseudo-limit  $\alpha \in A$  can be found such that

$$\gamma = \lim_{i \to \infty} \gamma_i$$

<u>Proof</u>: Let  $a \neq 0$  in A be chosen such that Va = d > 0. Then there exists a real number  $\sigma$  such that  $\sigma d = a$ . A sequence of integers  $\{n_i\}$  can be found such that

$$\frac{n_1}{10} < \frac{n_2}{10^2} < \dots < \frac{n_i}{10^i} < \dots$$

and

$$\lim_{i\to\infty}\frac{n_i}{10^i}=\sigma$$

Let  $\beta_i$  be any one of the roots of  $x'^{o^i} - a$ . Then  $d = Va = V \beta_i^{o^i} = 10^i V \beta_i$  or  $V \beta_i^{i} = 1/10^i d$ . Hence, the sequence  $\{V \beta_i^{n_i}\}$  is a strictly increasing sequence with

$$\lim_{i \to \infty} V \beta_i^{n_i} = \lim_{i \to \infty} \frac{n_i}{10^i} d = \sigma d = \lambda^4.$$

Let  $\alpha_i$  be defined by  $\alpha_i = \beta_i^{n_i} + \alpha$ . Since  $V(\alpha_i - \alpha_{i+i}) = V(\beta_i^{n_i} - \beta_{i+i}^{n_{i+i}}) = V\beta_i^{n_i}, \{\alpha_i\}$ is a pseudo-convergent sequence and since  $V(\alpha_i - \alpha) = V\beta_i^{n_i}$ ,  $\alpha$  is a pseudo-limit of this sequence. By Theorem 2.2, the sequence  $\{\alpha_i\}$  can have no limit in A, since the Ostrowski value defined by  $\{\alpha_i\}$  is also defined by the finite value V.

On combining Theorem 2.2 and Theorem 2.3 an equivalence is obtained between valuations defined by pseudo-convergent sequences with pseudo-limits and the inductive values of MacLane.

Theorem 2.4: If  $\{a_i\}$  is a pseudo-convergent sequence with respect to  $V_O$ , with no pseudo-limit in A , then the Ostrowski valuation V defined by  $\{\alpha_i\}$  is the same as the MacLane limit value  $V' = \int V_0, V_1(x-\alpha_1) = V_1, ..., V_i(x-\alpha_i) = V_i, ... ],$ where  $\lambda_i = V_o(\alpha_i - \alpha_{i+1})$ . Also, in a MacLane valuation V' the sequence  $\{\alpha_i\}$  is pseudo-convergent with no limit in A and the Ostrowski value V defined by  $\{\alpha_i\}$  is equal to V'. Proof: If necessary remove a finite number of terms from the beginning of the pseudo-convergent sequence  $\{a_i\}$  and renumber the a's, so that  $Y_i = V_0(\alpha_i - \alpha_{i+1})$  is strictly increasing. Since  $V_{o}(\alpha_{i \neq n} - \alpha_{i}) = \mathcal{X}_{i}$  for all  $n \geq 1$ ,  $V(x - \alpha_{i}) = \lim_{n \to \infty} V_{o}(\alpha_{i + n} - \alpha_{i}) = \mathcal{X}_{i}$ . Let  $V_1$  be defined by  $V_1 = [V_0, V_1(x-\alpha_1) = \mathcal{Y}_1]$ , then  $V_1$  defines a first-stage value of A [x] such that  $V_1 \leq V$ . Now, from MacLane's inductive argument used in the introduction, it follows immediately that  $V' = [V_{\alpha}, V_1(x-\alpha_1) = \mathcal{V}_1, \dots, V_i(x-\alpha_i) = \mathcal{V}_i, \dots]$ is a MacLane valuation and also satisfies V'  $\leq$  V . If there exists  $x - \sqrt{3} \in A[x]$  such that  $V'(x - \sqrt{3}) < V(x - \sqrt{3})$ =  $\lim_{i\to\infty} V_0(\alpha_i - \beta)$ , then there exists a positive integer N such that  $V'(x - \beta) < V_{\alpha}(\alpha_i - \beta)$  for all i > N. Therefore, from  $\begin{array}{l} \mathbb{V}\left(\mathbf{x}-\beta\right) = \lim_{i\to\infty} \left[\min\left\{\gamma_{i}, \mathbb{V}_{o}\left(\alpha_{i}-\beta\right)\right\}\right] & \text{it follows that } \lim_{i\to\infty} \gamma_{i} \\ < \mathbb{V}_{o}\left(\alpha_{i}-\beta\right) & \text{for } i > \mathbb{N} & \text{Hence, } \mathbb{V}_{o}\left(\alpha_{i+1}-\beta\right) = \mathbb{V}_{o}\left[\left(\alpha_{i+1}-\alpha_{i}\right)\right] \end{array} \right) \end{array}$ +  $(\alpha_i - \beta_i) = \lambda_i \leq \lim_{i \to \infty} \lambda_i$ , for i > N, which is a contradic-tion. Thefefore V' = V. Suppose now that V' is a Machane valuation. From  $x - \alpha_{i+1} = (x - \alpha_i) + (\alpha_i - \alpha_{i+1})$  it follows that  $V_o(\alpha_i - \alpha_{i+1}) \ge V_i(x - \alpha_i) = \mathcal{V}_i$ , for otherwise  $V_{i}\left[(x - \alpha_{i+1}) - (\alpha_{i} - \alpha_{i+1})\right] = V_{i}(x - \alpha_{i}) = \mathcal{V}_{i} > V_{o}(\alpha_{i} - \alpha_{i+1})$ 

and so  $x - \alpha_{i+1} \sim \alpha_i - \alpha_{i+1}(V_i)$ , which contradicts the minimal

condition (ii) of definition 1.4 for a key polynomial over  $V_i$ . Therefore  $V_i(x - \alpha_{i+1}) = \mathcal{V}_i$  and since  $V_i(x - \alpha_{i+1}) < V_{i+1}(x - \alpha_{i+1})$ ,  $\mathcal{V}_i < \mathcal{V}_{i+1}$  for all  $i \ge 1$ . Now

 $V_{0}(\alpha_{i} - \alpha_{i+1}) = V'[(x - \alpha_{i+1}) - (x - \alpha_{i})] = V_{i};$ hence {a<sub>i</sub>} is a pseudo-convergent sequence. If '/3 were a limit of {a<sub>i</sub>}, then lim  $V_{0}(\alpha_{i} - \beta) = \infty$ . Let k > i be chosen such that  $V_{0}(\alpha_{k} - \beta) > V_{1}$ , then

 $V_o(\alpha_i - \beta) = V_o[(\alpha_i - \alpha_k) + (\alpha_k - \beta)] = \lambda_i.$ 

Therefore

$$\mathbb{V}^{\dagger}(\mathbf{x} - \mathbf{3}) = \lim_{i \to \infty} \left[ \min \left\{ \mathcal{J}_{i}, \mathbb{V}_{0}(\alpha_{i} - \mathbf{3}) \right\} \right] = \lim_{i \to \infty} \mathbb{V}_{0}(\alpha_{i} - \mathbf{3}) = \infty;$$

but V' is a finite value. Hence  $\{\alpha_i\}$  has no limit in A and will, therefore, define an Ostrowski valuation which, by first part of theorem 2.4., must be the same as V'. NOTE; If  $\{\alpha_i\}$  has a pseudo-limit  $\alpha \in A$ , then V' may also be represented by

$$V^{n} = \left[ V_{0}, V^{n}(x - \alpha) = \lim_{i \to \infty} V_{i} \right].$$

The results of this section now provide a connection between the two methods of valuation A [x].

In sections 4 and 5 it will be shown how a MacLane valuation of A[x] reduces to a MacLane valuation of R [x], that is, the key polynomials and their assigned values will be found for the reduced valuation on R [x], and conversely how to extend a value on R [x] to A [x]. The connection between an Ostrowski valuation of R [x] and a MacLane valuation of R [x] will then be clear. 3. The key polynomials defining the restriction of a valuation of A [x] to R [x] are intimately related to the key polynomials used by MacLane ([3]) to extend a valuation  $V_{\Theta}$  on R to a valuation W of R ( $\alpha$ ), as separable extension of R. For this reason a description of the methods used by MacLane and the essential results will now be given.

As a particular example, consider the inductive value

$$\mathbf{V}_{\mathbf{k}} = \begin{bmatrix} \mathbf{V}_{\mathbf{0}}, \ \mathbf{V}_{\mathbf{1}}\mathbf{x} = \boldsymbol{\mu}_{\mathbf{1}}, \mathbf{V}_{\mathbf{2}}\boldsymbol{\varphi}_{\mathbf{2}} = \boldsymbol{\mu}_{\mathbf{2}}, \ \cdots, \ \mathbf{V}_{\mathbf{k}}\boldsymbol{\varphi}_{\mathbf{k}} = \boldsymbol{\mu}_{\mathbf{k}} \end{bmatrix}$$

of R[x] and reassign to  $\varphi_k$  the value +  $\infty$ . This defines a new, generalized valuation

 $V = \begin{bmatrix} V_0, V_1 x = \mu_1, V_2 \varphi_2 = \mu_2, \cdots, V_{k-1} \varphi_{k-1} = \mu_{k-1}, V \varphi_k = \infty \end{bmatrix}$ of R [x]. The generalized valuation V satisfies all the conditions of a valuation except that elements other than 0 are assigned the value + $\infty$ . If  $\alpha$  is a root of  $\varphi_k$ , the valuation V will define a valuation W on R ( $\alpha$ ). This is immediately seen upon noticing that

$$\mathbb{R} (\alpha) \cong \frac{\mathbb{R} [x]}{(\varphi_k)}$$

and defining W by  $Wf(\alpha) = Vf(x)$ . If the  $\varphi_i$ , for  $2 \le i \le k$ , above are homogeneous in the preceding inductive value  $V_{i-1}$ , MacLane has shown ([3], theorem 5.3) that this extension W of  $V_0$  is the only extension of  $V_0$  to R ( $\alpha$ ).

To facilitate the discussion of the remainder of this section and in view of sections 4 and 5, it is convenient at this point to define the terms projection and effective degree. <u>Definition 3.1</u>: Let  $V_k = \begin{bmatrix} V_0, V_1 x = \mu_1, \dots, V_k \phi_k = \mu_k \end{bmatrix}$  be an inductive value of R [x]. If

 $G(x) = g_n(x) \phi_k^n + g_{n-1}(x) \phi_k^{n-1} + \dots + g_{\phi}(x)$ ,

where deg  $g_i(x) < \deg \varphi_k$ , is a polynomial in R [x], then the projection of  $V_K$  with respect to G(x) is  $\alpha - \beta$ , written proj  $(V_k) = \alpha - \beta$ , where  $\alpha$  and  $\beta$  are the maximum and minimum values respectively of i such that

$$V_k G(x) = V_k [g_i(x) \phi_k^i].$$

<u>Definition 3.2</u>: The effective degree of  $G_{i}(x)$  in  $G_{i}(x)$  is a : written  $D_{\phi}G(x) = a$ .

Let W be a valuation of  $R(\alpha)$ , where W is an extension of  $V_0$  on R and  $\alpha$  has minimal polynomial  $G(x) \in R[x]$ . By the isomorphism  $R(\alpha) \cong R[x]/(G(x))$  it is clear that a generalized valuation V on R[x] may be defined by  $Vf[x] = Wf(\alpha)$ . The valuation V assigns the value  $+\infty$  only to the members of the ideal (G(x)). It would seem natural to construct V as MacLane does for finite valuations; that is, for valuations which assign the value  $+\infty$  only to 0. As before, a first-stage value  $V_1 = [V_0, V_1x = \mu_1]$ , where  $\mu_1 = Vx \neq \infty$ , is defined; again  $V_1 \leq V$ . It is worth noting that proj  $(V_1) > 0$ . For if proj  $(V_1) = 0$  then would be only one term in

 $G(x) = a_n x^n + a_{n-1} x^{n-1} + ... + a_0$ 

with minimum value, and  $VG(x) = V_1G(x) \neq \infty$ . To define a seconstage value a monic polynomial f(x) of minimal degree satisfying  $Vf(x) > V_1f(x)$  is chosen. If f(x) is not homogeneous in  $V_1$ , then its homogeneous part is to be chosen. Denote this homogeneous part by  $\varphi_2$ . As was mentioned in the introduction,  $\varphi_2$  is a key

polynomial over  $\tilde{y}_1$ . The second stage value  $V_2$  is then defined by  $V_2 = \begin{bmatrix} V_0, V_1 x = \mu_1, V_2 \varphi_2 = \mu_2 \end{bmatrix}$ , where  $\mu_2 \neq V \varphi_2$ . Now, if G(x) is a homogeneous key over  $V_1$ , then  $\varphi_2$  is chosen as G(x)and  $\mu_2 = \infty$ . That G(x) is a monic polynomial of minimal degree satisfying  $VG(x) > V_1G(x)$  will follow from lemmas 3.3 and 3.4. Lemma 3.3: Let  $V_k$  be a k-th stage inductive value of R[x]satisfying:

- (1)  $V_k f(x) \leq V f(x)$  for all  $f(x) \in R[x]$ ,
- (2) deg  $f(x) < \deg \varphi_k$  implies  $V_k f(x) = V f(x)$ ,
- (3)  $\nabla_k \varphi_i = \nabla \varphi_i = \mu_i$  for  $1 \le i \le k$ .

If  $\psi$  is a monic polynomial of minimal degree satisfying  $V_k \psi < V \psi$ , then  $V_k f(x) < V f(x)$  implies  $\psi | f(x) (V_k)$ . <u>Proof:</u> Let f(x) have the quotient remainder expression f(x)  $= q^{(x)}\psi + r(x)$ , where deg  $r(x) < \deg \psi$ . Then  $V_k [f - q\psi] = V [f - q\psi] \ge \min \{Vf, V[q\psi]\} > \min \{V_k f, V_k[q\psi]\},$ because of (2), the choice of  $\psi$  and the assumption (1) for q(x). Hence  $\psi | f(x) (V_k)$ . <u>Lemma 3.4</u>: Let  $V_k$  be an inductive value of R[x]. Any poly-

nomial  $G(x) \in R[x]$  has an equivalence decomposition

$$G(\mathbf{x}) \sim e(\mathbf{x}) \varphi_{\mathbf{k}}^{\mathbf{t}_{0}} \psi_{\mathbf{l}}^{\mathbf{t}_{1}} \psi_{\mathbf{2}}^{\mathbf{t}_{2}} \cdots \psi_{\mathbf{r}}^{\mathbf{t}_{\mathbf{r}}} (\mathbf{v}_{\mathbf{k}}),$$

where each  $\psi_i$  is a homogeneous key over  $V_k$ ,  $t_0 \ge 0$  and  $t_i > 0$  for  $1 \le i \le r$ , and e(x) is an equivalence unit, that is,  $D_{\phi}e(x) = 0$  in  $V_k$ . This decomposition is unique except for equivalence units.

<u>Proof</u>: Cf. [3], theorem 4.2.

Now suppose  $\varphi_2$  is a homogeneous monic polynomial of minimal degree satisfying  $V_1\varphi_2 < V \ \varphi_2$  when G(x) is a homogeneous key over  $V_1$ . Since  $V_1G(x) < VG$ ,  $G(x) \sim h(x) \ \varphi_2$   $(V_1)$  by lemma 3.3. But  $G(x) \sim G(x) \ (V_1)$ , and therefore by lemma 3.4  $G(x) = \varphi_2$ . Hence in this case the value V is given by

$$\mathbf{V} = \left[ \mathbf{V}_{\mathbf{0}}, \ \mathbf{V}_{\mathbf{1}}\mathbf{x} = \boldsymbol{\mu}_{\mathbf{1}}, \ \mathbf{V}_{\mathbf{G}}(\mathbf{x}) = \boldsymbol{\infty} \right].$$

If G (x) is not a homogeneous key over  $V_1$ , then the second-stage is given by

$$\mathbf{v}_2 = \begin{bmatrix} \mathbf{v}_0, \ \mathbf{v}_1 \mathbf{x} = \boldsymbol{\mu}_1, \ \mathbf{v}_2 \boldsymbol{\varphi}_2 = \boldsymbol{\mu}_2 \end{bmatrix},$$

where  $V_2 \leq V$ . It is noticed again that  $\operatorname{proj}(V_2) > 0$  for otherwise  $VG(x) \neq \infty$ . Also  $\varphi_2 | G(x) (V_1)$  by lemma 3.3. MacLane's inductive process is repeated until G(x) does become a homogeneous key over some  $V_k$  or, if this does not occur, it is repeated a countable number of steps. In the former case

$$\nabla = \left[ \nabla_{o}, \nabla_{1} \mathbf{x} = \mu_{1}, \dots, \nabla_{k} \phi_{k} = \mu_{k}, \nabla G(\mathbf{x}) = \infty \right]$$

by the preceding argument. Also  $\varphi_1|_{G(x)}(V_{i-1})$  for all i such that  $2 \le i \le k$ , and proj $(V_i) > 0$  for  $1 \le i \le k$ . If a countable number of steps are required, then

 $\mathbb{V} = \mathbb{V}' = \begin{bmatrix} \mathbb{V}_0, \ \mathbb{V}_1 \mathbf{x} = \mu_1, \ \cdots, \ \mathbb{V}_k \phi_k = \mu_k, \ \cdots \end{bmatrix} .$  Certainly,  $\mathbb{V}_k \leq \mathbb{V}$  and  $\phi_k | G(\mathbf{x}) (\mathbb{V}_{k-1})$  for all  $k \geq 2$ . Since each  $\phi_k$  is minimal over  $\mathbb{V}_{k-1}$ , deg  $\phi_k \leq \deg G(\mathbf{x})$ . So from some point on all the keys will have the same degree. In this case it can be shown ([2], lemma 6.3) that the value group of  $\mathbb{V}'$  is discrete if the value group: of  $\mathbb{R}$  is discrete; that is, the real numbers used as values for  $\mathbb{V}'$  form an isolated point set. If  $\mathbb{V}'f(\mathbf{x}) < \mathbb{V}f(\mathbf{x})$  for some  $f(\mathbf{x}) \in \mathbb{K}[\mathbf{x}]$ , then  $\mathbb{V}_k f(\mathbf{x}) < \mathbb{V}f(\mathbf{x})$  for all k > 0 by the monotone increasing character of the in-

ductive values and so  $\varphi_{k+1} | f(x) (V_k) \cdot T_{herefore} V_k f(x) < V_{k+1} f(x)$ for all k > 0 (Cf. the introduction). But since the value group of V' is discrete,

 $V_f(x) \ge \lim_{k \to \infty} V_k f(x) = \infty$ .

Therefore only polynomials in (G(x)) could satisfy V'(f(x) < Vf(x)) but since  $V'(f'(x)) = \infty = Vf(x)$  for  $f(x) \in (G(x))$ , V = V'.

It is seen that every discrete  $V_0$  of  $\mathbb{R}$  may be extended to a finite separable extension  $\mathbb{R}(\alpha)$  of  $\mathbb{R}$  by MacLane's inductive process, where the homogeneous keys can be further restricted to satisfy the conditions proj  $(V_1) > 0$  for i > 0 and  $\varphi_i | G(x) (V_{i-1})$  for  $i \ge 2$ , where G(x) is the minimal polyof  $\alpha$ . In fact, it is not difficult to see that these restrictions are necessary.

From the preceding arguments it follows almost immediately that every such sequence of values constructed by these restricted keys will give a valuation of  $R(\alpha)$ .

The construction of such a sequence of values may be accomplished in a systematic manner. Let  $V_1 = \begin{bmatrix} V_0, V_1 x = \mu_1 \end{bmatrix}$  be a first-stage value such that proj  $(V_1) > 0$  for G(x). If  $V_{i-1}$  has been defined, the next key  $\varphi_i$  is chosen as any one of the  $\psi_j$  occurring in the unique equivalence decomposition of lemma 3.4. The corresponding value  $\mu_i$  is chosen so that proj  $(V_i) > 0$  and  $\mu_i > V_{i-1}\varphi_i$ . In the sequence of valuations so defined, each  $V_i$  is called an i-th approximant to G(x). MacLane not only shows that every such sequence of values defines a valuation

W on  $R(\alpha)$  which is an extension of  $V_{0}$ , but that :

- (1) if G(x) eventually becomes a homogeneous key over the inductive value  $V_k$ , then the i-th approximant is unique for  $2 \le i \le k$  and also, the value  $\mu_1$  is unique ([3], theorem 5.3). This implies that  $V_9$ may be extended to  $R(\alpha)$  in only one way ([3], theorem 10.1),
- (2) if a countable sequence of keys are required, then there is at most a finite number of different sequences that can be constructed. Hence,  $V_0$  on R may be extended to  $R(\alpha)$  in at most a finite number of ways ([3], theorem 10.1).

4. The reduction, or restriction, of an inductive value W of A[x] to R[x] will first be found; following theorem 4.7 the reduction of a limit value will be found. These results will be established by mathematical induction.

<u>Theorem 4.1</u>: If  $W = [W_0, W(x - \alpha) = Y]$  is any inductive value of A[x] with  $W_0 \alpha \ge Y$ , then  $W = W' = [W_0, W'x = Y]$ . <u>Proof</u>: Let  $x - /3 \in A[x]$ . If  $W_0(\alpha - /3) < Y$ , then  $W_0(\alpha - /3) = W_0/3$ ; and

 $W(x - \beta) = \min \{ \delta, W_0(\alpha - \beta) \} = \min \{ \delta, W_0\beta \} = W(x - \beta).$ If  $W_0(\alpha - \beta) \ge \delta$ , then  $W_0\beta \ge \min \{ W_0(\alpha - \beta), W_0\alpha \} \ge \delta$ , and  $W(x - \beta) = \delta = W(x - \beta).$ 

An immediate consequence of theorem 4.1 is: <u>Theorem 4.2:</u> If  $W = [W_0, W(x - \alpha) = \mathcal{F}]$  is any inductive value of A[x] with  $W_0 \alpha \ge \mathcal{F}$ , then  $V = [V_0, V_1 x = \mathcal{F}]$  is the reduction of W to R[x] providing  $V_0 = W_0$  on R.

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<u>Theorem 4.3</u>: Let  $W = \begin{bmatrix} W_0, W(x - \alpha) = \delta \end{bmatrix}$  be an inductive value of A[x] with  $\delta > W_0 \alpha$ , then  $V_1 = \begin{bmatrix} V_0, V_1 x = W_0 \alpha \end{bmatrix}$ , where  $V_0 = W_0$  on R, is the first-stage of the reduction of W to R[x]. There exist polynomials  $f(x) \in R[x]$  such that  $V_1f(x) < Wf(x)$ . <u>Proof</u>: The value of x is  $Wx = \min \{\delta, W_0 \alpha\} = W_0 \alpha$ . Therefore  $W_1 = \begin{bmatrix} W_0, W_1 x = W_0 \alpha \end{bmatrix}$  is a first-stage value to W;  $W_1 \le W$ . Hence  $V_1 = \begin{bmatrix} V_0, V_1 x = W_0 \alpha \end{bmatrix}$  is the first-stage of the reduction to R[x]. Let  $G(x) = (x - \alpha)(x - \beta_1) \dots (x - \beta_t)$ be the minimal polynomial of  $\alpha$  in R[x]. Since  $W_1(x - \alpha) < W(x - \alpha)$  and  $W_1(x - \beta_1) \le W(x - \beta_1)$ ,  $V_1G(x) = W_1G(x) < WG(x)$ .

Theorem 4.3 shows that for  $\mathcal{V} > W_0 \alpha$  at least one more key is necessary to obtain the correct reduction of W to R[x]. Lemma 4.4: Let  $W = [W_0, W(x - \alpha) = \mathcal{V}]$  be any inductive value of A [x]. A polynomial  $f(x) \in A [x]$  is equivalence divisible by  $x - \alpha$  in W if and only if  $W_0 f(\alpha) > Wf(x)$ . Proof: Let  $f(x) = f_n(x - \alpha)^n + f_{n-1}(x - \alpha)^{n-1} + \dots + f_0$ be the expansion of f(x) in powers of  $x - \alpha$ ;  $f_1 \in A$ . Since  $Wf(x) = \min \{ W_0 f_1 + i \mathcal{V} \}, W_0 f_0 \ge Wf(x)$ ; and because  $f_0 = f(\alpha)$ , the relation  $W_0 f(\alpha) \ge Wf(x)$  always holds. Suppose  $W_0 f(\alpha) > Wf(x)$ . Then  $W [f(x) - \{ f_n(x - \alpha)^n + \dots + f_1(x - \alpha) \} ]$  $= W_0 f(\alpha) > Wf(x)$  and, therefore,  $f(x) \sim f_n(x - \alpha)^n + \dots + f_1(x - \alpha) \}$ in W; that is,  $x - \alpha | f(x)$  in W. Suppose, now,  $f(x) \sim q(x)(x - \alpha)$  in W. Then  $f(x) = q(x)(x - \alpha) + h(x)$ , where Wh(x) > Wf(x). But, since  $h_0$ , the last term in the expansion of h(x), is  $f_0$  it follows that

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 $W_{o}f_{o} = W_{o}h_{o} \ge Wh(x) > Wf(x)$ ; that is,  $W_{o}f(\alpha) > Wf(x)$ .

In the results to follow the polynomials.  $\varphi_i$  and the real numbers  $\mu_i$  will be the homogeneous key polynomials and their values which are used by MacLane to extend a value  $V_0$  on R to a value  $W_0$  on  $R(\alpha)$  (§3). Since the value  $W_0$  on A, and therefore on  $R(\alpha)$ , is given,  $V_0$  will be the restriction of  $W_0$  to R. However, there exist  $\varphi_i$  and  $\mu_i$  defining the extension of this  $V_0$  to the given  $W_0$  on  $R(\alpha)$ . <u>Theorem 4.5</u>: Let the polynomials  $\varphi_i$  and the numbers  $\mu_i$  be the keys and values which define the extension of  $V_0$  on R to  $W_0$  on  $R(\alpha)$ . The kath stage of the reduction of  $W = \left[W_0, W(x - \alpha) = \delta\right]$  to R[x] is given by  $V_k = \left[V_0, V_1 x = \mu_1, V_2 \varphi_2 = \mu_2, \dots, V_k \varphi_k = \mu_k\right],$ 

provided that  $x - \alpha | \varphi_i$  in W is false for all i in the interval  $1 \le i \le k$ .

<u>Proof:</u> By Theorem 4.3 and lemma 4.4 this theorem is true for k = 1. Suppose the result is true up to k - 1, then  $\nabla_{k-1} = \begin{bmatrix} \nabla_0, \nabla_1 x = \mu_1, \dots, \nabla_{k-1} \phi_{k-1} = \mu_{k-1} \end{bmatrix}$  is the (k-1)st stage of the reduced value and  $\nabla_{k-1} f(x) \leq W f(x)$  for all  $f(x) \in \mathbb{R}[x]$ . Let  $f(x) \in \mathbb{R}[x]$  be any polynomial such that deg  $f(x) < \deg \phi_k$ . Then  $\nabla_{k-1} f(x) = W_0 f(\alpha)$ , because in extending  $\nabla_0$  on  $\mathbb{R}$  to  $W_0$  on  $\mathbb{R}(\alpha)$  the value  $W_0 f(\alpha)$  is defined to be  $\nabla_{k-1} f(x)$  when deg  $f(x) < \deg \phi_k$ . But, since  $W f(x) \geq \nabla_{k-1} f(x) = W_0 f(\alpha) \geq W f(x)$  by lemma 4.4, it is concluded that  $W f(x) = \nabla_{k-1} f(x)$  for all  $f(x) \in \mathbb{R}[x]$  such that deg  $f(x) < \deg \phi_k$ . Høwever,  $W \phi_k = W_0 \phi_k (\alpha) = \mu_k > \nabla_{k-1} \phi_k$ , since  $\mathbf{x} - \alpha | \varphi_k$  in W is false, and  $\varphi_k$  may then be chosen the next key over  $V_{k-1}$  with value  $V_k \varphi_k = \mu_k$  (Cf. introduction). Also, the value

 $\begin{array}{l} \mathbb{V}_k = \left[ \mathbb{V}_0, \ \mathbb{V}_1 \mathbb{x} = \mu_1, \ \cdots, \ \mathbb{V}_{k-1} \varphi_{k-1} = \mu_{k-1}, \ \mathbb{V}_k \varphi_k = \mu_k \right] \\ \text{satisfies the relation } \mathbb{V}_k f(\mathbf{x}) \leq \mathbb{W} f(\mathbf{x}) \text{ for all } f(\mathbf{x}) \in \mathbb{R}[\mathbf{x}] \text{ and} \\ \text{is, therefore, a k-th stage of the restriction of } \mathbb{W} \text{ to } \mathbb{R}[\mathbf{x}]. \\ \underline{\text{Lemma } 4.6} \colon \text{ If } \varphi_1 \text{ are the keys used to extend } \mathbb{V}_0 \text{ on } \mathbb{R} \text{ to} \\ \mathbb{W}_0 \text{ on } \mathbb{R}(\alpha) \text{ , then there exists an } \text{ i such that } \mathbf{x} - \alpha \right] \varphi_1 \\ \text{ in } \mathbb{W} = \left[ \mathbb{W}_0, \ \mathbb{W}(\mathbf{x} - \alpha) = \mathbf{V} \right]. \end{array}$ 

<u>Proof</u>: There are two cases to consider:

- (a) W<sub>0</sub> is found by an inductive value then the last key is G(x), the minimal polynomial for α in R[x], which is divisible by x α and, therefore equivalence divisible by x α in W,
  - (b)  $W_0$  is found by a limit value if there exists no i such that  $x - \alpha | \phi_i$  in W, then by theorem 4.5 every  $\phi_i$ , with value  $\mu_i$ , occurs in the reduction of W to R[x]. But this implies the value of G(x) is  $+\infty$ ; while  $WG(x) < +\infty$ .

Lemma 4.6 implies the existence of a first key  $\varphi_{k+1}$  which is equivalence divisible by  $x - \alpha$  in W. By theorem 4.5,

$$V_{k} = [V_{o}, V_{1}x = \mu_{1}, \dots, V_{k}\phi_{k} = \mu_{k}]$$

is the k-th stage of the reduction of W to R [x]. There are two possibilities for the k-th stage value of  $\varphi_{k+1}$ , either

 $\begin{array}{l} \mathbb{V}_{k} \varphi_{k+1} < \mathbb{W}_{\phi_{k+1}} \quad \text{or} \quad \mathbb{V}_{k} \varphi_{k+1} = \mathbb{W} \varphi_{k+1} \\ \hline \\ \underline{\text{Theorem 4.7}}: \quad \text{Let} \quad \mathbb{W} = \left[ \mathbb{W}_{0}, \ \mathbb{W}(\mathbf{x}-\alpha) = \mathcal{F} \right] \quad \text{, with} \quad \mathcal{F} > \mathbb{W}_{0} \alpha \quad \text{, be} \end{array}$ 

given on A[x]. Let  $\{V_i\}$  be the sequence of approximants to G(x), the minimal polynomial of  $\alpha$ , defining the extension of  $V_{\phi}$  on R to  $W_{\phi}$  on R( $\alpha$ ). If  $\varphi_{k+1}$  is the first key in these approximants such that  $x - \alpha | \varphi_{k+1}(W)$ , then the reduction of W to R[x] is given by:

(1) 
$$V = \begin{bmatrix} V_0, V_{1x} = \mu_1, \dots, V_k \phi_k = \mu_k, V \phi_{k+1} = W \phi_{k+1} \end{bmatrix}$$
  
when  $V_k \phi_{k+1} < W \phi_{k+1}$ ,  
(2)  $V_k = \begin{bmatrix} V_0, V_{1x} = \mu_1, \dots, V_k \phi_k = \mu_k \end{bmatrix}$  when  
 $V_k \phi_{k+1} = W \phi_{k+1}$ .

<u>Proof</u>: As in theorem 4.5,  $\nabla_k f(x) = W_0 f(\alpha) = Wf(x)$  for all  $f(x) \in R[x]$  such that deg  $f(x) < \deg \varphi_{k+1}$ . Since in (1)  $\nabla_k \varphi_{k+1} < W \varphi_{k+1}$ ,  $\varphi_{k+1}$  may be chosen as the next key with value  $W \varphi_{k+1}$ ; this gives  $V \le W$  on R[x]. If  $\nabla_k$  in (1) and  $M_k$ d in (2) are both denoted by  $V^*$ , the two results may be given by one proof. Suppose the existence of

$$f(x) = f_{n}(x) \varphi_{k+1}^{n} + f_{n-1}(x) \varphi_{k+1}^{n-1} + \dots + f_{\varphi}(x)$$
  
=  $f_{\varphi}(x) - g(x) \varphi_{k+1}$ ,

where deg  $f_1(x) < \deg \varphi_{k+1}$  and  $f(x) \in R[x]$ , a monic polynomial of minimum degree such that V'f(x) < Wf(x). Then  $V'f_O(x) \ge V'f(x)$ . If  $V' = V_k$ , this is immediate from the definition of Vf(x). If  $V' = V_k$ , then  $\varphi_{k+1}$  is a key polynomial over  $V_k$  since it defines an approximant  $V_{k+1}$ . Suppose  $V_k f_O(x) < V_k f(x)$ , then  $V_k [f_O(x) - g(x) \varphi_{k+1}] = V_k f(x) > V_k f_O(x)$  and  $f_O(x) \sim g(x) \varphi_{k+1}(V_k)$ , which contradicts the minimal condition of the key  $\varphi_{k+1}$ . Now, V'g(x) = Wg(x) because deg  $g(x) < \deg f(x)$ ; therefore  $V'[f(x) = f_O(x)] = V((-g(x)) + V'\varphi_{k+1} = W(-g(x)) + W\varphi_{k+1} = W[f(x)-f_O(x)]$ That  $V(f_O(x) = V'f(x)$  may be seen by assuming  $V'f_O(x) > V'f(x)$ ;

for then

On combining theorems 4.2 and 4.7 and lemma 4.4 a picturesque description of the reduction can be given in terms of the size  $\mathcal{F}$ . For this purpose consider  $W = \left[ W_0, W(x - \alpha) = \mathcal{F} \right]$ A[x] with  $\gamma' < W_0 \alpha$  and examine the reduction to R[x] as  $\mathcal{V}$  continuously increases. For  $\mathcal{V} < W_0 \alpha$  the reduction is  $V_1 = [V_{\alpha}, V_1 = \delta]$ . As  $\delta$  increases to  $W_{\alpha} = \mu_1$  the reduction increases to  $V_1 = [V_0, V_1 x = \mu_1]$ . When  $\mathcal{X}$  is just larger than  $W_0\alpha$ , the key  $\varphi_2$  of the second approximant to G(x), the minimal polynomial of  $\alpha$ , is needed. The reduction is  $V_2^{\prime} = \left[ V_0, V_1 x = \mu_1, V_2^{\prime} \varphi_2 = \mu \right]$ . As  $\mathcal{X}$  increases again the value  $V_2^{*}\phi_2 = \mu$  increases to  $W\phi_2 = W_0\phi_2(\alpha) = \mu_2$  and the reduction increases to  $V_2 = [V_0, V_1x = \mu_1, V_2\varphi_2 = \mu_2]$ . If this process is continued it is seen that as  $\mathcal X$  increases the reduction sweeps through the approximants to G(x) which describe the extension of  $V_O$  to  $W_O$  on  $R(\alpha)$  . The only difference between the reduction and the corresponding approximant defined by the same keys is that the value assigned to the last key in the reduction may be less than its value in the approximant. But as  $\mathcal{X}$  is increased this value will increase to the corresponding approximant value if the last key is not G(x) . For then  $\mathcal{Y}$  would have to increase to  $\infty$  It will now be shown what this reduction is when W is a limit value. In the remainder of this section W is defined by

 $W = \begin{bmatrix} W_0, & W_1(x - \alpha_1) = \mathscr{V}_1, \dots, & W_1(x - \alpha_1) = \mathscr{V}_1, \dots \end{bmatrix},$ where the pseudo-convergent sequence  $\{\alpha_i\}$  has no pseudo-limit in A. It should be noticed that

 $W_{i} = \begin{bmatrix} W_{0}, W_{i}(x - \alpha_{i}) = \mathscr{V}_{1}, \dots, W_{i}(x - \alpha_{i}) = \mathscr{V}_{i} \end{bmatrix}$ can also be represented by  $W_{i} = \begin{bmatrix} W_{0}, W_{i}(x - \alpha_{i}) = \mathscr{V}_{i} \end{bmatrix}$ . For, to find  $W_{i}f(x)$ , f(x) is expanded in powers of  $x - \alpha_{i}$  and the coefficients are valued with  $W_{i-1}$ . But the coefficients are in A and are therefore actually valued by  $W_{0}$ .

Lemma 4.8: The value  $W_i = W_{i+1} = \begin{bmatrix} W_0, W_{i+1}'(x - \alpha_{i+1}) = Y_i \end{bmatrix}$ on A[x].

Lemma 4.9: For each  $x - \beta \in A[x]$  there exists a positive integer N such that  $W(x - \beta) = W_i(x - \beta)$  for all  $i \ge N$ . Proof: If no such N exists, then  $W(x - \beta) > W_i(x - \beta)$  for all  $i \ge N$ . all i > 0 since  $W_i \le W_{i+1}$ . Now  $W_0(\alpha_{i+1} - \beta) > W_i(x - \beta)$  for otherwise from  $x - \beta = (x - \alpha_{i+1}) + (\alpha_{i+1} - \beta)$  it follows that  $\delta_i = W_i(x - \alpha_{i+1}) \ge W_0(\alpha_{i+1} - \beta)$ ; and therefore  $\delta_{i+1} > W_0(\alpha_{i+1} - \beta)$ . Hence,  $W(x - \beta) = W_0(\alpha_{i+1} - \beta)$  $= W_{i+1}(x - \beta)$ , which contradicts  $W(x - \beta) \ge W_i(x - \beta)$  for all i > 0. Since  $W_0(\alpha_{i+1} - \beta) \ge W_i(x - \beta) = W_i(x - \alpha_{i+1}) = \delta_i$ , therefore  $W_{o}(\alpha_{i+1} - 3) = W_{o}\left[(\alpha_{i+1} - \alpha_{i+2}) + (\alpha_{i+2} - 3)\right] = \delta_{i+1}$ for all i > 0. But this implies 3 is a pseudo-limit of  $\{\alpha_i\}$ .

<u>Theorem 4.10</u>: The k(i) keys occurring in the reduction of  $W_i = \left[W_s, W_i(x - \alpha_i) = Y_i\right]$  on A[x] to R [x] are the first k(i) keys in the reduction of

 $W = \begin{bmatrix} W_0, W_1(x - \alpha_1) = Y_1, \dots, W_i(x - \alpha_i) = Y_i, \cdots \end{bmatrix}$ on A[x] to R [x]. Also, the values  $\mu_v$ , for  $1 \le v \le k(i)$ , in the reduction of  $W_i$  are the first k(i) - 1 values for the keys in the reduction of W.

<u>Proof:</u> By lemma 4.9, for any given  $f(x) \in \mathbb{R}[x]$  there exists N such that  $Wf(x) = W_if(x)$  for all  $i \ge N$ . If  $V_{k(i)}$ an is the reduction of  $W_i$ , then  $V_{k(i)} f(x) = Wf(x)$  for all  $i \ge N$ . Hence, the sequence of values  $\{V_{k(i)}\}$  on R [x] gives every polynomial in R[x] the correct W value. It is only necessary then to show that the k(i) keys in the reduction of  $W_i$  are the first k(i) keys in the reduction of  $W_{i+i}$  and that the values, with the possible exception of  $\mu_k(i)$ , are the same. By lemma 4.8,  $W_i$ and  $W_{i+1} = \left[ W_0, W_{i+1}(x - \alpha_{i+1}) = Y_i \right]$  define the same valuation of A[x]. Hence, they will have the same reduction on R[x] and, because the keys in the reductions are homogeneous, each reduction will be identical with respect to keys and values ([2], theorem 16.4) . As the value of  $x - \alpha_{i+1}$  is increased from  $\mathscr{S}_i$ to  $\mathcal{F}_{i+1}$  the value:  $\mu_{k(i)}$  might increase and the keys, if any, appearing in  $V_{k(i+1)}$  but not in  $V_{k(i)}$  are used to augment  $V_{k(i)}$  to  $V_{k(i+1)}$ . These are the only changes that can happen; and at least one of these changes must happen. The truth of this follows from the discussion immediately after theorem 4.7

and the fact that the minimal polynomial of  $\alpha_{i+1}$  definitely increases as the value of  $x - \alpha_{i+1}$  increases. <u>Theorem 4.11</u>: The reduction of a limit value W on A[x] to R[x], as described in theorem 4.10, is a limit value. <u>Proof</u>: Suppose the reduction is an inductive value  $V_k = \begin{bmatrix} V_o, V_1 x = \mu_1, \dots, V_k \phi_k = \mu_k \end{bmatrix}$ . By lemma 4.9 there exists a smallest i such that  $W_i \phi_k = W \phi_k = \mu_k$ . The reduction of  $W_i$  to R[x] must be  $V_k$  since this is the first stage in which  $\phi_k$  assumes the value  $\mu_k$  and  $W_i \leq W$ . But for G(x), the minimal polynomial of  $\alpha_{j+1}$ , in R[x],  $V_k G(x) = W_i G(x) < W_{i+1} G(x)$  $\leq WG(x)$ . Since by assumption  $V_k = W$ , this contradiction establishes the theorem.

5. In section 4 the connection between a value W of A[x] and its reduction to R[x] was established. The converse problem will now be solved; that is, given a value V of R[x], to find an extension of V to A[x]. First, however, it will be shown that a value W on A[x] may be written in a standard form.

In the following theorem the notation

 $V = \begin{bmatrix} V_0, \ V_1 x = \mu_1, \ \cdots, \ V_k \phi_k = \mu_k, \end{bmatrix}$  is to mean that a least the keys up to  $\phi_k$  occur in the representation of V; however, V may be an inductive value with keys past  $\phi_k$ , or V may even be a limit value. This notation will also be used for

 $W = \left[ W_0, W_1 x = \mu_1, \dots, W_k (x - \alpha_k) = \mathcal{X}_k, \dots \right]_{\mathcal{X}_k}$ Lemma 5.1: If a value W' of A[x] reduces to

 $V = V_0, V_1 x = \mu_1, \dots, V_k \phi_k = \mu_k,$ on R[x], then V may be extended to \_24.

$$W = \begin{bmatrix} W_0, & W_1 x = \mu_1, & W_2(x - \alpha_2) = \mathcal{V}_2, & \dots, & W_k(x - \alpha_k) \neq \mathcal{V}_k, \end{bmatrix}$$
on A [x] where:

- (1)  $\varphi_i(\alpha_i) = 0$  for i = 1, 2, ..., k, (2)  $Y_i = W^i(x - \alpha_i)$  for i = 1, 2, ..., k,
  - (3)  $W_0/3 = W'/3$  for all  $/3 \epsilon A$ ,
- (4)  $W'f(x) \ge Wf(x)$  for all  $f(x) \in A[x]$ .

Further, the reduction of  $W_i = \begin{bmatrix} W_0, W_i(x - \alpha_i) = \aleph_i \end{bmatrix}$ is  $V_i = \begin{bmatrix} V_0, V_{1x} = \mu_1, \dots, V_i \varphi_i = \mu_i \end{bmatrix}$  for  $i = 1, 2, \dots, k$ . <u>Proof</u>: Let  $W_0$  on A be defined by  $W_0 / 3 = W' / 3$  for  $/ 3 \in A$  and  $W_1$  defined by  $W_1 = \begin{bmatrix} W_0, W_{1x} = \mu_1 \end{bmatrix}$ . The reduction of  $W_i$  is certainly  $V_1 = \begin{bmatrix} V_0, V_{1x} = \mu_1 \end{bmatrix}$ . Let  $f(x) = \sum_i a_i x^i$ ,  $f(x) \in A[x]$ ,

then  $W'f(x) \ge \min \{W_0a_1 + iW'x\} = W_1f(x)$ . The value  $W_1$ satisfies (2), (3), (4) and also (1) since  $\varphi_1 = x$  and  $\alpha_1 = 0$ . Assume theorem true up to  $W_{k-1}$ . Then  $W_{k-1}$  reduces to  $V_{k-1}$ and  $W'\varphi_k > V_{k-1}\varphi_k = W_{k-1}\varphi_k$ . For a key to augment  $W_{k-1}$  let a factor  $x - \alpha_k$  of  $\varphi_k$  be chosen so that  $W'(x-\alpha_k) > W_{k-1}(x-\alpha_k)$ (Cf. introduction) and  $W(x - \alpha_k) \ge W'(x - \beta)$  for any factor  $x - \beta$  of  $\varphi_k$  such that  $W'(x - \beta) > W_{k-1}(x - \beta)$ . Now define  $W_k$  by  $W_k = \left[W_{k-1}, W_k(x - \alpha_k) = \mathcal{K}_k = W'(x - \alpha_k)\right]$ . For  $x - \beta$ , any factor of  $\varphi_k$ ,

$$W_{k}(x - \beta) = \min \{ \mathcal{V}_{k}, W_{0}(\alpha_{k} - \beta) \}, \text{ and}$$
$$W^{\dagger}(x - \beta) \geq \min \{ \mathcal{V}_{k}, W_{0}(\alpha_{k} - \beta) \}.$$

The inequality cannot hold; for, then,

$$\begin{split} \mathbb{W}(\mathbf{x} - \mathbf{\beta}) > \mathbb{V}_{k}^{t} &= \mathbb{W}_{k}(\mathbf{x} - \alpha_{k}) \geq \mathbb{W}_{k}(\mathbf{x} - \mathbf{\beta}) \geq \mathbb{W}_{k-1}(\mathbf{x} - \mathbf{\beta}) ,\\ \text{which contradicts the choice of } \mathbf{x} - \alpha_{k}. \text{ Therefore} \\ \mathbb{W}^{t}(\mathbf{x} - \mathbf{\beta}) &= \mathbb{W}_{k}(\mathbf{x} - \mathbf{\beta}) \text{ for all factors of } \varphi_{k} ; \text{ so} \\ \mathbb{W}^{t}\varphi_{k} &= \mathbb{W}_{k}\varphi_{k} . \text{ Since } \varphi_{k}(\alpha_{k}) = 0 , \text{ certainly } \mathbf{x} - \alpha_{k} | \varphi_{k} \text{ in } \mathbb{W}_{k} . \end{split}$$

This means the reduction will use only keys  $\varphi_i$  for  $i \leq k$ . But  $W_k \varphi_k = W' \varphi_k = \mu_k > W_{k-1} \varphi_k$ , hence the reduction of  $W_k$ must be  $V_k$ .

<u>Theorem 5.2</u>: If a value  $W' = [W_0, W'(x - 3) = \delta]$  on A[x] reduces to  $V_k = [V_0, V_{1x} = \mu_1, \dots, V_k \phi_k = \mu_k]$  on R[x], then there exists an  $\alpha$ , such that  $\phi_k(\alpha) = 0$ , and a  $\delta$  so that  $W' = W = [W_0, W(x - \alpha) = \delta]$ . Also  $\delta = \delta$ . <u>Proof</u>: Let  $V_k$  be extended to W on A[x] as in theorem 5.1, where  $\alpha = \alpha_k$ , then  $W \le W'$ . If there exists an  $x - \Theta \in A[x]$ such that  $W(x - \Theta) < W'(x - \Theta)$ , then for  $G(x) \in R[x]$ , where  $G(\Theta) = 0$ , it follows that

$$\begin{split} \mathbb{V}_k \mathbb{G}(\mathbf{x}) &= \mathbb{W}\mathbb{G}(\mathbf{x}) < \mathbb{W}^*\mathbb{G}(\mathbf{x}) = \mathbb{V}_k\mathbb{G}(\mathbf{x}); \quad \text{hence,} \quad \mathbb{W} = \mathbb{W}^* \text{ .} \\ \text{Also, since} \quad & \mathbf{\delta} = \mathbb{W}(\mathbf{x} - \alpha) = \mathbb{W}^*(\mathbf{x} - \alpha) \leq \delta \quad \text{and} \\ \delta &= \mathbb{W}^*(\mathbf{x} - \mathbf{\beta}) = \mathbb{W}(\mathbf{x} - \mathbf{\beta}) \leq \mathbf{\delta}, \end{split}$$

then  $\mathcal{F} = \delta$ .

<u>Theorem 5.3</u>: Let  $W' = \left[W_0, W_1(x - \beta_1) = \delta_1, \dots, W_1(x - \beta_1) = \delta_2, \dots\right]$ be a value on A[x] which reduces to

$$V = \left[V_0, V_1 x = \mu_1, \dots, V_k \varphi_k = \mu_k, \dots\right] \text{ on } \mathbb{R}[x],$$

then W' may be represented by

 $W = \begin{bmatrix} W_0, & W_1 x = \mu_1, & W_2(x - \alpha_2) = \mathscr{V}_2, & \dots, & W_k(x - \alpha_k) = \mathscr{V}_k & \dots \end{bmatrix}$ where:

(1) 
$$\varphi_k(\alpha_k) = 0$$
 for all  $k \ge 1$ 

(2) reduction of  $W_k$  to R[x] is  $V_k$ ,

<u>Proof</u>: The proof is similar to that for theorem 5.2.

From theroems 5.2 and 5.3 it is seen that every valuation of A(x) may be put into a form such that each  $a_k$  is a root of the

corresponding key  $\varphi_k$  appearing in the reduction of the valuation to R [x]. This information indicates how a valuation V of R[x] may be extended to some valuation of A[x]. It will now be shown how this extension can be accomplished. Lemma 5.4: Let  $\alpha$  be a root of some polynomial  $\varphi \in A[x]$ , then a valuation  $W = \begin{bmatrix} W_0, W(x - \alpha) = Y \end{bmatrix}$  can be defined on A[x] such that  $\varphi$  has a prescribed value  $\mu$ . The value  $\lambda$ is uniquely determined. <u>Proof</u>: Let  $\varphi = \beta_n(x-\alpha)^n + \beta_{n-1}(x-\alpha)^{n-1} + \ldots + \beta_1(x-\alpha)$ , where  $\beta_i \in A$ , then  $W_{\phi} = \min \{ W_{\phi} | \beta_i + i \}$ . Let the numbers  $\gamma_i$ be defined by  $W_0/3_i + i Y_i = \mu$  for i = 1, 2, ..., m and Ydefined by  $\delta = \max \delta_i$ . For this value of  $\delta$ ,  $\mathbb{W}_{0}/\mathbf{i} + \mathbf{i} \mathbf{i} \geq \mu$  for  $\mathbf{i} = 1, 2, ..., n$  and the equality holds for at least one value of i. Suppose there were two values  $\gamma$ and  $\mathcal{F}'$ , with  $\mathcal{F} > \mathcal{F}'$ , with the desired property. Since  $\varphi(\alpha) = 0$ , there exists an  $i \neq 0$  such that

$$u = W_0 / i + i \rangle > W_0 / i + i \rangle'.$$

So W defined by  $\mathcal{F}'$  would give  $\varphi$  a value  $W\varphi < \mu$ . There-fore the value  $W(x - \alpha)$  is unique.

Lemma 5.5: In an inductive value

 $V_k = \left[V_0, V_1 x = \mu_1, \dots, V_{k-1} \varphi_{k-1} = \mu_{k-1} V_k \varphi_k = \mu_k\right]$  of R[x] the  $V_1$  for  $1 \le i < k$  are the complete, and only, set of approximants to  $\varphi_k$ .

<u>Proof</u>: This follows immediately from [3], theorem 5.3 (Cf. end of  $\S$ 3) if V is defined by

 $\mathbb{V} = \left[ \mathbb{V}_{o}, \ \mathbb{V}_{1} \mathbf{x} = \mu_{1}, \ \cdots, \ \mathbb{V}_{k-1} \varphi_{k-1} = \mu_{k-1} \mathbf{v} \varphi_{k} = \infty \right].$ 

Theorem 5.6: Let  $W_0$  on A be an extension of  $V_0$  on R, then  $V_k = [V_0, V_{1x} = \mu_1, \dots, V_k \varphi_k = \mu_k]$  on R[x] may be extended to  $W = [W_0, W(x - \alpha) = \mathcal{V}]$  on A[x], where  $\varphi_k(\alpha) = 0$ . <u>Proof</u>: By lemma 5.4 the value  $W(x - \alpha) = \mathcal{V}$  is uniquely determined from  $W\varphi_k = \mu_k$ . Now, W defined by this  $\mathcal{V}$  reduces to

 $\begin{array}{l} V_{l} = \left[ V_{0}, \ V_{l} x = \mu_{l}, \ \cdots, \ V_{\ell-1} \varphi_{\ell-1} = \mu_{\ell-1}, \ V_{\ell} \varphi = \nu \right], \\ \text{where } \mathcal{I} \leq k \ \text{, by theorem 4.7 since } \varphi_{k}(\alpha) = 0 \ \text{. But } W \varphi_{k} = \mu_{k} \\ \text{and, so, } V_{l} \ \varphi_{k} = \mu_{k} = V_{k} \varphi_{k} \ \text{. This implies } \mathcal{I} = k \ \text{and } \nu = \mu_{k} \ \text{.} \\ \text{That is, the reduction of } W \ \text{is } V_{k} \text{ or } W \ \text{is the extension of } V_{k} \ \text{.} \end{array}$ 

Since theorem 5.2 claims every valuation W may be defined by some a where  $\varphi_k(\alpha) = 0$ , it is seen that for a given  $W_0$ the maximum number of extensions of  $V_k$  on R[x] to W on A[x]is the degree of  $\varphi_k$ ; and, every extension may be found by the method of theorem 5.6.

<u>Theorem 5.7</u>: Let  $W_0$  on A be an extension of  $V_0$  on R and let  $V = \begin{bmatrix} V_0, V_1 x = \mu_1, \dots, V_k \phi_k = \mu_k, \dots \end{bmatrix}$  be a limit value of R[x]. The value V may be extended to the MacLane value  $W = \begin{bmatrix} W_0, W_1 x = \mu_1, W_2(x - \alpha_2) = V_2, \dots, W_k(x - \alpha_k) = V_k, \dots \end{bmatrix}$ on A[x] where:

(1)  $\varphi_i(\alpha_i) = 0$  for i = 1, 2, ...,

(2) 
$$Y_i$$
 is uniquely determined by the factor  $x - \alpha_i$   
and the value  $\mu_i$  of  $\varphi_i$ ,

(3) 
$$\mathbb{D}_{\varphi}(x - \alpha_i) = 1$$
 in  $x - \alpha_{i-1}$  for all  $i > 1$   
(Cf. definition 3.2),

(4)  $W_i$  reduces to  $V_i$  on R[x] for i = 1, 2, ...<u>Proof</u>: It will be possible to formally construct the sequence of values  $\{W_i\}$  if it can be shown that there always exists a factor

 $\phi_{\rm i}$  satisfying (3). The construction of W<sub>i</sub> with this of factor may be accomplished by the method of theorem 5.6. However, it will be necessary to show that W defined by this sequence of values is actually a MacLane value. In order to prove that property (3) may be satisfied it will be shown that, given a value  $W_k$  which satisfies (1), (2), and (4) a value  $W_{k+1}$  can be defined over  $W_k$  also satisfying (1), (2) and (4) and such that  $D_{\phi}(x - \alpha_{k+1}) = 1$  in  $x - \alpha_k$ . Then, since  $W_1$  satisfies (1), (2) and (4) it follows by induction that property (3) can be satisfied for all i > 1. Let  $W_k$  satisfy (1), (2) and (4) and let  $W_k^i$  be defined by  $W_k^i = \left[W_0, W_k^i(x - \alpha_k) = \delta\right]$  with  $\pmb{\mathcal{X}}<\pmb{\mathcal{Y}}_k$  but such that  $\forall_{k=1}\phi_k<\forall_k^{\dagger}\phi_k=\mu<\mu_k$  . Then as  $\delta \rightarrow \delta_k$ ,  $\mu \rightarrow \mu_k$ . Therefore, if  $\varphi_{k+1}$  is expanded in powers of  $\varphi_k$ , as  $\gamma \neq \gamma_k$  the value of  $\varphi_{k+1}$  must continuously increase, since proj  $V_k > 0$  with respect to  $\varphi_{k+1}$ . So, if  $\varphi_{k+1}$ is expanded in powers of  $x - \alpha_k$ , the value of  $\varphi_{k+1}$  must also continuously increase as  $\ensuremath{\mathcal{X}} \ensuremath{\rightarrow} \ensuremath{\mathcal{X}}_k$  . This could only happen if  $D_{\phi}\phi_{k+1} \ge 1$  in  $x - \alpha_k$ . Hence, there exists a factor  $x - \alpha_{k+1}$ of  $\varphi_{k+1}$  such that  $D_{\varphi}(x - \alpha_{k+1}) = 1$  in  $x - \alpha_k$ . Now  $W_{k+1}$ may be defined by  $W_{k+1} = \begin{bmatrix} W_k, W_{k+1}(x - \alpha_{k+1}) = Y_{k+1} \end{bmatrix}$  where  $W_{k+1}$ satisfies (1), (2), (3) and (4). It only remains to show that W defined by this sequence of values is a MacLane value. Since  $D_{\varphi}(x - \alpha_{k+1}) = 1$  in  $x - \alpha_k$ ,  $\mathcal{Y}_k = W_k(x - \alpha_k) = W_k(x - \alpha_{k+1})$  $\leq W_0(\alpha_k - \alpha_{k+1})$ ;

that is,  $\mathscr{F}_k \leq W_0(\alpha_k - \alpha_{k+1})$ . Let  $W_{k+1}^i$  be defined by  $W_{k+1}^i = \left[W_0, W_{k+1}^i(x - \alpha_{k+1}) = \mathscr{F}_k\right]$ ,

then  $W_k = W_{k+1}^{\dagger}$ ; for, let  $x - /3 \in A[x]$ , then

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$$\begin{array}{l} & W_k(x-\beta) = \min\left\{ \mathscr{Y}_k, \ W_0(\alpha_k-\beta) \right\} \quad \text{and} \\ & W_{k+1}'(x-\beta) = \min\left\{ \mathscr{Y}_k, \ W_0(\alpha_{k+1}-\beta) \right\} \quad \text{but} \\ & W_0(\alpha_{k+1}-\beta) \geq \min\left\{ W_0(\alpha_{k+1}-\alpha_k), \ W_0(\alpha_k-\beta) \right\} \geq \min\left\{ \mathscr{Y}_k, \ W_0(\alpha_k-\beta) \right\} \\ & \text{Therefore } W_k \quad \text{and} \quad W_{k+1}' \quad \text{have the same reduction } V_k \quad \text{on } \mathbb{R}[x] \\ & \text{Hence,} \quad \mathscr{Y}_{k+1} \geq \mathscr{Y}_k' \quad \text{in order that } W_{k+1} \quad \text{reduce to } V_{k+1} \\ & \text{Let} \quad \Psi_k \quad \text{have the expansion} \end{array}$$

$$\varphi_{k} = f_{n}(x - \alpha_{k})^{n} + f_{n-1}(x - \alpha_{k})^{n-1} + \dots + f_{1}(x - \alpha_{k}) \text{,then}$$

$$W_{k}\varphi_{k} = \min_{i} \left\{ W_{0}f_{i} + i \bigotimes_{k} \right\} \text{ and}$$

$$W_{k+1}\varphi_{k} \geq \min_{i} \left\{ W_{0}f_{i} + iW_{k+1}(x - \alpha_{k}) \right\}.$$

Let i, which cannot be zero, be chosen such that a minimum term is actually obtained in the second inequality; then

 $W_{o}f_{i} + i \bigvee_{k} \geq W_{k}\phi_{k} = W_{k+1}\phi_{k} \geq W_{o}f_{i} + iW_{k+1}(x - \alpha_{k}) .$ Therefore, since  $\mathbf{1} \neq 0$ ,  $\bigvee_{k} \geq W_{k+1}(\mathbf{x} - \alpha_{k}) .$  Now,

$$\begin{split} & \mathbb{W}_{o}(\alpha_{k} - \alpha_{k+1}) \geq \min \left\{ \mathbb{W}_{k+1}(x - \alpha_{k+1}), \ \mathbb{W}_{k+1}(x - \alpha_{k}) \right\}, \text{ but} \\ & \mathbb{W}_{k+1}(x - \alpha_{k+1}) = \mathcal{Y}_{k+1} > \mathcal{Y}_{k} \geq \mathbb{W}_{k+1}(x - \alpha_{k}), \text{ and, so,} \\ & \mathbb{W}_{o}(\alpha_{k} - \alpha_{k+1}) = \mathbb{W}_{k+1}(x - \alpha_{k}) \leq \mathcal{Y}_{k}. \end{split}$$

But it is known from above that  $W_0(a_k - a_{k+1}) \ge \delta'_k$ . Therefore  $W_0(a_k - a_{k+1}) = \delta'_k$  and since  $\delta'_{k+1} \ge \delta'_k$  for all  $k \ge 1$  the sequence  $\{a_1\}$  is pseudo-convergent. The sequence has no pseudohimit since W reduces to a limit value; otherwise V would be an inductive value ([2], theorem 16.4). There can be no limit for the sequence  $\{a_1\}$  in A since V is a finite value; alternatively, every limit of pseudo-convergent sequence is a pseudolimit. Therefore, W is a MacLane valuation of A [x] satisfying properties (1), (2), (3) and (44). Because of theorem 5.3 and since for every limit value W of A[x],  $D_{\phi}(x - \alpha_{k+1}) = 1$  in  $x - \alpha_k$ , every extension of a limit value V of R[x] to A [x] may be found as in Theorem 5.7.

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