BING'S DOGBONE SPACE AND CURTIS' CONJECTURE

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M.A., University of British Columbia, 1966

A THESIS SUBMITTED IN PARTIAL FULFILMENT OF
THE REQUIREMENTS FOR THE DEGREE OF
DOCTOR OF PHILOSOPHY

in the Department

of

MATHEMATICS

We accept this thesis as conforming to the required standard

THE UNIVERSITY OF BRITISH COLUMBIA

April, 1973

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ABSTRACT

Bing's dogbone space $\mathcal D$ is an upper semi continuous decomposition space of $\mathbf E^3$ which fails to be $\mathbf E^3$ although the associated decomposition consists only of points and tame arcs. It has proved difficult to find topological properties of $\mathcal D$ which distinguish it from $\mathbf E^3$. In this paper, we prove a conjecture of Morton Curtis in 1961 that certain points of $\mathcal D$ fail to possess small simply connected neighbourhoods.

I wish to acknowledge my gratitude to my supervisor Dr. Whittaker for his unselfish and often indispensible aid during my graduate studies at UBC, and to Dr. Luft for his support and enthusiasm. I am grateful also for some conversations and a blizzard of letters from R. H. Bing.

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INTRODUCTION

Bing's dogbone space (which is denoted by \mathcal{D} in this paper) is a decomposition space of E^3 which fails to be homeomorphic to E^3 even though the associated decomposition space is upper semicontinuous and point-like, and each element of the decomposition is either a point or a tame arc. The appearance of $\mathcal D$ in [12] caused some surprise since it was thought at the time that all usc point-like decomposition spaces of E^3 would turn out to be E^3 . Although $\mathcal D$ dates from 1955 and has become rather well-known, it has been found hard to determine those topological properties of the space which distinguish it from $\ensuremath{\text{E}}^3$. Bing's original paper [12] showed that $\,\mathcal{D}\,$ is a non-manifold; but $\,\mathcal{D}\,$ is a simply connected homology manifold and locally simply connected. This paper contains a proof of a conjecture of Morton Curtis that $\,\mathcal{D}\,$ fails to possess small simply connected open neighbourhoods about certain points. This property is stronger than local simple connectivity (see our comments in II §1). A proof of Curtis' Conjecture was anounced in 1964 [14]; however the detailed proof has not appeared. Only one other topological property distinguishing $\mathcal D$ and E^3 is known: some points of $\mathcal D$ cannot be enclosed in 2-spheres [11], [13]. The general state of affairs seems to be that some points of \mathcal{D} have no closed or open 3-cell neighbourhood systems, but do have systems of neighbourhoods bounded by double tori.

Our arguments use elementary methods exclusively (except for an easily circumvented reference to the Hopf property of knot groups) and may well appear old-fashioned. We are less than proud of much of the exposition, which was intended to combine the detail appropriate to a

thesis with the directness of a journal paper and somehow didn't. The reader will probably share our pain at the length of the argument (the whole paper is essentially one theorem). The reader who is unfamiliar with pathological decomposition spaces is advised to read [3], which is brief and exceptionally entertaining, and then skim Ch. II.. We will mention some notational peculiarities: we follow common practise in describing geometric constructions, even complicated ones, by the use of diagrams. "Theorem' in this paper means 'working theorem'; thus 'theorems' appear in the introductory chapter conly.

CHAPTER ONE

0. Introduction.

This first chapter gives preliminary material for the arguments in Ch III and especially Ch IV. The reader who wishes to skim the paper will find that Ch II, which contains the discussion of Curtis' Conjecture, is largely independent of this first chapter. In this paper, our approach to elementary topology is along the lines of the easier chapters of [10[, in particular, we always assume a separable metric space. In this chapter, sections 1 and 2 are elementary, §3 contains working theorems for Ch IV, and §4 is essentially a comment on Bing's Theorems 6 and 7 of [12]. Section 5 is part of the argument of Ch IV which is self-contained and has been smuggled into the preliminary material, although it could have been left until it appeared naturally in the main argument.

1. Notation.

The arguments in this paper use elementary methods exclusively, so that notation should present morproblems. We use of for the null set and the symbol Informathe end of the proof notation numbered result. The expression 'Bd A' may mean either the manifold boundary of the mainifold-with-boundary A, or the point-set boundary of the set A. A similar comment applies to the expression 'Int A'. This reflects common practise; we will comment whenever the meaning is unclear. As mentioned in the preface, our attitude to the construction of tame sets will be cavalier; we will construct many important tame sets simply by describing the set and perhaps giving a picture of it. We advise against the intuitive approach of imagining our constructions as straight-sided polyhedra whose structural detail is so fine that the polyhedra approximate the figures closely. Several of our arguments will require extensive repair if our geometric constructions are interpreted in this way. If neccessary, methods in [4] could be used to

show that each of our constructions is in fact a curvilinear polyhedron.

2. Elementary Results.

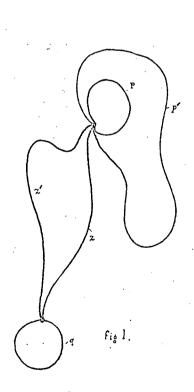
In this section we give some 'obvious' results which we have found hard to justify by simple references. This may be a matter of ignorance, especially in the case of (2.1) and (2.3). We define an annulus to be a topological sphere with two holes. The proofs are omitted.

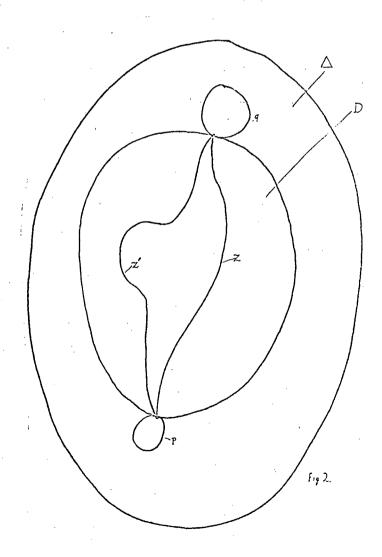
- (2.1). Let a be an arc which intersects two disjoint closed sets s_1, s_2 . Then there is a sub-arc a* of a which connects s_1 and s_2 and meets $s_1 \cup s_2$ only at the end points of a*.
- (2.2). Any two annuli A_1 , A_2 are homeomorphic. Any homeomorphism of one boundary component of A_1 onto a boundary component of A_2 may be extended to a homeomorphism of A_1 onto A_2 .
- (2.3). The union of two locally connected (1c) continua which intersect is a 1c continuum.
- (2.4). Let 0 be a bounded connected open set in the plane whose boundary is 1c. Then any two points x and y in $\overline{0}$ may be connected by an arc which lies in 0 except possibly for its end points.
- (2.5). Let A be a 2-manifold with boundary, and K a continuum in A. Then any two points of K may be connected by an arc in Int A (except possibly for end points) which lies within a distance ϵ of K.
- (2.6). Let C_1 , C_2 be disjoint simple closed curves in E^2 . Then

one of the following exclusive alternatives is true:

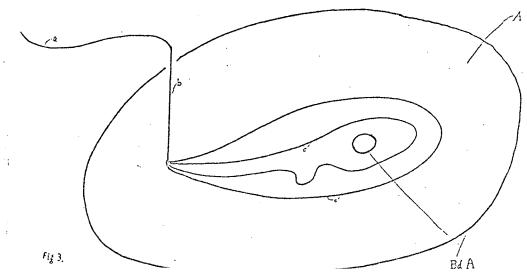
- a) $C_1 \subset Int C_2$ or equivalently $\overline{Int C_1} \subset Int C_2$.
- b) $C_2 \subset Int C_1$ or equivalently $\overline{Int C_1} \subset Int C_2$.
- c) Each of C_1 , C_2 lies in the others exterior, or equivalently $\frac{1}{\text{Int } C_1} \cap \frac{1}{\text{Int } C_2} = \emptyset$
- (2.7). Let A be an annulus, and C a simple closed curve in Int A which bounds no disk in A . Then C separates A into components B_1 , B_2 such that $B_1 \cup C$ and $B_2 \cup C$ are annuli.

3. Sliding Curves on Spheres.





Suppose p (/ g (/ z lies on a disk $\Delta \subset E^3$. We ask what properties the homeomorphism H should have in order to reflect the intuitive idea of sliding z to z' on Δ while keeping p ψ q fixed. One way to do this would be to construct a new disk $D \subset \Delta$ (see fig. 2) so that Dcontains $z \cup z'$ and misses p and q except where they hit $z \cup z'$. Then we could require that H carry z onto z $^{\prime}$, H[D] = D , and that H be the identity on Δ - D and on Bd D, (thus H will fix p \cup q). It seems a good idea to specify a number of standard moves, prove that they can always be made and stick to these in the sequel. When, as commonly happens, an arc or loop moves only a short distance and has explicit initial and final locations, then our idea of 'standard moves' is probably too formal. However our standard moves are intended for the case that the initial position of the set is unknown. In this case the existence of the required move is less obvious, especially when, as in §5, Th 6, a base point must be held fixed during the move. If S is a sphere with n holes in E^3 , then a collar of S is the image of an embedding h of $S \times [-1,1]$ into E^3 so that h(x,0) = x. Evidently a collar of S may not exist (S could be wild). A set upon which a collar has been constructed is called a collared set. Note that a collar of S is not a neighbourhood of S.



(3.2). A- B- and B'-moves.

We give three standard moves in Theorems 1 and 2.

Theorem 1. Let D be a disk in E^3 , J a collar of D, and a', a two arcs which have common end points and lie in Int D except for these end points, which lie in Bd D. Then there is a homeomorphism A(a,a',D,J) of E^3 onto itself which carries a onto a', D onto itself, and which fixes $E^3 - J$.

We call A(a,a',D,J) 'the A-move' and say that A(a,a',D,J) moves a to a'. (Of course the fact that a moves to a' is only one of a number of things that have to be kept in mind. We write the move as a function of D and J to emphasize that the trick of using the move depends on the right definition of D and J).

Theorem 2. Let A be an annulus in E^3 . Let c, c', be simple closed curves which lie in the interior of A and bound no disks in A. Let Q be a collar of A. Then there is a homeomorphism B(c,c',A,Q), also called a B-move, of E^3 onto itself which carries c onto c', A onto itself, and which fixes Bd Q and $E^3 - Q$. If, in addition, c and c' have a common base point y, then there is a homeomorphism B'(c,c',A,Q) and the following additional property: if h is the embedding associated with Q, so that $Q = h[A \times [-1,1]]$, then the B'-move fixes y and in fact all of $h[y \times [-1,1]]$.

The B'move is a move 'keeping the base point fixed'. One could probably fix the base point by providing that c \cup c' could hit Bd A \subseteq Bd Q

so that $y \in Bd \ Q$ (the B-move does not permit this), however the B'-move as given above fits the intended applications better and is easier to prove. We will give an example which shows why we want the B'-move to fix $h[y \times [-1,1]]$. Fig. 3 shows c, c', A, and an arc $a \cup b$ such that a misses A and b is a straight arc perpendicular to A. We want to move c to c' while leaving $a \cup b$ fixed. We do this with a B'-move B'(c, c', y, A, Q) in which Q is defined so that all points of Q lie near A (i.e. for $x \in A$, $h[x \times [-1,1]]$ is short) and so that each arc $h[x \times [-1,1]]$ with $x \in A$ is perpendicular to A. For a sufficiently 'thin' Q, the B'-move will fix a because $a \subset E^3 - Q$, and b will be fixed because b lies in $h[y \times [-1,1]]$ wherever it hits Q. Evidently the utility of the B'-move is limited. However subsequent use of the B'-move will be very much along the lines of this example.

4. The Phragmen-Brouwer Properties. The Zoretti Theorem.

The Phragmen-Brouwer Properties are usually given for the n-sphere, but hold also on a disk. We quote from Wilder, [I, II 4.1]. Let S be a locally connected metric space. Then the following properties of S are equivalent.

- (4.11). If A, B are disjoint, closed subsets of S, and $x,y \in S$ such that neither A nor B separates x and y in S, then A \cup B does not separate x and y in S. (By 'X separates x and y in S' is meant 'x and y are in different components of S X').
- (4.12). If $S = A \cup B$, where A, B are closed and connected, then

A \(\text{B} \) is connected.

(4.13). If A, B are disjoint closed subsets of S and a \in A, b \in B, then there exists a closed connected subset C of S - (A \cup B) which separates a and b.

Theorem II 4.12 of [1] states that these properties are equivalent in a locally connected metric space. From VII, 9.3 of [1] (note also 9.2), a disk D will have properties (4.11), (4.12), (4.13), if its first Betti number is zero; thus (4.11) ... (4.13) hold on D.

(4.2). We get the following important working theorems from (4.11). These theorems resemble Theorems 6 and 7 of [12].

Theorem 3. Let D be a 2-cell in E^3 and F_1 , F_2 closed disjoint subsets of E^3 . Let pxq, pyq be arcs in D which share the end points p and q, and such that arc pxq misses F_1 , arc pyq misses F_2 . Then there exists an arc pzq with end points p, q such that arc pzq c D - F_1 - F_2 .

Theorem 4. Let D, F_1 , F_2 , p, q, arcs pxq, pyq be defined as in Th 3 except that arc px v arc yq misses F_1 , arc py v arc xq misses F_2 . Then there exists an arc pzq D with end points p, q, such that arc misses either F_1 , or F_2 .

Proofs of Th 3 and Th 4. Since D is simply connected, pxq and pyq are homotopic in D by a homotopy which fixes p and q.

Using this fact, the proofs of Th 6 and Th 7 of [12] may be used word for word to prove Th 3 and Th 4 respectively, reading D for M in [12] D.

- (4.4). The Plane Separation Theorem and the Zoretti Theorem.
 We quote these results, slightly simplified, from [10, VI §3].
- (4.41). The Plane Separation Theorem. Let A, B be compact sets in E^2 which intersect in at most one point. Let $a \in A B$, $b \in B A$, and let $\varepsilon > 0$. Then there is a simple closed curve J which separates a and b in E^2 , lies within an ε -neighbourhood of A, and misses $A \cup B$ except possibly at the point $A \cap B$.
- (4.42). The Zoretti Theorem. If K is a component of a compact set M in the plane, then there is a simple closed curve J whose interior contains K, which misses M, and which lies in an ϵ -neighbourhood of K.

5. Annulus Dodging Theorems.

Suppose A is an annulus and F is a closed set in A . When can we say that a simple closed curve which looks like c in fig. 4 exists so as to miss F? The answer is about what would be expected. We say that \underline{F} bridges \underline{A} iff the two boundary components of \underline{A} are in the same component of \underline{Bd} \underline{A} \underline{U} F, or equivalently, iff some component of F meets both boundary components of \underline{A} .

We will prove the equivalence. Let the boundary components of A be ℓ and m . ' + ' is obvious. If no component of F meets both ℓ and m, then no component of F meets both $\ell \cap F$ and m $\cap F$, and by I(9.3) of [10] (taking A, B, K to be $\ell \cap F$, m $\cap F$, F), there is a separation of F into compacta F_ℓ , F_m such that F_ℓ meets only ℓ , F_m meets only m in Bd A. Evidently this denies the existence of a connected subset of

Fulum which meets both ℓ and m.

(5.1). If F fails to bridge A, then there is a simple closed curve c in Int A such that c bounds no disk in A and c misses F.

Proof: We can assume that A is the set $1 \le x^2 + y^2 \le 2$ in E^2 . Let D be the set $x^2 + y^2 \le 1$. Let ℓ , m be the boundary components $x^2 + y^2 = 1$, $x^2 + y^2 = 2$ respectively.

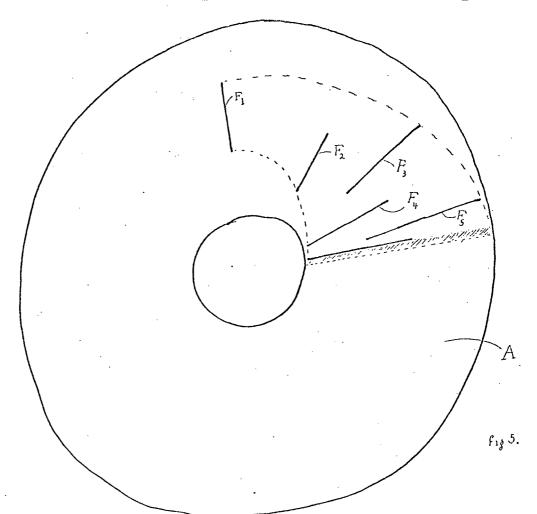
Consider the component K of $\ell \cup m \cup F$ which contains (the connected set) ℓ . The set $\ell \cup m \cup F$ is clearly compact, and by the Zoretti theorem (4.4.2) there is a simple closed curve c which lies in $E^2 - F - \ell - m$, contains K in its interior and lies in an ε -neighbourhood of K. We will show that c has the properties required by (5.1). To see that c Int A: K contains ℓ and misses m, since otherwise F bridges A. Thus $K \subset (A \cup D) - m = Int(A \cup D)$. Since K is compact, K has an ε -neighbourhood in $Int(A \cup D)$, and we can assume that c lies in this neighbourhood. Thus $c \in Int(A \cup D)$. But c encloses $k \supset \ell$ and hence D (by (2.6)); therefore $c \in Int(A \cup D) - D = Int A$. We know that c bounds no disk in A because, from the Schoenflies theorem, c bounds just one disk in E^2 , This disk is $\overline{Int c}$ which is not a subset of A since it contains D. Since c misses F (by construction), lies in Int A, and bounds no disk in A, the proof of (5.1) is complete \square .

Remark: the converse of 5.1 is true and easily proved.

We will look at some generalizations, the choice being influenced by later applications.

Theorem 5. Let F_1 , F_2 be disjoint closed sets in the annulus A . If each of F_1 , F_2 fails to bridge A , then there is a simple closed curve c in Int A - F_1 - F_2 such that c bounds no disk in A .

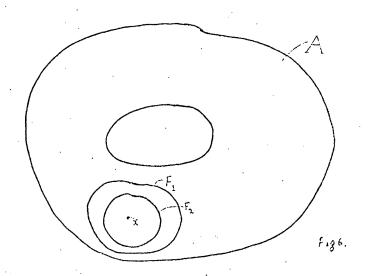
Proof: This result is trivial once we show that if neither of F_1 , F_2 bridges A , then $F_1 \cup F_2$ fails to bridge A . Once this is done, the proof of Th. 5 is completed by applying (5.1) taking F to be $F_1 \cup F_2$. To see that $F_1 \cup F_2$ fails to bridge A : since F_1 does not bridge A , no component of $\ell \cup m \cup F_1$ intersects both ℓ and m , (for otherwise some component of $\ell \cup m \cup F_1$ would contain ℓ and m). By I (9.3) of [10], taking A , B , K in that theorem to be ℓ , m , $\ell \cup m \cup F_1$, there is a separation of $\ell \cup m \cup F_1$ into disjoint compact sets U_1 ,



U_2 so that $\ell \subset U_1$, $m \subset U_2$. Similarly there is a separation of $\ell \cup m \cup F_2$ into disjoint compact sets V_1 , V_2 , with $\ell \subset V_1$, $m \subset V_2$. It is easily checked that $U_1 \cup V_1$ misses $U_2 \cup V_2$. Evidently $\ell \cup m \cup F_1 \cup F_2$ may be separated into the disjoint closed sets $U_1 \cup V_1$ and $U_2 \cup V_2$ with $\ell \subset U_1 \cup V_1$, $m \subset U_2 \cup V_2$. Therefore ℓ and m are not in the same component of $\ell \cup m \cup F_1 \cup F_2$ and $F_1 \cup F_2$ fails to bridge A \square .

We remark that ${}^tF_1 \cup {}^tF_2 {}^t$ may be replaced by a finite union of disjoint closed sets with a few trivial changes in the proof. Theorem 5 is false for a non-compact union of sets $F_1, F_2, \ldots, Fig. 5$ shows A and a collection F_1, F_2, \ldots such that A is the set $1 \le r \le 2$ in polar coordinates and for $1 = 1, 2, 3, \ldots, F_i$ is a subset of the ray $\theta = 1/i$. Although each F_i does not bridge A (nor does the union $\bigcup_{i=1}^{\infty} F_i$), the curve c in Th. 5 cannot be constructed.

We next look at the case where the curve c is constructed as in Th. 5 but with the further property that c contains a given base point ${\bf x}$. In this case c cannot in general miss either of ${\bf F_1}$, ${\bf F_2}$, as Fig. 6 shows.



We will give a characterization of those placements of x, F_1 , F_2 , so that c can be made to miss one of F_1 , F_2 . We say that a simple closed curve c with base point x has Property $^{\text{L}}P$ (read 'property not - P') with respect to closed sets F_1 , F_2 iff one of the following is true:

 $^{\sim}$ P(a): c misses on of F_1 , F_2 . $^{\sim}$ P(b): There exists a point $y \in c - x$ and a decomposition of c into $arcs c_1, c_2, with c_1 \cup c_2 = c$ and $c_1 \cap c_2 = \{x,y\}$, such that F_1 misses c_1 , F_2 misses c_2 (see fig. 7).

This is an ugly and awkward definition. An equivalent and prettier statement is 'c has Property $^{\prime}$ P iff any point in c - x may be joined to x by an arc which misses one of F_1 , F_2 '; however we will not prove this, and we will use the earlier statement exclusively. The odd name of this

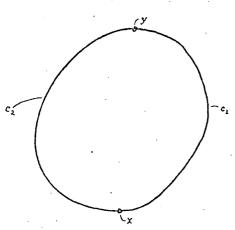


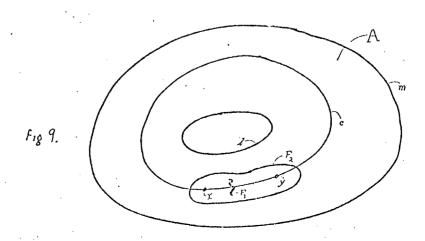
Fig 7.

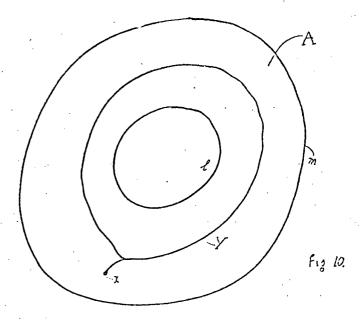
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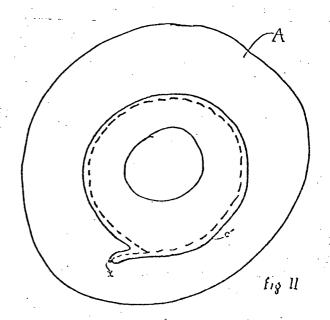
property is intended to recall Bing's Property P in [12]. This property is defined on double ended lassos (see fig. 8). Later we will define Property $^{\circ}$ P on double ended lassos and it will turn out that the loops of such lassos, with the obvious base points, have Property $^{\circ}$ P in the present sense. The next theorem says that if c with base point x has Property $^{\circ}$ P, then there is a loop c' which behaves like c and misses one of F_1 , F_2 .

Theorem 6. Let A, F_1 , F_2 be defined as in Th. 5, including the condition that neither F_1 nor F_2 bridges A. Let $x \in Int A$. Let c be a simple closed curve which lies in Int A and bounds no disk in A and contains x. If c has Property $^{\circ}$ P with respect to x, F_1 , F_2 , then there exists a simple closed curve c' which lies in Int A, bounds no disk in A, has base point x, and misses one of F_1 , F_2 .

This result cannot be improved so as to allow us to specify which of F_1 , F_2 is to be missed by c'. Fig. 9 shows a case where c' in Th. 6 cannot be made to miss F_2 although F_1 U F_2 fails to bridge A , and c exists with Property ${}^{\circ}$ P . (There are simpler counter examples in which only F_2 hits c . One of these may be derived by removing F_1







from fig. 9. However fig 9 shows that matters do not improve if we insist that both $\, \mathbf{F}_1 \,$ and $\, \mathbf{F}_2 \,$ hit c .)

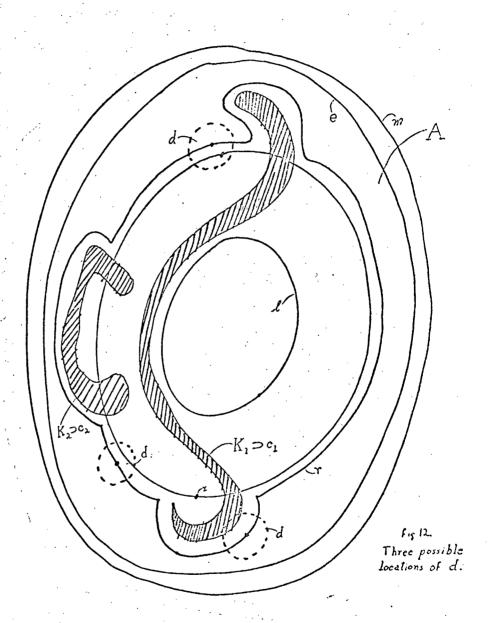
Proof of Th 6. We can assume that A is the set $1 \le x^2 + y^2 \le 2$ in E^2 . The inner and outer boundary components of A will be called ℓ and m respectively. Since neither of F_1 , F_2 , bridges A, it follows from Th 5 that there is a simple closed curve $e \in Int A$ which bounds no disk in A and misses $F_1 \cup F_2$. If $x \in e$, then the proof is completed by letting e be e ; thus we assume that e . We make the further assumption that e interpretation is easy to remove. Assuming that e interpretation e is either e or a curve which behaves like e and is constructed similarly, while the 'handle' of Y joins the loop to e . The whole of Y misses one of e its near Y and meets e as shown in fig 11.

Construction of Y . The lasso Y consists of the union of a simple closed curve $\, r \,$ and an arc $\, s \,$, and is constructed so as to have the following properties:

Y C Int A, $\begin{tabular}{ll} Y & misses & one & of & F_1, & F_2, \\ \\ the & circle & r & bounds & no & disk & in & A & , \\ \\ the & end & points & of & s & are & x & and & a & point & z & c & r & , \\ \\ & & and & s-z & misses & r & . \\ \end{tabular}$

The construction of Y is divided into two cases.

Case one: e meets c-x. We assume that c satisfies Property v P(b), since if c satisfies Property v P(a), we immediately let



c' = c . Thus we take c to be the union of arcs c_1 , c_2 which meet only at their end points, and for i = 1, 2, c_i misses F_i . If e meets, say, c_1 (it will do no harm if e meets both c_i), then use (2.5) to construct an arc s which joins x and e \cap c_1 , and lies so near c_1 that s misses F_1 (or take the obvious sub-arc of c_1). Let r = e, Y = $r \circ s$. To check that Y has the required properties; Y misses one of F_1 , F_2 because s misses one of F_1 , F_2 and r misses both; r = e bounds no disk in A by construction; and $Y \subset Int$ A because $e \circ c \subset Int$ A. Finally, from (2.1), we can assume that s meets r only at a single point z.

Perty \sim P(b) . Outline of proof: a) As usual we take c to be the union of arcs c_1 , c_2 ; let K_1 be c_1 plus those components of F_2 which hit c_1 and let K_2 be c_2 plus those components of F_1 which hit c_2 . $K_1 \cup K_2$ is a component of $c \cup F_1 \cup F_2$. b) A 'Zoretti curve' r is constructed so that r misses $c \cup F_1 \cup F_2$, lies in Int A, encloses $K_1 \cup K_2$, and bounds no disk in A . c) Some care needs to be taken to attach the tail s to r so that s misses one of the F_1 . Construct a disk $d \subset Int A$ with centre on r, (see fig 12) so that d is big enough to hit $K_1 \cup K_2$ but small enough to miss one F_1 . This is managed by a careful choice of the ε associated with the Zoretti curve. d) There is an arc s near $K_1 \cup K_2 \cup d$ which has the required properties.

Details of proof.

a) Let $K_1 = c_1$ plus those components of F_2 which hit c_1 . Let $K_2 = c_2$ plus those components of F_1 which hit c_2 . We will show that $K_1 \cup K_2$ is a component of $c \cup F_1 \cup F_2$. Let K be the component of $c \cup F_1 \cup F_2$ which contains the connected set $K_1 \cup K_2$, and suppose

that some point p exists in $K - (K_1 \cup K_2)$. Then p lies in one of F_1 , F_2 , say F_1 . Since $p \notin K_1 \cup K_2$, no component of F_1 meets both p and $c \cap F_1$. By I(9.3) of [10], there is a separation of F_1 into compacta U_1 , U_2 containing $c \cap F_1$ and p respectively. Evidently U_2 misses not only U_1 but F_2 and the whole of c; thus $U_1 \cup F_2 \cup c$ is a compactum disjoint from U_2 , and there is therefore a separation of $c \cup F_1 \cup F_2 = U_1 \cup U_2 \cup c \cup F_2$ into compacta containing c and p severally. This denies the assumption that p lies with c in a connected subset of $c \cup F_1 \cup F_2$.

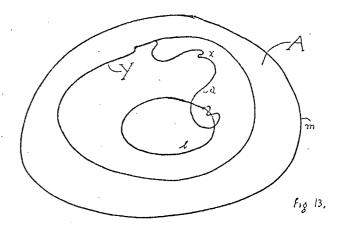
- b) Since $x \in Int e$ and e misses $c \cup F_1 \cup F_2$, all of $K_1 \cup K_2$ lies in Int e by the usual argument. Since $K_1 \cup K_2$ is a component of $c \cup F_1 \cup F_2$ we can construct a Zoretti curve r which misses $c \cup F_1 \cup F_2$, encloses $K_1 \cup K_2$, and lies withing a distance e of $K_1 \cup K_2$. The following argument shows that r bounds no disk in A: since e bounds no disk in e in e, which is a disk, must meet points of e in e in e. Since e encloses e in e in
- c) We construct a (closed) disk of radius 2ϵ with centre anywhere on r . Clearly d will hit $K_1 \cup K_2$. We show that d C Int A by showing

that dc Int m and d misses $\overline{\operatorname{Int} \ell}$. The distance ϵ could have been chosen so that 4ϵ (i.e. the diameter of d) is less than the distance separating c and Int ℓ , and the distance separating $K_1 \cup K_2$ and m; and we assume that this was done (the last distance is positive because $K_1 \cup K_2$ is compact and lies in Int m). Since d hits $K_1 \cup K_2$, $d \subset \operatorname{Int} m$ by the choice of ϵ . If d hits $\overline{\operatorname{Int} \ell}$, then d must also hit c, since points of $d \cap r$ lie in the exterior of c which encloses $\overline{\operatorname{Int} \ell}$ as we saw. By the choice of ϵ , d cannot meet both c and $\overline{\operatorname{Int} \ell}$; thus d misses $\overline{\operatorname{Int} \ell}$.

We also assume that $^{\epsilon}$ was chosen so that for $i=1, 2, 4\epsilon$ is less than the distance separating K_i from F_i . The disk d must hit one of K_1 , K_2 , say K_1 (it does no harm if d hits both K_i). Since d hits K_1 , d misses F_1 since otherwise F_1 would be closer to K_1 than the diameter of d.

d) The continuum $K_1 \cup d$ meets x and r and misses F_1 . Using (2.5), let s be an arc in Int A which joins x and r and lies so near $K_1 \cup d$ that s misses F_1 . (note that although K_1 may not miss Bd A, (2.5) provides that s misses Bd A). To see that $Y = r \cup s$ has the required properties: $r \cup s \subset Int$ A by construction; Y misses one of F_1 , F_2 because s misses one F_1 and r misses both. The circle r bounds no disk in A as we saw in b); and finally we can assume that s has end points x and $z \in r$ with $r \cap s - z = \emptyset$ by (2.1).

This completes the construction of Y assuming that $x \in Int e$.



Construction of c'. We maintain the assumption that $\times \mathcal{E}$ Int e during this construction. We will first construct a continuum a C A which joins $\overline{\operatorname{Int} \ell}$ to x so that a meets Y only at x. As suggested by fig 13, the plane separation theorem can then be used to separate Y - x and $\overline{\operatorname{Int}}\ \ell\ \sigma$ a - x . Construction of a: since r encloses c, x \in Int r. Let Q be the open set Int r-s. Q is connected because s-r does not disconnect Int r ([10, VI(3.4)]). Evidently Bd Q = Y, and since each of r, s is a 1c continuum, so is Y. Using (2.4), connect x to a point in $\overline{\operatorname{Int}\,\ell}$ by an arc a which lies in Q except for x . Since Y misses Int ℓ , Y and Int ℓ U a are continua in Int m which meet only at the point x . We now use (4.41) to separate Y - x and $(\overline{\operatorname{Int}\,\ell}\,\,\cup\,\,a)$ - x by a circle c' which lies so close to Y that it misses one of F_1 , F_2 . Evidently c' must pass through \mathbf{x} (since otherwise $(\overline{\operatorname{Int}\,\ell}\,\cup\,\mathbf{a})$ - x and Y - x are subsets of a connected set in E^2 - c'). We know that c' \subset Int A because c' misses Int ℓ by construction and c' lies so near Y C Int m that c' C Int m. It remains to show that c' bounds no disk in A . To see this: we know that rencloses ℓ . This means that c' cannot enclose r, since this would imply that Int c'cru ℓ (from (2.6)), whereas we know that c' separates r and

Thus Ext $c'\supset r$ and Int $c'\supset \ell$. The fact that Int $c'\supset \ell$ implies that c' bounds no disk in A by the usual argument.

The construction of c' is now complete except that the restriction $x \in I$ intermined in the case that $x \in I$ is easy if $x \in I$, we only look at the case that $x \in I$ is easy to construct a homeomorphism f of f onto itself which exchanges f and f i.e. f i.e. f if f is f intermined in the fact that f is easy to construct a homeomorphism f of f onto itself which exchanges f and f i.e. f i.e. f if f is f intermined in the fact that f in the fact that f intermined in the fact that f intermined in the fact that f is now an analysis of the fact that f in the fact that f is now an analysis of the fact that f in the fact that f is now an analysis of the fact that f is now an analysis of the fact that f is now an analysis of the fact that f is now an analysis of the fact that f is now an analysis of the fact that f is now an analysis of the fact that f is now an analysis of the fact that f is now an analysis of the fact that f is now an analysis of the fact that f is now an

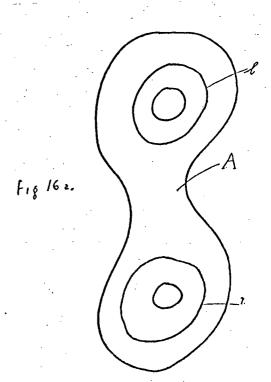
Theorem 7. Let A, F_1 , F_2 be defined as in Th 5 and Th 6 except that F_1 bridges A while F_2 does not. Let c be a simple closed curve in Int A which bounds no disk in A and contains the base point x. Then if c has Property $^{\circ}$ P with respect to x, F_1 F_2 , there exists a simple closed curve c' which meets x, lies in Int A, bounds no disk in A, and misses F_2 .

Th 7 is proved in the same way as Th 6. At first glance one might think that one of Th 6, Th 7 is stronger than the other; but in fact this is not true. If F_1 , F_2 fail to bridge A, one might wish to add pieces to, say, F_1 so that the enlarged F_1 would bridge A; this would obtain the conclusion of Th 7 which is stronger than that of Th 6 (since it predicts which F_1 is hit by c'). However it may not be possible to do this $(F_1 \cup F_2)$ might be a number of circles concentric with $F_1 \cup F_2$ might be a number of circles concentric with $F_2 \cup F_3$

Proof of Th 7. Use (5.1) with F taken to be F_2 to construct

a simple closed curve e which lies in Int A, misses F_2 , and bounds no disk in A . If e meets x, then e is the required c'. If $x \notin e$, assume that $x \in Int e$ as before. Since e bounds no disk in A, e separates ℓ and m by (2.7); in particular e meets some component k of $\boldsymbol{F_1}$ which hits both ℓ and \boldsymbol{m} (there must be at least one since ${\bf F}_1$ bridges A). For similar reasons, c meets the same component ${\bf k}$. Let $y' \in k \cap e$. Because c has Property \circ P, there is an arc $b \in c$ such that b misses one of F_1 , F_2 and connects a point of k to x. Since $k \subseteq F_1$, evidently bulk misses F_2 . Since bulk is a continuum which connects $y' \in e$ to x and misses F_2 (as does e), we can construct the lasso Y as in the proof of Th 6, reading b \cup k for c \cup d and e for r . In the proof of Th 6, the curve r misses $\mathbf{F}_1 \cup \mathbf{F}_2$, whereas here e misses just F_2 ; however following the procedure of the proof of Th 6 will yield a lasso Y which misses F_2 . The lasso Y is used to construct c' precisely as in the proof of Th 6, keeping in mind the fact that Y misses F_2 , so that the resulting c'also misses F_2 . The assumption that $x \in Int c$ is removed just as in the proof of Th 6 \square .

1. An upper sémicontinuous decomposition G of E into compact sets (or simply a decomposition of E³) is a collection of disjoint compact sets Λ of E^3 such-that the union of the elements of the decomposition is E^3 , and each element $\Lambda \in G$ possesses a system of open neighbourhoods which are unions of elements of G . The decomposition space G associated with G is a topological space in which each point is as element $\Lambda \in G$, and the open sets are just those subsets of Gthe union of whose elements is open when considered as subset of \mathbf{E}^3 . Thus each point Λ of G has a system of neighbourhoods each of which is open 'both in G and in E^3 '. One can use this intuitive idea to get a certain geometric grasp of the topology of G simply by remembering that 'some points are sets' and keeping an eye on the neighbourhoods; for example one often does geometry on a torus or Klein bottle by looking at the equivalent decomposition space of a rectangle 'with certain sides identified. If an element $\Lambda \in G$ contains more than one point of E^3 , then Λ is called a big element of G . If Λ is a singleton, then Λ is a small element of G . In G, the corresponding points are called big and small points. The decompositions G in which we will be interested are all pointlike, which is to say that the complement of each $\Lambda \in G$ is topologically equivalent to that of a point; in particular, each Λ is connected. We definitely assume some acquaintance with these ideas and do not regard the present text as an adequate introduction. The classical approach to decompositions and decomposition spaces may be found in Ch VII of [10]. Our approach will be more along the lines of [3,56]. We will use two main classical results: i) an upper semicontinuous decomposition space (i.e. the decomposition space associated with an upper semicontinuous decomposition) of E^3 is a separable metric space. ii) there is an



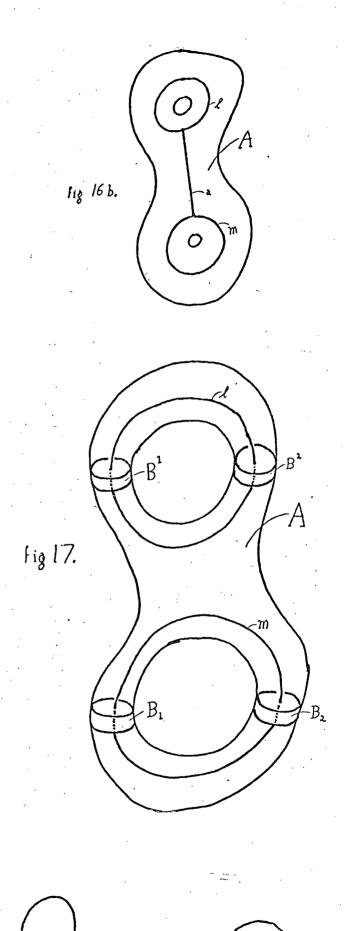
obvious way of expressing G as a quotient space. In this case the quotient topology turns out to be the decomposition space topology, and the canonical mapping ϕ of the quotient space carries each Λ \mathcal{E} G onto the corresponding point $\phi[\Lambda]$ in G. We will often write A^* for $\phi[A]$ if A is a subset of E^3 . In the sequel, 'decomposition space' will mean 'pointlike upper semi continuous decomposition space of E^3 .'

An important question is: if G is a decomposition space, is G homeomorphic to E^3 ? That G is homeomorphic to E^3 is <u>Wardwell's conjecture</u> (in [8]) and is known to be false. R. H. Bing showed this in 1957 with a celebrated example ([12]) which reinforced everyone's worst prejudices against the analytic topology of E^3 . In Bing's example, the dogbone space of our title, most of the elements of the decomposition are small. Each big element is a tame arc (so that the example refutes a very strong form of Wardwell's conjecture), and the big points in the decomposition space form a totally disconnected set.

Detailed construction of the dogbone decomposition.

We will describe an infinite sequence of compact sets whose elements intersect to form the set of big elements of the dogbone decomposition G. Our construction differs slightly from Bing's, but we assume an acquaintance with the original construction in [12] and will not prove, for example, that the various embeddings to be described can be assumed to be polyhedral.

Dogbone space takes its name from the distinctive shape of the double handlecube A depicted in fig 16a. We imagine A imbedded in E 3 . A path ℓ \subset Int A, which makes one circuit of the circle marked ℓ



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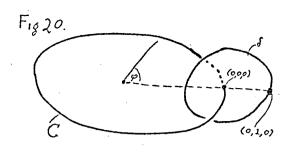
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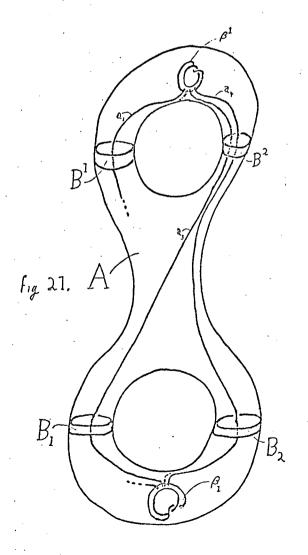
Facing page 26 (π) double twist holes of A double tvist F18 19

in fig 16a is called the upper eye of A . A path m ⊂ Int A which makes one circuit of the curve marked m in the figure is called the lower eye of A (we imagine the dogbone placed vertically in ${\ensuremath{\text{E}}}^3$ so that it makes sense to talk about 'upper' and 'lower' here). One could imagine A to be a closed r-neighbourhood of a planar double ended lasso consisting of the eyes ℓ and m laid out as nice circles plus a straight connecting arc a (with r of course, taken sufficiently small, say less than onethird of the common diameter of the nice circles ℓ and m). We call ℓ υ m υ a the centre of A . The centre of a dogbone will not be important in this chapter (but will be needed in Chapters III, IV). The idea of A as an r-neighbourhood of its centre k is introduced mainly to pin down the embedding of $\,k\,$ in $\,A$; we usually draw $\,k\,$ and $\,A\,$ as in fig 16b. Fig 17 shows four short solid cylinders B^1 , B_2 , B_1 , B_2 , which are subsets of A and cut into the eyes of A as the figure suggests. The removal of one of B^1 , B^2 and one of B_1 , B_2 from A leaves a set whose closure is a cube. A dogbone can be imagined in the topologically equivalent form of a thick double ended lasso as shown in fig 18. In a sense, we are pictorially confusing the dogbone with its centre. Let A_1 , A_2 , A_3 , A_4 be four dogbones embedded as shown in fig 19 by embeddings h_i : $A \rightarrow A$, j = 1, 2, 3, 4 so that the / $A_{i} = h_{i}[A]$ are mutually disjoint and lie in Int A . In fig 19, two double twisted bands β^1 and β_1 are placed so that $\beta^1(\beta_1)$ lies in the interior of the upper (lower) component of $A - B_1 - B_2 - B^1 - B^2$. In the obvious way, the centre of A_j is called k_j , j = 1, 2, 3, 4, with upper loop ℓ_{j} and lower loop m_{j} . The ℓ_{j} are placed so as to lie as parallels on $\,\beta_{1}^{}$. The connecting arcs $\,a_{_{\dot{1}}}^{}\,$ are laid out in a peculiar

Toroidal Coordinates (p, r, 0).

(r,0) defines a point on the disk S. As p increases in Osps27, S sweeps a toroid which is a figure of revolution about the planar circle C.





way which is characteristic of the dogbone construction. Using toroidal coordinates (which we recall in fig 20), we could define β^1 and β_1 to be appropriate translations of the set $r \leq 1$, $\theta = \phi$ and thus construct a band with an even double twist. However the bands in the drawing are translations of the set

$$r \le 1$$
, $\theta = 0$: $\pi/3 \le \phi \le 2\pi$
 $r \le 1$, $\theta = 6\phi$: $0 \le \phi \le \pi/3$

This gives a 'flatter' band and a better picture. Another concession to art appreciation is the placing of β^1 and β_1 so that their 'flat' parts lie on the plane of k. This necessitates a right angled bend in the a_j near β^1 and again near β_1 . The additional conditions are imposed on a_j that a_j misses $\beta^1 \vee \beta_1$ except at $a_j \wedge \ell_j$ and $a_j \wedge m_j$, and that the part of a_j lying within a distance ϵ of $\beta^1(\beta_1)$ consists of a single straight arc perpendicular to $\beta^1(\beta_1)$. Note that the order of ℓ_j 's on β^1 is $\ell_1, \ell_2, \ell_3, \ell_4$ while the order of m_j 's on β_1 , due to the unusual embedding of the A_j , is m_1, m_3, m_2, m_4 . The B_i and B^i locate the A_j in the following way (see fig. 21):

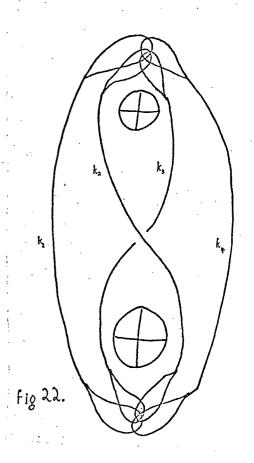
The closure of the component of $A-B^1-B^2-B_1-B_2$ which contains $\beta^1(\beta_1)$ is called $K_1(K_2)$. Finally we let $A_1\cup A_2\cup A_3\cup A_4=\mathcal{A}_1$.

Now since each dogbone A_j is homeomorphic to A, we can embed four dogbones A_{j1} , A_{j2} , A_{j3} , A_{j4} in each A_i just as the A_j are embedded in A. We could write $A_{jk} = h_j h_k [A]$. The union of the 16 A_{jk} , j, k, chosen from 1, 2, 3, 4, is called \mathcal{A}_2 . The construction proceeds as in [12] with the definition of 64 $A_{jk\ell} = h_j h_k h_\ell [A]$ where $h_j h_k h_\ell$ embeds A in A_{jk} just as A_ℓ is embedded in A. The union of the 64 $A_{jk\ell}$ is called \mathcal{A}_3 . The construction proceeds in this way, defining at each m-th stage $A_j = h_j h_k h_\ell [A]$ where the intersection $A \cap \mathcal{A}_1 \cap \mathcal{A}_2 \cap \mathcal{A}_3 \cap \ldots = A_0$. The components of A_0 are compact and are defined to be the big elements of A_0 while the remaining points of A_0 are the small elements. The dogbone space A_0 is the associated decomposition space of A_0 .

Remark 1. In $k_1 \cup \ldots \cup k_4$, each upper (lower) eye fails to shrink to a point in the complement of any other upper (lower) eye. This is easily checked using, say, Ch XV of [6].

Remark 2. We are sure that the construction of \mathcal{D} here is the same as that given by Bing in [12]. In the Appendix we show a deformation of the upper part of $k_1 \cup \ldots \cup k_4$ to look like the upper part of Bing's construction. We think that the reader will see the plausibility, but we give no strict proof that our embedding of $A_1 \cup \ldots \cup A_4$ is the same as the corresponding embedding in [12], and our attitude in this paper will be that Dogbone space has been redefined.

Remark 3. We know little about the h_j except that they embed A in certain ways. We cannot, for example be sure of the location of the 64 $h_j h_k h_\ell [k]_k$. However the various subsets of $A_{jk\ell}$ are images of sub-



sets of $A_{jk\ell}$ are images of subsets of A and continue to be related to each other in all the ways which are preserved by a homeomorphism of A; and we will usually apply results obtained for A to any $A_{jk...r}$ without further justification. Note that k_j has a property which is not preserved by homeomorphism: a_j is perpendicular to β^1 or β_1 wherever it lies near these sets. This property is lost after the first stage of the dogbone construction. This does not prevent the construction of $\mathcal D$, but further comment will be required when we use the property in Chapters III and IV.

Remark 4. Partly out of adherence to the traditional representation in [3] and partly because the use of β^1 and β_1 will not become apparent until Ch III (apart from the fact that they cause the eyes to link together) we will often use the picture in fig 22 to describe the embedding of $k_1 \cup \ldots \cup k_4$ in A. We will use pictures like fig 22 in which the crossovers of the links are ignored, whenever the exact manner of linking is unimportant. In this chapter, the only thing which needs to be kept in mind concerning the linking of the ℓ_j and m_j is that no $\ell_j(m_j)$ will shrink to a point in the complement of any other $\ell_j(m_j)$. Another pictorial abbreviation shown in fig 22 is the omission of much of the boundary of A, even though the figure purports to describe the embedding of the four centres in A. As in fig 22, we will often show only the holes of A which will be represented by the symbol Φ .

Intuitively it often helps to see a decomposition space as ${\tt E}^3$ with certain sets identified. One typically finds the small elements distributed so that it is easy to define a neighbourhood system for the

big elements. Thus a lot can be learned about the topology of the decomposition space by looking at elements of the associated decomposition. However if we try to approach $\mathcal D$ in this way, we find that the components of A_0 , which constitute the big elements of G, are hard to see. To find a big element, note that each big element of G is the limit of a sequence of dogbones A, A_1 , A_{1k} , ... Evidently each big element may be specified by an infinite sequence j, k, ... of integers chosen from 1, 2, 3, 4; and the A, A_i , A_{ik} , ... constitute a neighbourhood system of this big element. Because A is compact, we know that if Λ is a big element of G, then if Λ lies in an open set V, some member $A_{ik...r}$ of the neighbourhood system lies in V (see I, 7.2 of [10]). It is known that each big element of G is a tame arc (see [12, §2]). The canonical mapping ϕ is a local homeomorphism near small elements of (because A_0 is compact) but not of course in general. The fact that is monotone means that ϕ^{-1} preserves connectedness (VIII (2.2) of [10]). Simple connectivity properties are more complicated. As will appear later, any open set $V*\subset \mathcal{D}$ which lies in A* and contains a big point of $\,\mathcal{D}\,$ cannot be simply connected. We must expect a proof of this property to be delicate since it is known that $\,\mathcal{D}\,$ is locally simply ([5]). (Roughly, what happens is that any mapping of S^{\perp} connected. into small neighbourhood V* of a big point of $\mathcal D$ will shrink to a point in the second smallest dogbone which contains V* . Thus one can satisfy the definition of 'locally simply connected' by taking a smaller neighbourhood V* although V* itself will never be simply connected.)

For the rest of this section we will prove a result which relates simple connectivity in $\mathcal D$ to the same property in $\mathbb E^3$. A mapping

f of S^1 into a space X shrinks to a point in X iff f is homotopic in X to a constant mapping or represents the identity in $\pi_1(X)$ for an appropriate base point. A third equivalent statement is: consider S^1 to be the boundary of a disk Δ : then $f:S^1 \to X$ shrinks to a point iff f can be extended to a mapping \overline{f} of Δ into X.

(1.1). Let V* be an open set in $\mathcal D$. If $f:S^1\to V*$ so that rng f consists of small points, then $\phi^{-1}f$ will shrink to a point in V, where ϕ is the canonical mapping of E^3 onto $\mathcal D$.

Corollary: if V* is simply connected then so is V .

We can use this result hto examine sets V* which we suspect not to be simply connected, by looking at the associated V = E³. The result (1.1) and its corollary are not new and are particular cases of Lemma 1 of [2]. The proof the (1.1) introduces methods which will recur frequently in the sequel, and we will complicate the (pretty easy) proof slightly by introducing more generality in the method than is needed for the present argument.

Outline of proof. a) Assume that f maps the boundary of a disk Δ into V*. Since f shrinks to a point, there is a mapping $\overline{f}:\Delta\to V^*$ such that $\overline{f}_{\mid Bd\Delta}=f$. Recall that A_0^* is the union of the big points of $\mathcal D$. The set $f^{-1}[A_0^*]$ is compact. Let Q be a disk with holes such that $Q\subset \Delta$, the outer boundary of Q is $Bd\Delta$, and the (open) holes of Q contain $f^{-1}[A_0^*]$. b) The mapping \overline{f} maps Q into small points of $\mathcal D$; thus $\phi^{-1}\overline{f}=\overline{f}$ on Q. Let the (open) holes of Q be $\phi^{-1}\overline{f}_{\mid Bd}$ ure to a mapping $\phi^{-1}\overline{f}_{\mid Bd}$ ure to a point in a certain cube in $\mathbb V$.

c) Glue the γ_r , r=1, ... n, to $\phi^{-1}\overline{f}_{\mid Q}$ to form a mapping of Δ into V .

Details of proof. a) We know that $\overline{f}:\Delta \to V^*$ so that $f=\overline{f}_{\mid Bd\Delta}$. Since A_0 is compact, so is A_0^* and $\overline{f}^{-1}[A_0^*]$. Note that $\overline{f}^{-1}[A_0^*]$ misses $Bd\Delta$ because $\overline{f}[Bd\Delta]$ consists of small points, from the hypothesis. To obtain the disk with holes Q, we use the following result which will be needed several times in the sequel.

- (1.2) Let Δ be a disk in E² and S a compact set in Δ . Then there exist n disks W₁, ... W_n such that W_r Δ and
 - i) $W_r \cap W_s = \emptyset$, $r \neq s$.
 - ii) $S \subset W_1 \cup \ldots \cup W_n$.
 - iii) Each point of Bd W lies withing a positive distance ϵ of S .
 - iv) If S misses Bd Δ , then S \subset Int W $_1$ \cup Int W $_2$ \cup ... \cup Int W $_n$, and Δ Int W $_1$... Int W $_n$ is a disk with holes. If S hits Bd Δ , then S misses Bd W $_r$ Bd Δ for each $r=1,\ldots n$.

Proof of (1.2). We can assume that $S \neq \emptyset$ and that Δ has the form of an equilateral triangle. Triangulate Δ into a finite number of 2 - simplexes (i.e. closed triangular disks) whose diameter is less than $\epsilon/2$, and whose edges are parallel to the three sides of the big triangle Δ . Note that the three vertices of Δ each belong to one 2 - simplex only so that the three vertices of Δ cannot be cut points of any union of 2 - simplexes. The only properties of the 2 - simplexes which will be used are that each 2 - simplex has an edge of length less than $\epsilon/2$, and if two 2 - simplexes meet, they meet either along the whole of one edge or only at a vertex. Let \hat{S} be the finite union of those 2 - simplexes

which meet S . Evidently \hat{S} is 1c and each component of \hat{S} is a 1c continuum. For later reference, note that S cannot meet Bd S at a point interior to $\,\Delta;\,\,$ for assume that S meets $\,\hat{S}\,\,$ at a vertex v ϵ Int Δ . Then by construction of \hat{S} , the entire star of v lies in \hat{S} and $v\notin Bd\ \hat{S}$. S cannot meet $Bd\ \hat{S}$ at the interior of an edge in Int Δ by a similar argument. We alter \hat{S} to a set \hat{S} which has no cut poin nts in this way: a cut point of S cannot lie in the interior of a 2-simplex in S, nor in the interior of an edge belonging to one 2-simplex, nor in the interior of an edge belonging to two 2-simplexes. Thus the cut points of \hat{S} are a (finite) subset of the vertexes. Let the cut points be $t_1, \dots t_k$, and cover each t_s with a set b_s , s = 1,...k, which is a disk of radius $\epsilon/6$ and centre t_s if $t_s \subset Int \Delta$; and is a semi-disk of the same centre and radius if $t_{\rm s}$ lies on $Bd\Delta$ and is not a vertex of the big triangle Δ (thus b is a 'disk relative to Δ '). We do not define b for the three remaining points of Δ since these points are never cut points of \hat{S} . Note that the \hat{S} are disjoint. Define \hat{S} to be $\hat{\mathbf{S}} \ \upsilon \ \mathbf{b_1} \ \upsilon \ \ldots \ \upsilon \ \mathbf{b_k}$. It will turn out that the Bd $\mathbf{W_r}$ are some of the components of Bd $\hat{\hat{S}}$. We know the following facts about $\hat{\hat{S}}$: the components of $\hat{\hat{S}}$ are 1c continua and are consequently bounded apart. Every point of $\hat{\hat{S}}$ (in particular every point of Bd $\hat{\hat{S}}$) lies within a distance ϵ of S; the boundary of $\hat{\hat{S}}$ consists of the union of a finite number of straight arcs (which are either edges of 2 - simplexes or edges minus the interior of one or two $b_{\rm g}$) and a finite number of segments of circles (i.e. proper subsets of various $Bd b_s$). Such a subset is precisely Bd b_c - Bd Δ intersected with a connected subset of St t_c ; a suitable upper bound for the number of segments is the number of $b_{_{\mathbf{S}}}$

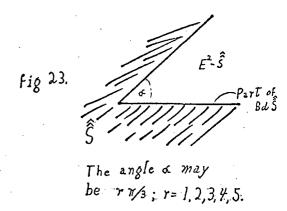


Fig 24.

Fig 24.

E2-\$

times the number of subsets of 2 - simplexes). Two straight arcs in Bd S meet as in fig 23; a straight arc meets a segment as in fig 24. Segments never meet because the b_S are disjoint. Evidently \hat{S} has no cut point on its boundary and hence no cut point at all. Let the components of \hat{S} be W_1' , ... W_m' . These will be reordered so that the first n components will lie in the disks required in (1.2). Since each W_a' , a=1, ... m_s ats a lc continuum with no cut point, by IV (9.3) and VI (2.5) of [10], the unbounded complementary domain of W_a' is bounded by a simple closed curve which will be called c_a . Evidently $c_a \in W_a'$. Reorder the W_a' and corresponding c_a so that c_1 , ... c_n are contained in the interior of no other circle c_a , while each of c_{n+1} , ... c_m , if they exist, is contained in the interior of a c_a . Let a if they exist, is contained in the interior of a a and a and a and a in the interior of a a and a and a and a in the interior of a a and a and a in the interior of a a and a and a in the interior of a a and a are interior of a a and a and a in the interior of a a and a and a and a in the interior of a a and a and a in the interior of a a in the interior of

Proof of i). If $r \neq s$ then neither of c_r , c_s is contained in the interior of the other. Then $\overline{\text{Int } c_r} \cap \overline{\text{Int } c_s} = W_r \cap W_s = \emptyset$ by I(1.6).

Proof of ii). $S \subset \hat{S} = W_1' \cup \ldots \cup W_m'$. Each $W_a' \subset \overline{Int c_a}$, $a = 1, \ldots m$, since $\operatorname{Ext} c_a$ is the unbounded complementary domain of W_a' . Hence $S \subset W_1 \cup \ldots \cup W_m$. And in fact $S \subset W_1 \cup \ldots \cup W_n$ because for $n+1 \leq a \leq m$, c_a lies in some $\operatorname{Int} c_r$, $r=1, \ldots n$, and by $\operatorname{I}(1.6)$, $\operatorname{Int} c_a = W_a \subset W_r$.

Proof of iii). Each Bd W is a c C Bd $\hat{\hat{S}}$, and we saw earlier that all of (the closed set) $\hat{\hat{S}}$ lies near S .

Proof of iv). Take ϵ less than the distance from (compact) S to $Bd\Delta$ if S misses $Bd\Delta$. The rest follows from the definition of a disk with holes, from i), and from the fact that a point of Bd \hat{S} and hence a point of Bd \hat{S} misses S wherever it lies in Int Δ \Box .

We now return to a) in the proof of (1.1). Since $f^{-1}[A_0^*]$ is compact, from (1.2) there are disks W_1 , ... W_n which lie in Δ and such that $f^{-1}[A_0^*] \subset W_1 \cup \ldots \cup W_n$ and each Bd W_r lies within a distance ϵ of $f^{-1}[A_0^*]$. Since $f^{-1}[A_0^*]$ misses Bd Δ , by (1.2)iv, the set Δ - Int W_1 - ... - Int W_n = Q is a disk with holes. Let W_r be the 'holes' of Q, $r = 1, \ldots, n$, i.e. $W_r = Int W_r$.

b) Since $\overline{f}^{-1}[A_0^*]$ lies in the holes u_r of Q, $\overline{f}[Q]$ consists only of small points. Thus ϕ^{-1} is a well defined mapping when restricted to $\overline{f}[Q]$ and $\phi^{-1}\overline{f}_{|Q}$ maps Q into $V \subset E^3$. We now find cubes in which to shrink $\phi^{-1}\overline{f}_{|Bd|u_r}$, $r=1,\ldots,n$. Since $\overline{f}[\Delta] \cap A_0^*$ is compact and the dogbones (considered as sets in \mathcal{D}) evidently form a neighbourhood system of A_0^* , there is a covering of $f[\Delta] \cap A_0^*$ by a finite number of dogbones J_1^* , ... J_q^* each of which lies in V^* . Look at the corresponding J_1 , ... J_q in E^3 . If J_s , $s=1,\ldots,q$, belongs to the mth stage of the dogbone construction, define J_{s1} , J_{s2} , J_{s3} , J_{s4} to be the four dogbones of the m+1st stage lying in J_s . Note that since $J_s \subset V$, each J_{sj} lies in a cube M_{sj} which is a subset of J_s and hence of V (if J_s were the dogbone $A \subset V$ then $J_{s1} \subset M_{s1} = A - B^2 - B_2$, J_{s2} lies in $M_{s2} = A - B^2 - B_1$ etc.) Now in a) above, we could have chosen ε so small that each $\overline{f}[Bd W_r] = \overline{f}[Bd u_r]$ lies in some J_{s1}^* (for $\mathcal D$ is a separable metric

space, and there is a minimum distance in the dogbone metric separating the compact set $\overline{f}[\Delta] \cap A_0^*$ from the complement of the union of the J_{sj}^* . Clearly $\phi^{-1}\overline{f}[Bd\ u_r]$ is defined and lies in the union of the J_{sj} . We can assume the J_{sj} are disjoint because we could have removed from the covering J_1,\ldots,J_q any J_s which was contained in any other member of the covering. Since $\phi^{-1}\overline{f}[Bd\ u_r]$ is connected and the closed sets J_{sj} are separated, $\phi^{-1}\overline{f}[Bd\ u_r]$ lies entirely in some one J_{sj} and $\phi^{-1}\overline{f}_{|Bd\ u_r}$ shrinks to a point in $M_{sj}\subset V$. Thus there is an extension γ_r of $\phi^{-1}\overline{f}_{|Bd\ u_r}$ to all of $\overline{u_r}$, i.e. $\gamma_r:\overline{u_r}+M_{sj}\subset V$ and $\gamma_r|Bd\ u_r=\phi^{-1}\overline{f}_{|Bd\ u_r}$.

- c) In view of the set-theoretic definition of function we can express the idea of 'mappings glued together' by unions of mappings. Consider the union $\phi^{-1}f_{|Q} \vee \gamma_1 \vee \ldots \vee \gamma_n$. This is a well-defined mapping of $Q \cup \text{dom } \gamma_1 \vee \ldots \vee \text{dom } \gamma_n = Q \cup u_1 \vee \ldots \vee u_n$ into V because each mapping in the union has its image in V and because where the domains intersect the intersection is closed and the mappings agree on the intersection; in fact every point of domain intersection occurs on a $Bd u_r$ where we know that γ_r agrees with $\phi^{-1}\overline{f}_{|Bd} u_r$ by construction of γ_r . Finally we note that the new mapping $\phi^{-1}\overline{f}_{|Q} \vee \gamma_1 \vee \ldots \vee \gamma_n$ agrees with $\phi^{-1}\overline{f}$ on $Bd\Delta \subset Q$ and is thus a homotopy which shrinks $\phi^{-1}f$ to a constant mapping into V. This completes the proof of (1.1) \square . We will record the argument in this paragraph as a separate result.
 - (1.3). Let Δ , W_1 , ... W_n be defined as in (1.2) including the fact that Δ Int W_1 ... Int W_n is not necessarily a disk with holes.

Let $g:\Delta$ - Int W_1 - ... - Int W_n \to E^3 . Then each $g_{\mid Bd \mid W_r}$ is defined; and if $g_{\mid Bd \mid W_r}$ shrinks to a constant mapping in a space P_r , r = 1, ... n, then there is a mapping of Δ into rng $g \cup P_1 \cup \ldots \cup P_n$. In particular, $g_{\mid Bd\Delta}$ will shrink to a point in rng $g \cup P_1 \cup \ldots \cup P_n$.

Proof: The argument of c) in the proof of (1.1) is sused and is valid even if Δ - Int W_1 - ... - Int W_n is not a disk with holes. It is easy to see that $Bd\Delta \subset \Delta$ - Int W_1 - ... - Int W_n since $W_r \subset \Delta$; then g is defined on $Bd\Delta$. Since $BdW_r \subset \Delta$, g_{BdW_r} is always defined D.

Proof of the Corollary to (1.1). Let $\psi:S^1 \to V$. If V^* is simply connected, we can use (1.1) to show that ψ shrinks to a point in V only if rng ψ misses A_0 . Evidently in order to apply (1.1), it is sufficient to show that ψ is homotopic in V to a mapping $\psi: S^1 \to V - A_0$. We use ' \simeq ' to mean 'is homotopic in V to'. To construct ψ' : using an argument like that of b) in the proof of (1.1), cover $\psi[S^1] \cap A_0$ with dogbones J_1, \ldots, J_q which are disjoint and lie in V. Dogbones J_{sj} , j = 1, 2, 3, 4, are defined just as in b) of (1.1) so that $\bigcup_{r,j} J_{r,j}$ covers $\psi[S^1] \cap A_0$ and each $J_{s,j}$ lies in a cube $M_{ri} \subset V$ (the construction of the J_{si} here is not identical to that of the J_{si} in b) of (1.1), but the construction here is easier since we need not consider sets in \mathcal{D}). We assume that some point zexists in $\psi[S'] \cap (E^3 - \bigcup_{sj} Int J_{sj})$, for otherwise, since S^1 is connected and the Int J are separated, rng ψ lies in one J M si, the proof is concluded by shrinking ψ to a point in M V . Now choose $\delta > 0$ so that if x and y are closer together on S^1 than the distance δ , then $\psi(x)$ and $\psi(y)$ are closer together than the

distance from $\Psi[S'] \wedge A_0$ to $E^3 - \bigcup_{s,j} J_{s,j}$ (remember that $\bigcup_{s,j} J_{s,j}$ is a neighbourhood of $\psi[S^*] \cap A_0$ so that this distance is positive). If every point of S¹ lies closer than δ to $\psi^{-1}[E^3 - \bigcup_{s_1} J_{s_1}]$, then by the definition of δ no point of S^1 maps under ψ into A_0 , and we can let $\psi' = \psi$. If some point of S¹ fails to lie withing δ of $\psi^{-1}[E^3 - \bigcup_{s_i} J_{s_i}]$, then there is an open interval; e_1 in $\psi^{-1}[\bigcup_{s_1}^{U} \text{Int J}_{s_1}]$ such that the length of e_1' is greater than δ . Let e_1 be the largest open interval such that $e_1 \subset e_1 \subset \psi^{-1}[\bigcup_{sj} \text{ Int } J_{sj}]$. Then e_1 is a closed interval of length greater than δ whose end points p_1 and q_1 map into Bd \bigvee_{si} J_{si} by the usual continuity argument. (Since S^1 is a circle, we must make an easy allowance for the possibility that $\psi(p_1) = \psi(q_1) = z$.) Because the J_{sj} are separated and $\psi[\overline{e}_1]$ is connected, $\psi[e_1]$, lies in the interior of just one J_{sj} which we will call R_1 , while $\psi(p_1)$ and $\psi(q_1)$ lie in Bd R₁. Define the mapping $\psi_1: S^1 \to E^3$ so that $\psi = \psi_1$ on $S^1 - e_1$ while $\psi_1 | e_1$ is a path in (connected) Bd R₁ with end points $\psi(p_1)$ and $\psi(q_1)$ (this is well defined because p_1 and q_1 map into Bd R_1 under ψ). Both ψ and ψ_1 are paths in V and $\psi_{|e_1} \simeq \psi_1|_{e_1}$ because they share end points and both map into the same cube $M_{sj} \supset R_1$ with $M_{sj} \subset V$. Evidently $\psi \simeq \psi_1$. Since rng $\psi_1|_{e_1}$ \subset Bd R_1 \subset E³ - A₀, the homotopy has moved images of points in e_1 away from A_0 . We now look for an open interval e_2 in $S^1 - e_1$ where \mathbf{e}_2^{\prime} is of length greater than δ and such that $\mathbf{e}_2^{\prime\prime}$ maps into \bigcup_{s_1} Int J_{s_1} under ψ_1 (and in fact under ψ , since $\psi = \psi_1$ on $S^1 - e_1$). If e_2' does not exist, let $\psi_1 = \psi_1'$. If e_2' exists, then there is an open interval e_2 of maximal length such that $e_2 \subset e_2 \subset S^1 = \overline{e_1}$ and $\psi_1[e_2] \subset \bigcup_{s_1} \text{ Int } \mathbb{J}_{s_1}$. The end points p_2 , q_2 of e_2 map under ψ_1 into $E^3 - \bigcup_{si}$ Int J_{si} , either because of the maximality of e_2 if the

end point is in $S^1-\overline{e_1}$, or, if the end point is in Bd $(S^1-\overline{e_1})=Bd$ e_1 , because ψ [Bd e_1] \subset Bd R_1 . By a continuity argument, $\psi_1[e_2]$ lies in the interior of some one J_{sj} called R_2 and $\psi_1(p_2)$ and $\psi_1(q_2)$ lie in Bd R_2 . Define $\psi_2:S^1\to E^3$ so that ψ_2 agrees with ψ_1 on S^1-e_2 (note that this means that ψ_2 agrees with ψ on $S^1-e_1-e_2$), and so that $\psi_2[e_2]$ is a path in Bd R_2 with end points $\psi_1(p_2)$ and $\psi_1(q_2)$. By a previous argument, $\psi_2\simeq\psi_1\simeq\psi$. Note that the fact that $\psi_2=\psi$ on $S^1-e_1-e_2$ means that $\psi_2=\psi$ on the end points of both e_1 and e_2 .

In general, suppose that mappings $\psi_1 \simeq \ldots \simeq \psi_{r-1}$ of s^1 into V, intervals $e_1, \dots e_{r-1}$ and components $R_1, \dots R_{r-1}$ of $\bigcup_{sj} J_{sj}$ have been defined so that each $e_s \subset S^1 - e_1 - \dots - e_{s-1}$, $s = 1, \dots r - 1$, and for each ψ_s , $\psi_s = \psi_{s-1}$ on $s^1 - e_s$, $\psi_s[\overline{e}_s] \subset Bd$ $R_s \subset E^3 - A_o$. Now look for an open interval $e_r \in S^1 - e_1 - e_2 - \dots - e_{r-1}$ such that the length of e_r is greater than δ and $\psi_{r-1}[e_r] \subset \bigcup_{si} \text{Int } J_{si}$, or equivalently $\psi_{r-1}[e_r] \subset \operatorname{Int} R_r$, where R_r is a single J_{si} (and hence lies in a cube $M_{si} \subset V$). If there is no such interval, let ψ_{r-1} be ψ^{\prime} . If e_{r}^{\prime} exists, then let e_{r}^{\prime} be the largest open interval in $S^1-\overset{-}{e_1}-\ldots-\overset{-}{e_{r-1}}$ such that $\psi_{r-1}[e_r]\subset \operatorname{Int}\, R_r$. We know that ψ_{r-1} carries the end points p_r , q_r of e_r into E^3 - Int R_r by the maximality of e_r if the end point is in $S^1 - e_1 - \dots - e_{r-1}$, or, if the end point is in $Bd(S^1 - \overline{e}_1 - \dots - \overline{e}_{r-1}) \subset \overline{e}_1 \cup \dots \cup \overline{e}_{r-1}$, because ψ_{r-1} carries $e_1 \cup \ldots \cup e_{r-1}$ into $\bigcup_{s_1} \operatorname{Bd} J_{s_1}$. (To see that $\psi_{r-1}[\overline{e}_s]$ lies in some Bd J_{si} : for $s = 1, \ldots r-1, \psi_s[\overline{e}_s] \subset Bd R_s$ by construction. ψ_{s+1} agrees with ψ_s on $s^1 - e_{s+1} \supset e_s$ since

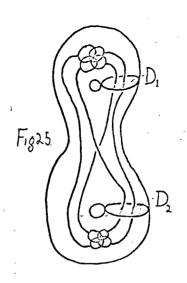
Since each e_r is of length greater than δ and $r\delta$ must be less than the circumference of S^1 , (it is easily checked that the e_r are disjoint) the sequence $\psi_1, \ldots, \psi_r, \ldots$ ends at ψ_k . Let ψ be ψ_k . We know that $\psi_k \colon S^1 \to V$ because each ψ_r does this; and $\psi \simeq \psi_1 \simeq \ldots \simeq \psi_k = \psi$ in V. Before we can show that $\operatorname{rng} \psi$ misses A_0 , we must show that $\psi = \psi' = \psi_k$ on $S^1 - e_1 - \ldots - e_k$. To see this:

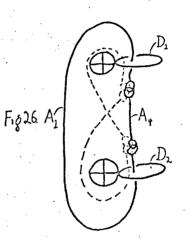
$$\psi_1 = \psi$$
 on $S^1 - e_1$,

$$\psi_2 = \psi_1$$
 on $S^1 - e_2$ and $\psi_2 = \psi - on S^1 - e_1 - e_2$,

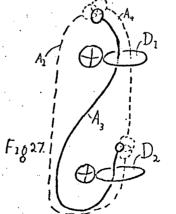
$$\psi_k = \psi_{k-1}$$
 on $S^1 - e_k$ and $\psi_k = \psi$ on $S^1 - e_1 - \dots - e_k$.

It is now easy to show that $\operatorname{rng} \psi' = \operatorname{rng} \psi_k$ misses A_0 , for ψ_k carries every $\overline{\operatorname{e}}_s$ into $\operatorname{Bd} \operatorname{R}_s \subset \operatorname{E}^3 - \operatorname{A}_0$ by a previous result; and if $x \in \operatorname{S}^1 - \overline{\operatorname{e}}_1 - \ldots - \overline{\operatorname{e}}_k$, then x lies within a distance δ of a point y such that $\psi_k(y) \in \operatorname{E}^3 - \bigcup_{sj} \operatorname{Int} \operatorname{J}_{sj}$. We can assume that $y \in \operatorname{S}^1 - \operatorname{e}_1 - \operatorname{e}_2 - \ldots - \operatorname{e}_k$ because otherwise $y \in \operatorname{e}_1 \circ \ldots \circ \operatorname{e}_k$ and





Ay may be shrunk;
Az, Az, Az now miss at
I cast one Dz. But Az
must hit both.



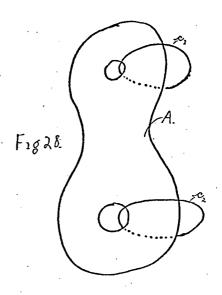
A₁ l A₄ may be moved so that each of A₁, A₄.

Misses one D₁.

But now A₃ meets both.

some point y' of $\operatorname{Bd}(e_1 \cup \ldots \cup e_k)$ $\operatorname{S}^1 - e_1 - \ldots - e_k$ lies closer to x than y does. Since ψ_k carries $y' \in \operatorname{Bd}(e_1 \cup \ldots \cup e_k)$ $\subset e_1 \cup \ldots \cup e_k$ into $\bigcup_{sj} \operatorname{Bd} J_{sj} \subset \operatorname{E}^3 - \bigcup_{sj} \operatorname{Int} J_{sj}$, we could have originally chosen y' instead of y. But if both x and y' lie in $\operatorname{S}^1 - e_1 - \ldots - e_k$, then since $\psi = \psi_k$ on $\operatorname{S}^1 - e_1 - \ldots - e_k$ and $\psi_k(y') \subset \operatorname{E}^3 - \bigcup_{sj} \operatorname{Int} J_{sj}$, $\psi(x) = \psi_k(x)$ lies so near to $\psi(y') = \psi_k(y')$ that $\psi(x)$ misses A_0 by our definition of δ \square .

In his paper [12], Bing was concerned with an interesting property of G which we will make use of here and in Ch IV. The formidable aspect of G lies in what might be called its 'topological idiom', as shown in fig 25: four double ended lassos strung in a special way inside a 2-holed torus. Bing's intent in using this idiom to construct G was to utilize this property: let D_1 , D_2 , be the planar disks shown in fig 25. Then, no matter how the four lassos are deformed (provided that they remain linked and stay in the interior of the 2-holed torus), some one lasso will hit both $\,\mathrm{D}_{1}\,$ and $\,\mathrm{D}_{2}\,$. Figs 26, 27 show unsuccesful attempts by the lassos to avoid this necessity, and there is a proof of a very similar idea in §7 of [12]. Bing hoped to show that this property was induced through the construction of G in the following sense: assume that fig 25 shows D_1 , D_2 in relation to the first stage of the construction of G, then, no matter how A_1 , A_2 , A_3 , A_4 were deformed, one of these, say A_1 , would hit both of D_1 , \boldsymbol{D}_2 . Additionally, however, it might turn out that for any deformation of A_2 inside A, one of the 16 A_{1k} would hit both D_1 and D_2 , and so on for the 64 $\,^{\mathrm{A}}_{\mathrm{i}k\ell}$ etc. Bing found that there was no easy



proof of this (see §7 of [12]); however he was able to define a property which he called Q on the dogbones of the decomposition and show that A had this property. If a dogbone had property Q, this implied trivially that it intersected both of D_1 , D_2 ; at the same time it could be shown that if a dogbone B had property Q, then one of the four dogbones of the next stage of the dogbone construction lying in B would have property Q. Evidently there would be a descending intersection chain of dogbones each with property Q and the limit of the chain would be a big element of G which touched both D_1 and D_2 . We can express this idea in a slightly different way:

(2.1) (Bing). Let D_1 , D_2 be topological disks whose boundaries C_1 , C_2 lie in A and link the upper and lower eyes respectively of A as shown in fig 28. Then either D_1 metts D_2 in A, or some big element of G meets both D_1 and D_2 .

We will refer frequently to fig 28, which shows the relation—ship of c_1 , c_2 to A. Strictly speaking, we take c_i , i=1,2, to be an embedding of s^1 in e^3 ; however we frequently will confuse the embedding with the circle which is its range (at a the same time reserving the right to write $rng c_i$ when we wish to make the distinction clear).

Bing showed that (2.1) was inconsistant with the existence of a homeomorphism between $\mathcal D$ and E^3 (Th 12 of [12]). In this paper we will be interested in this conjecture:

(2.2). Let Δ be a 2 - simplex. For i=1, 2, let $f_i: \Delta \to E^3$ so that $f_{i\mid Bd\Delta}=c_1$ and $f_{2\mid Bd\Delta}=c_2$ are paths whose ranges lie in E^3 - A and which will not shrink to a point in the complements of the

upper and lower eyes respectively of A . Then either $f_1[\Delta]$ and $f_2[\Delta]$ intersect in A, or some big element of G meets both $f_1[\Delta]$ and $f_2[\Delta]$.

In (2.2), we replace the disks D_i of (2.1) with singular disks $f_i[\Delta]$. The conjecture is plausible and lacks earthshaking surprise. It is interesting because it leads directly to the following topological property of \mathcal{D} :

(2.3). If (2.2) is true, then \mathcal{D} fails to possess arbitrarily small simply connected open neighbourhoods about any big point.

The conclusion of (2.3) is called <u>Curtis' conjecture</u> (see 3, §6), and (2.3) reduces it to the somewhat more plausible conjecture (2.2). The remainder of this chapter consists of a proof of (2.3). The pleasures of (2.2) will be deferred to Chapters III and IV.

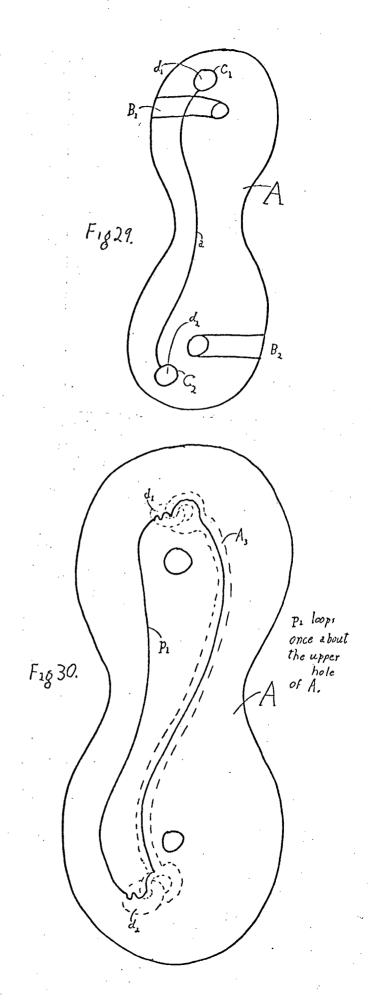
Proof of (2.3). Suppose that Λ is a big element of G and that in \mathcal{D} , $\Lambda^* = \phi[g]$ lies in a simply connected open neighbourhood V^* such that $\Lambda^* \subset V^* \subset \Lambda^*$. Clearly $\Lambda \subset V \subset \Lambda$, and V is open in E^3 . We could write $\Lambda = A \cap A_j \cap A_{jk} \cap \ldots$ for some sequence of dogbones A, A_j , A_{jk} , \ldots . By the Corollary to (1.1) (of lemma 1 of [2]), V is simply connected if V^* is. Thus our assumption implies that V is simply connected. We will demonstrate that this is false by showing that $\Lambda \subset V \subset A$ with V simply connected, implies that the upper eye ℓ and the lower eye m of A shrink to a point in A. We define an upper (lower) principal path of A_j to be a mapping of S^1 into Int A_j which is homotopic in A_j to the upper eye ℓ_j (the lower eye m_j) of A_j . Upper and lower principal paths of other dogbones, including

A, are defined analogously; this a mapping of S^1 into Int A_{jk} which is homotopic in A_{jk} to $h_{jk}[\ell]$ is an upper principal path of A_{jk} . As usual, we will often confuse the mapping with its range. We know that A, A_j , A_{jk} , ... is a neighbourhood system of Λ and, by a previous remark, that some member A_{jk} ... rs of the system lies in V. However this fact plus the following lemma leads to a contradiction.

Lemma for (2.3). If one of A_j , j=1,2,3,4 contains an upper principal path e_1 and a lower principal path e_2 which intersect and lie in V, then A contains upper and lower principal paths which intersect and lie in V. In general, if A_j ... rs contains intersecting upper and lower principal paths which lie in V, then so does A_j ... r

To apply the lemma we look at the neighbourhood $A_j \dots rs$ which we know to be a neighbourhood of A in V. Obviously $V \cap A_j \dots rs$ contains intersecting upper and lower principal paths of $A_j \dots rs$ since any intersecting principal paths will qualify. The lemma implies that the dogbone $A_{ij} \dots r$ contains intersecting principal paths in V. Repeated application of the lemma leads to the conclusion that A contains an upper principal path which lies in V. Since $V \subseteq A$, V is simply connected, and the upper principal paths of A are all homotopic to ℓ in A, therefore ℓ must shrink to a point in A. This is clearly false from fig 16a. Thus the proof of (2.3) will be complete when we have proved the lemma.

Proof of the lemma for (2.3): Simplified version. Suppose that \mathbf{e}_1 and \mathbf{e}_2 lie in \mathbf{A}_1 . The following outline reflects our original intuition of the proof. Although the 'proof' we give now is



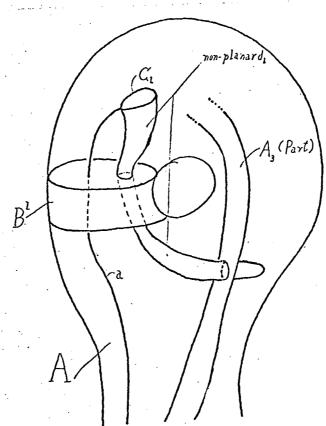
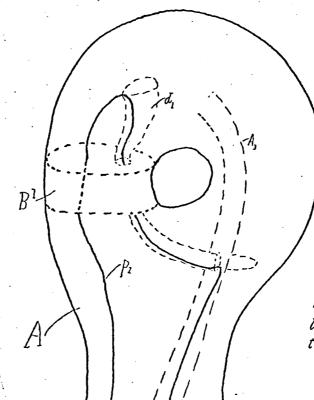


Fig 3la.

A non-planar di may meet A3 as shown,
It is now possible to construct p2 so that p2 does not loop about the upper hole of A 2s required.

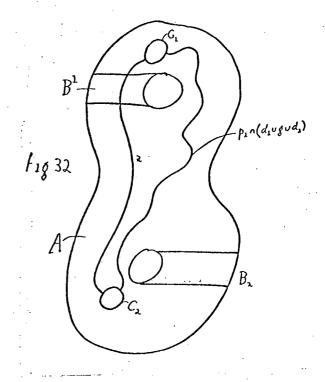


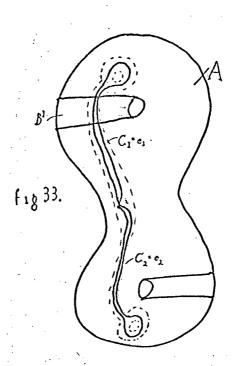
F1g 31 b.

The curve p_1 may lie in a ud₁ uA₃ ug as constructed in fig 31 a, and fail to loop once about the hole.

simple minded and needs much patching, we give the crude version because we think that it clarifies basic ideas which tend to be submerged in the final version of the proof. Suppose that by good fortune the paths e_1 , e, take the form of the double ended lasso J in fig 29. J consists of circles $\mathbf{C_1}$, $\mathbf{C_2}$ and connecting arc a as shown. The circle $\mathbf{C_1}$ will not shrink to a point in the complement of the upper eye of A_2 . Similarly C_2 will not shrink to a point in the complement of the lower eye of A_3 . We also pretend that J lies in V and that disjoint planar disks d_1 , d_2 bounded by C_1 , C_2 also lie in V . By (2.1) some big element g in A_3 meets both d_1 and d_2 ; and g lies in $V \supset d_i$ since V contains every element of G that it intersects (remember that $\, { t V} \,$ is the pre-image of an open set in $\, { t D} \, { t)} \, . \,$ We can now construct the upper principal path p_1 shown in fig 30 from parts of V lying in a, d_1 , d_2 , g . A similar procedure using A_2 instead of A_3 will yield the lower principal path p_2 . The paths p_1 , p_2 intersect in A so that $\mathbf{p}_1 \cup \mathbf{p}_2$ is the set required by the conclusion of the lemma.

The above 'proof' is far to easy and will fail if we allow d_1 , d_2 to be non-planar, for then p_1 may not be a principal path as figs 31a, b show. We ensure that p_1 makes one circuit about the upper hole of A by trapping $p_1 \cap a$ in the cube $A - B^2 - B^2$ (which is easy) and $p_1 \cap (d_1 \cup g \cup d_2) = q$ in the cube $A - B^1 - B_2$ (see fig 32). This last step is hard since one would fear that the connectivity would be spoilt by parts of $d_1 \cup d_2$ projecting from the cube. The trick of controlling the homotopy class of p_1 by constructing certain arcs in





cubes only works if the C_i lie in $A - B^1 - B_2$. But if we use the obvious candidate for J, viz. $C_i = e_i$ with arc a degenerate, then fig 33 shows that this may not happen, and in fact J cannot usually be $e_1 \cup e_2$. However we show that, provided that intersecting principal paths exist in $A_1 \cap V$, there is a double ended lasso (perhaps with singularities) in $A \cap V$ which has just the properties which we assigned to J. We will now give the final version of the proof of the lemma for (2.3). This proof uses the ideas of the earlier crude version, but incorporates the various improvements suggested in this paragraph.

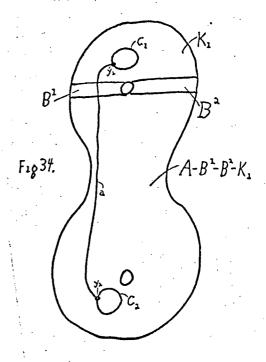
Outline of final version of proof. We first give the proof assuming that $e_1 \subset e_2 \subset A_1$, then indicate alterations in the case that $e_1 \cup e_2$ lies in A_2 , A_3 or A_4 . a) Let e_1 , e_2 be upper and lower principal paths of A_1 which lie in V and intersect at least at p. We follow the sketch of the 'proof' already given, but as previously explained, we cannot use $e_1 \cup e_2$ for J in fig 32. We construct $J = C_1 \cup C_2 \cup A$ so as to satisfy five properties i), ... v). Sometimes we will regard C_1 as a mapping (not necessarily an embedding) of S^1 and sometimes as the range of this mapping. The set J must satisfy the following properties: for i = 1, 2,

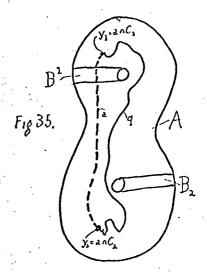
i) $C_1[Bd\Delta] \subset V \cap Int(A - B^1 - B_2)$

ii) rng C_i misses A_3 .

iii) C_1 (C_2) fails to shrink to a point in E^3 - ℓ_3 (E^3 - m_3) .

iv) There is a point $y_1 \in \operatorname{rng} C_1 \cap \operatorname{rng} e_1 \cap K_1$ and a point y_2 in $\operatorname{rng} C_2 \cap \operatorname{rng} e_2 \cap A - B - B_2 - K_1$ (recall that K_1 is the topological cube which is the closure of the upper component of $A - B^1 - B^2$, see fig 34).



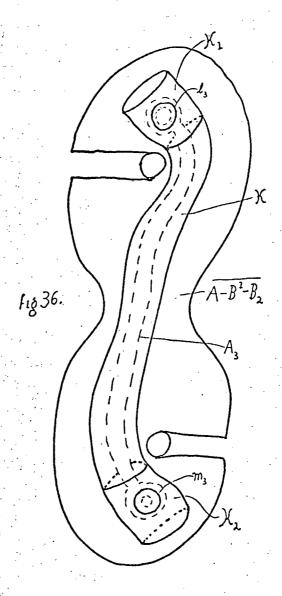


v) the arc $a \in V \cap A_1$ and the points y_1, y_2 of iv) are the end points of a.

The idea of ii) and iii) is that we want the C_i to act like the circles $\mathbf{c_1, \ c_2}$ in fig 28 with respect to $\mathbf{A_3}$. Property iv) provides the endpoints of a 'above and below B^1 ' . This plus v) and the fact that the C_1 are trapped in $A - B^1 - B_2$ allows us eventually to construct an upper principal path of A which winds one around the upper hole of A. This happens because we will join y_1 and y_2 by a path like $q \subset A - B^1 - B_2$ in fig 35. b) For i = 1, 2, let $f_i: \Delta \to V$ so that $f_i \mid \dot{B} d\Delta = C_i$. Using (1.2) and (1.3), obtain a new mapping \overline{f}_{i} which agrees with f_{i} on Δ - Int W_1 - ... - Int W_n , where Int W_r , $r = 1, \ldots, n$, are holes in Δ ; in particular $\overline{f}_i = f_i$ in $Bd\Delta$. The $\overline{f}_i[W_r]$ may leave V (!) but this does not harm the proof. c) By (2.2), $\overline{f}_1[\Delta]$ and $\overline{f}_2[\Delta]$ either intersect in A_3 or hit the same big element g in A_3 . There is a path $\ \mathbf{q}$ from a 0 C $_1$ to a 0 C $_2$ which resembles $\ \mathbf{q}$ in fig 35 and lies in V and in Int $(A - B^1 - B_2)$. The path q travels to A_3 in $\overline{f}_1[\Delta]$, passes from $\overline{f}_1[\Delta]$ to $\overline{f}_2[\Delta]$ in A_3 either at the intersection of $\overline{f}_1[\Delta]$ and $\overline{f}_2[\Delta]$ or using the element g, and then proceeds to a \cap C_2 by means of $\overline{f}_2[\Delta]$. d) The path which begins at a \cap C_1 , travels to a n c_2 on q and returns to aa c_1 on a, is an upper principal path of $\,A\,$ which lies in $\,V\,$. $\,e)\,$ The lower principal path of $\,A\,$ in V may be constructed as in a), b), c), d) above, using A_2 and $A - B_1 - B^2$ instead of A_3 and $A - B^1 - B_2$. f) If k = 2, 3, 4 the lemma remains true.

Details of Proof. Suppose that e_1 , e_2 are upper and lower

principal paths respectively of A_1 which lie in V and intersect at p . For i = 1, 2, since e_i shrinks to a point in V, there are mappings $e_i: \Delta \to V$ such that $e_{i \mid Bd\Delta} = e_i$. We do not claim that rng $\frac{1}{e}$ lies in A_1 or even A . We use (1.2), taking f, S, to be e_1 , $e_1^{-1}[A - B^1 - B_2]$. Thus there are disjoint disks W_1 , ... W_m in Δ such that $e_1^{-1}[A - B^1 - B_2] \subset W_1 \cup W_2 \cup ... \cup W_m$, each point x of Bd W_r lies near $e_1^{-1}[A - B^1 - B_2]$, and $x \in Bd W_r$ misses $e_1^{-1}[A - B^1 - B_2]$ if $x \notin Bd\Delta$. Those W_r which hit $Bd\Delta$ are called W_1 , ... W_n ; those W_r which miss Bd Δ are W_{n+1} , ... W_m (with obvious adjustments of one or the other class does not exist). We now apply (1.3) with g taken to be the restriction of e_1 to Δ - Int W_1 - ... - Int W_m . For $r = n+1, \dots m, e_{1|Bd W_r} = g_{Bd W_r} maps into Ext(A - B^1 - B_2)$ because Bd W_r misses Bd Δ for r > n and hence misses $e_1^{-1}[A - B^1 - B_2]$. Thus for r = n+1, ... m, $e_{1|Bd W_r}$ shrinks to a point in $Ext(A - B^1 - B_2)$ which is the exterior of a cube in E^3 ; and we can let $Ext(A - B^1 - B_2)$ be $P_{n+1} = \dots = P_m$ in the hypothesis of (1.3). There is no chance that $e_{1|Bd|W_r}$ is C_1 for r > n, since ℓ_3 misses $Ext(A - B^1 - B_2) = P_r$. We suspect that c_1 is an $e_{1|Bd\ W_r}$ for $r \le n$. Assume that every $\frac{1}{e_{1|Bd W_r}} = g_{Bd W_r}$ will shrink to a point in $E^3 - \ell_3$. Use (1.3) again, letting $P_1 = P_2 = \dots = P_n$ be $E^3 - \ell_3$. Then with $Ext(A - B^1 - B_2)$ taken to be $P_{n+1} = \dots = P_m$, $g_{\mid Bd\Delta} = e_1$ will shrink to a point in rng g \cup P $_1$ \cup ... \cup P $_n$ \cup P $_{n+1}$ \cup ... \cup P $_m$. Each P $_r$ misses ℓ_3 either by definition of because P_r misses $A - B^1 - B_2$. And rng g misses



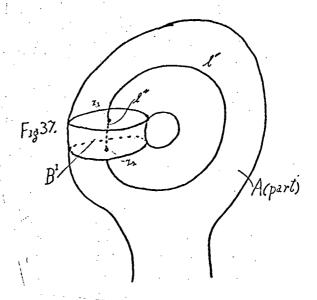
 ℓ_3 as well; for $g = e_1 | \tilde{\Delta} - Int W_1 - ... - Int W_m$ and from (1.2) ii, iv, the only points of Δ - Int W_1 - ... - Int W_m which can map into $A - B^1 - B_2$ are those in $Bd\Delta$. Such points are in dom e_1 and map into A_1 . Hence rng g $\subset A_1 \cup Ext(A - B^1 - B_2) \subset E^3 - \ell_3$. Therefore $g_{\mid Bd\Delta} = e_1$ shrinks to a point in rng g $\cup P_1 \cup \ldots \cup P_m \subset E^3 - \ell_3$, which contradicts the fact the $\,e_1^{}\,$ is an upper principal path. Thus it is false that every $[e]_{Bd\ W_r}$, $r \leq n$, shrinks to a point in $E^3 - \ell_3$. Let C_1 be one of the e_1 Bd W_2 which fails to shrink to a point in E^3 - ℓ_3 . As regards C_1 : the above argument plus the fact that rng $C_i \subset rng \stackrel{-}{e_i} \subset V$ shows that iii) is true; ii) is true because from (1.2) iv, every point x in Bd W_r is either in Int Δ , in which case $C_1(x) \in E^3 - \overline{(A - B^1 - B_2)}$ $\subset E^3 - A_3$, or $x \in Bd\Delta$, when $C_1(x) = e_1(x) \in A_1 \subset E^3 - A_3$. In general,i) is not true because some candidates for $C_1(x)$ lie outside of $A - B^1 - B_2$ as we have just seen. However we can assume that $C_1[Bd\ W_r]$ lies in Int $(A - B^1 - B_2)$ by the following argument: By (1.2), we assume that dom C_1 (which is one of the Bd W_r) lies so near $\overline{e_1^{-1}}[A - B^1 - B_2]$ that rng C_1 lies within ε of $A - B^1 - B_2$ (remember that $C_1 = g = e_1$ on dom C_1). In this paragraph a) so far, we could have replaced $A - B^1 - B_2$ by a cube $K \supset A_3$ such that an ϵ -neighbourhood of K lies in Int $(A - B^1 - B_2)$. Such a cube is shown in fig 36. If this had been done, we would have $\ \operatorname{rng}\ \operatorname{\textbf{C}}_1$ in the $\ \epsilon$ -neighbourhood of K, i.e. rng $C_1 \subset Int (A - B^1 - B_2)$. We assume that this was done and that rng $C_1 \subseteq Int (A - B^1 - B_2)$. Proof of iv): In (1.2) iv, each Bd W misses S (in (1.2)) except where Bd W hits Bd Δ . In the

present context, with $e_1^{-1}[K]$ for S (i.e. continuing to use K for $A - B^1 - B_2$), the domain of C_1 is a Bd W_r and $C_1[Bd W_r - Bd\Delta]$ misses K. To show that there is a $y_1 \in \text{rng } C_1$ rng $e_1 - K_1 = C_1[Bd W_r] \cap e_1[Bd\Delta] \cap K_1 = C_1[Bd W_r \cap Bd\Delta] \cap K_1$, assume that $C_1[Bd\Delta \cap Bd W_r] \cap K_1 = \emptyset$. In fig 36, note the two cubes K_1 , K_2 which are placed so that $\ell_3 \subseteq K_1 \subseteq K \cap K_1$ and $m_3 \subseteq K_2 \subseteq K \cap K_2$. Then $C_1[Bd W_r \cap Bd\Delta] \cap K_1 = \emptyset$ because $K_1 \supset K_1$; and $C_1[Bd W_r - Bd\Delta] \cap K_1 = \emptyset$ because $C_1[Bd W_r - Bd\Delta] \cap K_1$

We repeat the entire procedure of this paragraph a) taking e_2 , e_2 , m_1 , m_3 , K_2 , $A - B^1 - B_2 - K_1$, for e_1 , e_1 , ℓ_1 , ℓ_3 , K_1 , K_1 . This is just the preceding argument 'upside down' and constructs the path $c_2 \ni y_2$ as required. The only unexpected thing is the use of the cube $A - B^1 - B_2 - K_1$ for the original cube K_1 ; this reflects the fact that y_1 should be found in K_1 and y_2 not necessarily in K_2 but merely 'in A_1 and below B^1 '. We now have y_1 and y_2 as required by iv). To construct a_1 join y_1 to p_1 by a path p_2 and p_3 in rng p_4 and p_4 in rng p_4 and p_5 in rng p_6 and p_6 in rng p_6

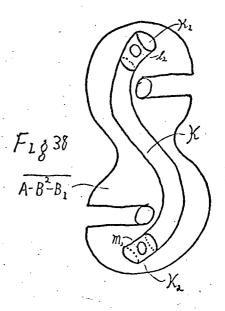
(unless they intersect), but we are not sure that there is a connected set in $f_i[\Delta]$ that will join g and y and stay in $A - B^1 - B_2$, so that it is not yet possible to build $q \subset V \cap (A - B^1 - B_2)$ as in fig 35. By (1.2), taking S to be $Ext(A - B^1 - B_2)$, there are disks W_1^i , ... W_n^i in Δ such that $f_i^{-1}[\overline{\operatorname{Ext}(A-B^1-B_2)}] \subset W_1^i$, \cup ... \cup W_n^i . Since BdA misses $f_i^{-1}[Ext(A - B^1 - B_2)]$ (because $f_i[BdA] \subset Int(A - B^1 - B_2)$), Δ - Int W_1^i - ... - Int W_n^i is a disk with holes. We assume that each Bd W_r^i lies so near $f_i^{-1}[\text{Ext}(A-B^1-B_2)]$ that $f_i[Bd\ W_r^i]$ lies close to $Ext(A - B^1 - B_2)$. Since we know exactly what $A - B^1 - B_2$ looks like, we can construct an ε -neighbourhood N of Ext(A - B¹ - B₂) so that N is simply connected. We can assume that each $f_i[BdW_r^i] \subset N$; then $f_{i}[BdW_{r}^{i}]$ shrinks to a point in N; and by (1.3), taking $N = P_1 = P_2 = \dots = P_n$, and $g = f_{i|\Delta} - Int W_1^i - \dots - Int W_n^i$, there. is a mapping $\overline{f}_i: \Delta \rightarrow \text{rng g} \vee P_1 \cup \ldots \cup P_n = f_i[\Delta - \text{Int } W_1^i - \ldots - \text{Int } W_n^i] \cup N$ such that $\overline{f}_i = f_i$ on $\Delta - Int W_1^i - \dots - Int W_n^i$. In particular $\overline{f}_{i} = f_{i}$ on Bd \triangle . It is important that $\overline{f}_{i}[\triangle - Int W_{1}^{i} - ... - Int W_{n}^{i}]$ = $f_i[\Delta - Int W_1^i - ... - Int W_n^i] \subset V$.

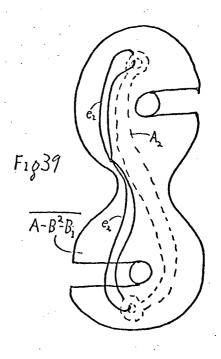
c) Since $\operatorname{rng} C_i = f_i \ [\operatorname{Bd}\Delta] = \overline{f_i} \ [\operatorname{Bd}\Delta]$ misses A_3 and fails to shrink to a point in the absence of the appropriate eyes of A_3 , by (2.2), $\overline{f_1}[\Delta]$ and $\overline{f_2}[\Delta]$ either intersect in A_3 or hit the same big element λ in A_3 . We can combine these ideas by saying ' $f_1[\Delta]$ and $f_2[\Delta]$ meet the same element λ in A_3 ' and allowing λ to be either a big element or a small element. Since for a small ε , N misses A_3 , $\lambda \wedge \operatorname{rng} \overline{f_i}$ must lie in $\operatorname{rng} \overline{f_i} - N \subset \overline{f_i}[\Delta - \operatorname{Int} W_1^i - \ldots - \operatorname{Int} W_n^i] \subset V$, and $f_1^{-1}[\lambda] \subset \Delta - \operatorname{Int} W_1^i - \ldots - \operatorname{Int} W_n^i$, which we saw was a disk with



holes and which contains $C_{\mathbf{i}}^{-1}(y_{\mathbf{i}}) \subseteq \text{dom } C_{\mathbf{i}} = \text{Bd}\Delta$. Since $\lambda \cap f_{\mathbf{i}}[\Delta]$ and $y_{\mathbf{i}}$ lie in the image under $\overline{f}_{\mathbf{i}}$ of a disk with holes which maps into V, there is a path $v_{\mathbf{i}}$ which joins $y_{\mathbf{i}}$ and λ in V. Futhermore $v_{\mathbf{i}} \subseteq A - B^{1} - B_{2}$, because $v_{\mathbf{i}}$ may be constructed in $\overline{f}_{\mathbf{i}}[\Delta - \text{Int } W_{\mathbf{i}} - \dots - \text{Int } W_{\mathbf{n}}] = f_{\mathbf{i}}[\Delta - \text{Int } W_{\mathbf{i}} - \dots - \text{Int } W_{\mathbf{n}}]$ which misses $\text{Ext}(A - B^{1} - B_{2})$ by the construction of the $W_{\mathbf{i}}$. Hence $v_{\mathbf{i}} \subseteq V \cap (A - B^{1} - B_{2})$. Let q be a path joing $y_{\mathbf{i}}$ and $y_{\mathbf{i}}$ in $v_{\mathbf{i}} \cup \lambda \cup v_{\mathbf{i}}$. Clearly $q \subseteq V \cap (A - B^{1} - B_{2})$.

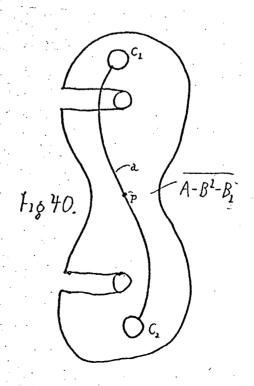
d) We will show that the path ξ_1 which travels from y_1 to \mathbf{y}_2 in \mathbf{q} and returns to \mathbf{y}_1 in \mathbf{a} is an upper principal path of A in V by showing that $\xi_1 \subset A \cap V$ and that ξ_1 is homotopic to ℓ in A . Let ℓ be decomposed into two paths ℓ and ℓ " such that ℓ " \subset B^1 and ℓ \subset E^3 - B^1 . We assume that ℓ pierces Bd B^1 in just two points z_1 , z_2 as shown in fig 37. We can do this because Bd ${ t B}^1$ is horizontal near ℓ and because ℓ can be a nice circle. Construct arcs z_1 y_1 and z_2 y_2 in the cubes K_1 and $A - B^{\frac{1}{2}} B_2 - B^2 - K_1$ respectively. (The idea here is that both z_i y_i will lie in $A - B^1 - B_2$, the cube which locates $A_3 \supset q$, and in $A - B^2 - B_2$, the cube which locates $A_1 \supset a$). The path which begins at z_1 and travels to z_2 through z_1 y_1 , a, and z_2 y_2 is homotopic in the cube $A - B^2 - B_2$ to ℓ " . The path which begins at \mathbf{z}_2 and travels to \mathbf{z}_1 via \mathbf{z}_2 y₂, q, and $z_1 y_2$ is homotopic in the cube $\overline{A - B^1 - B_2}$ to ℓ . Combining homotopies, the path ξ_1 which begins at z_1 , travels to z_2 in z_1 y_1 v q v z_2 y_2 and returns to z_1 in z_2 y_2 v a v z_1 y_1 , is





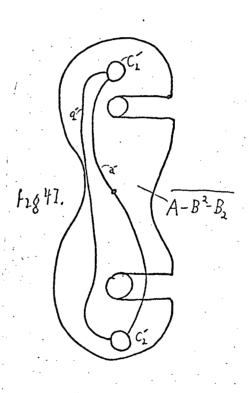
homotopic in A to ℓ . The path ξ_1' is evidently homotopic to ξ_1 . Note that ξ_1 passes through the point p ϵ a. Eventually p will be the 'official' intersection point of the principal paths ξ_1 and ξ_2 of A.

There is no difficulty in altering the argument to construct a lower principal path ξ_2 if one keeps in mind the fact that 'the pictures are different' and that everything in the construction of $\ \xi_1$ must be repeated. We cannot, for example, use the C_{i} from a) because they were defined with respect to $A - B^1 - B_2$ and we must replace $A - B^1 - B_2$ (the cube which located A_3 and 'shaped' the right side of ξ_1) with $A - B^2 - B_1$ which locates A_2 . The idea is to start with e_1 and e_2 as before, but to use A_2 rather than A_3 as suggested in fig 38 which, in a sense, is a replacement for fig 35. The new ξ_2 turns out to contain $p \in e_1 \cap e_2$ just as ξ_1 does; this establishes that $\xi_1 \cap \xi_2 \neq \emptyset$. We begin by finding a new lasso $J' = C_1' \cup C_2' \cup A'$ so that $C_1 \cup C_2 \subset V \cap Int(A - B^2 - B_1)$, $C_1 \subset A_3$, and C_1 contains a point y_i such that $y_1 \in \text{rng } C_1 \cap \text{rng } e_1 \cap (A - B^2 - B_1 - K_2)$ and $y_2 \in \text{rng } C_2 \cap \text{rng } e_2 \cap K_2$. The arc a lies in $V \cap A_1$ and has end points y_1^2 , y_2^2 . One finds $C_1^2 \cup C_2^2 \cup A^2$ by adapting the procedure in a); there is very little more involved than reading m, m_1 , B_1 , B^2 , K_2 , $A - B^2 - B_1 - K_2$ for ℓ , ℓ_1 , B^1 , B_2 , $A - B^1 - B_2 - K_1$, K_2 , and priming every new construction. It will be found that the arc a contains p just as the original a does. For the construction of K', K_1^a , K_2^a , replace fig 36 by fig 39. It is quite easy to adapt b) and



c) by keeping in mind that the important cube is $A - B^2 - B_1$ which replaces $A - B^1 - B_2$ in the construction of ξ_1 . (The point is that in b), c), one must use a cube whose boundary encloses the 'important' dogbone A_2 , see fig 38). Finally we construct a path q which joins y_1' , y_2' in $A - B^2 - B_1$. This plus $a' \subset A - B^2 - B_2$ can be combined into the path ξ_2 which can be shown to be homotopic to m by adapting d) above, decomposing m into paths $m' \subset B_1$ and $m' \subset E^3 - B_1$ etc. The path ξ_2 lies in V by an argument which should appear naturally from the adaptation of a), b), c) to construct ξ_2 ; and ξ_2 is clearly in A. Since the point p lies in both a and a and hence in both ξ_1 and ξ_2 ; therefore $\xi_1 \cap \xi_2 \neq \emptyset$.

f) The proof is now complete if $e_1 \vee e_2$ lies in A_1 , i.e. if j=1. The is no difficulty in constructing a proof for the lemma when j=4 in view of the symmetry of the construction of the A_j . We will give only an outline of the proof for j=2 (and by symmetry for j=3) for these reasons: 1) the details can be filled in along the lines of a) ... e) above, and, 2) the argument in a), ... e) is sufficient to prove the 'meat' of (2.2), viz. that there are uncountably many big points of $\mathcal D$ which fail to possess arbitrarily small simply connected open neighbourhoods, these being images of elements of the form $A \cap A_i$ A_{ij} A_{ijk} ... where i, j, k, ... are chosen from 1 or 4. To construct the principal path a_i when a_i when a_i and the cube a_i and a_i are a_i which we now assume to lie in a_i and a_i and the cube a_i are fig 40) which acts toward a_i just as a_i just under the

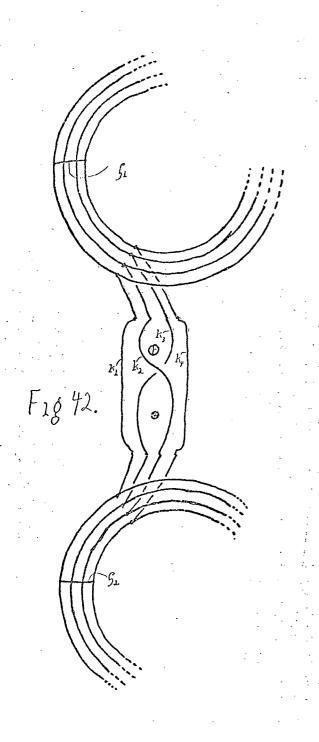


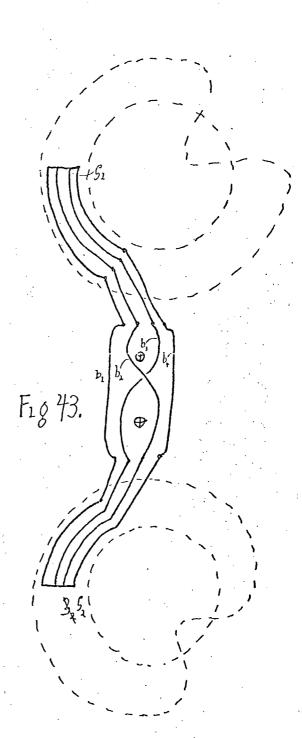
upper eye while $A - B^1 - B_2$ does the same for A_1). Using the argument of a), construct a lasso $J = C_1 \cup a \cup C_2$ which is related to V, $e_1 \cup e_2$, A_4 , and $A - B^1 - B_1$ just as J in a) was related to V, $e_1 \cup e_2$, A_3 , $A - B_1^1 - B_2$. Fig 40 shows the new J. When C_1 ${\tt C_2}$ shrink to a point in V they hit the same element $\,\lambda\,$ of $\,{\tt A_4}\,$ ($\,\lambda\,$ may be a big element or a point). The path q joining the end points of a in $V \cap (A = B^1 - B_1)$ may be constructed by adapting the argument of b), c), and ξ_1 = a U q may be shown to be an upper principal path of A by an argument like that of d). Just as in the case j = 1, the arc a contains a point $p \in e_1 \cap e_2$. Thus $p \notin \xi_1$. To construct the lower principal path ξ_2 when j = 2, we start as before with $e_1 \cup e_2 \subset A_2$, but we use the cube $A - B^2 - B_2 \supset A_1$ and construct $J' = C_1' \cup a' \cup C_2'$ so that J' is related to V, $e_1 \cup e_2$, A_1 , $A - B^2 - B_2$, just as J' in e) is related to V, $e_1 \cup e_2$, A_2 , $A - B^2 - B_1$, see fig. 41. Fig. 41 also shows q which is used with a' to form ξ_2 . The arc a' and hence the path ξ_2 turns out to contain p; hence $\xi_1 \cap \xi_2 \neq \emptyset$ as before \square .

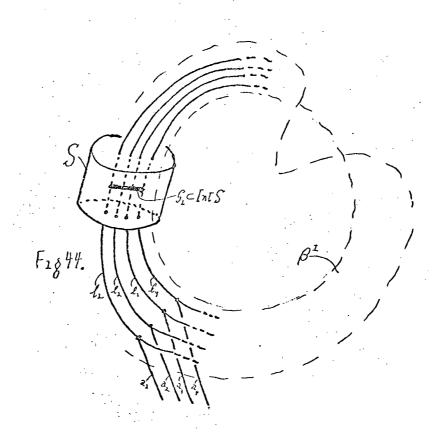
CHAPTER THREE: GENERALIZATION OF A THEOREM OF BING: LEMMAS.

1. In this chapter, we give two lemmas for the proof of II (2.2), the generalization of Bing's theorem II (2.1). In proving II (2.1), Bing defined a property Q such that A had Property Q, and if a dogbone $A_{jk} \ldots r$ had Property Q, then one of $A_{jk} \ldots r1$, $A_{jk} \ldots r2$, $A_{jk} \ldots r3$, $A_{jk} \ldots r4$ had Property Q. This meant that there was a chain $A \supset A_{j} \supset A_{jk} \supset \ldots$ of dogbones with Property Q. Since the possession of Property Q implied intersection with both disks D_{j} in II (2.1), the limit of the chain was a big element A which hit both D_{1} and D_{2} (see the discussion in IIS2). We follow Bing's proof closely (in spite of the fact that we alter Property Q to a property which has to be applied to a whole \mathcal{Q}_{m} to be of any use) and in fact depend on the reader's familiarity with [12] for the motivation in this chapter and the next. In the remainder of this paper, i = 1, 2, and j = 1, 2, 3, 4.

In the proof of II (2.1) in [12], it is evident that the crucial part of the argument is the proof of [12, Th 10], where it is shown that if the four centres of A_1 , A_2 , A_3 , A_4 fail to have Property P, then some set homotopic to the centre of A also fails to have Property P, (The precise definition of Property P is unimportant until Ch IV). In Ch IV we will prove just this result with the disks D_i in [1257] replaced by the singular disks $f_i[\Delta]$ in II (2.2). Our proof will differ from the proof of [12, Th 10] in that whereas in [12 Th 10] the disks D_i remain unchanged during the proof, in our proof of the analogous result the $f_i[\Delta]$ are replaced by new singular disks $f_i[\Delta]$ which retain the desirable properties of the $f_i[\Delta]$. Although

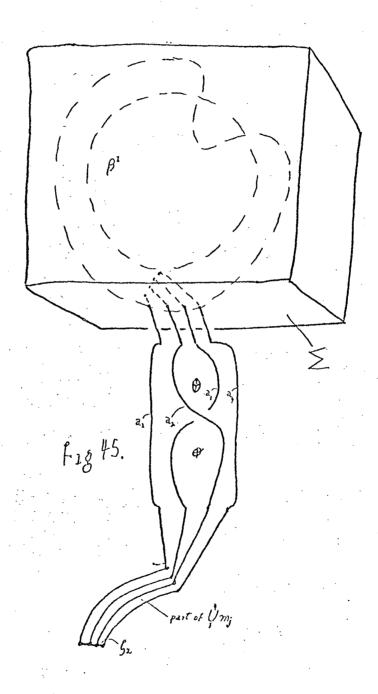






this is a considerable change, it turns out that our Ch IV resembles the argument of [12§7] very closely. In the present chapter, we prove an important lemma which shows that if each k_{\downarrow} (in fig 19) misses one of $f_1[\Delta]$, $f_2[\Delta]$, then the new $f_1[\Delta]$ may be constructed so that not only does each k_i miss one $f_i[\Delta]$, but both $f_i[\Delta]$ miss each of the arcs ζ_1 and ζ_2 shown in fig 42. The ζ_1 lie in β^1 and β_1 and tie the upper and lower loops of the k_{i} together as shown in the figure. If we can obtain such singular disks $f_i^*[\Delta]$, the reward is considerable, for then parts of the k, can be erased as shown in fig 43, leaving the set $b_1 \cup b_2 \cup b_3 \cup b_4 \cup c_1 \cup c_2$ shown in this figure. Since each $b_i \subset k_i$, each b_i misses one $f_i[\Delta]$ while $\zeta_1 \cup \zeta_2$ misses both. One can now apply Part II of the proof of Th 10 of [12] to b₁ \cup ... υ b₄ \cup ζ 1 \cup ζ 2 instead of to $U_{pq_jr_js}$ in [12, fig 2]. This can be done with very little change in the argument of [12] and results in the construction of a centre of A which fails to have Property P. We say that mappings $g_1: \Delta \to E^3$ are Z-disjoint iff $Z \subset E^3$ and rng $g_1 \cap rng g_2 \cap Z = \emptyset$, i.e. iff the ranges are disjoint at least in $\ensuremath{\text{Z}}$.

Lemma One. Consider A, A_j , β^1 , k_j as defined in Ch II (see fig 19). Let $Z\supset A$ and let $c_1\colon Bd\Delta\to E^3-Z$. Let $g_i\colon \Delta\to E^3$ be Z-disjoint mappings such that $g_i\mid Bd\Delta = c_i$. Let S be the sphere shown in fig 44 consisting of the cylindrical annulus Ω with disks d_1 , d_2 for end caps. Each ℓ_j pierces each d_i exactly once and β^1 misses Ω . Let $S\subset I$ Int A and let N be an η -neighbourhood of S such that $N\subset I$ Int A. The arc ζ_1 shown in fig 42 lies in Int A - N. Then there exist Z-disjoint mappings $g_i: \Delta\to E^3$ such that



i)
$$\frac{1}{g_i} = c_i$$
 on $Bd\Delta$,

ii)
$$\overline{g}_i: \Delta \rightarrow (\text{rng } g_i - \text{Int } S) \cup N$$
,

iii) If
$$\ell_j \cup \Omega$$
 misses rng g_i , then ℓ_j misses rng g_i .

Corollary. Let K_1 be the cube defined in Ch II (see fig 21). Then ii) and iii) in Lemma One may be replaced by

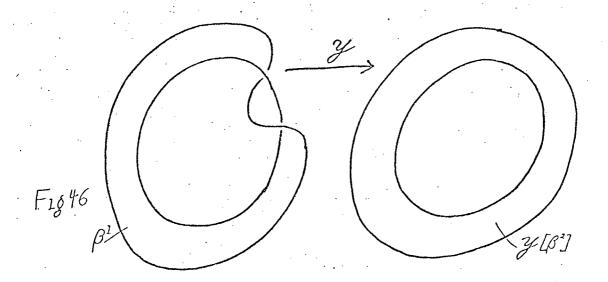
ii)
$$\frac{1}{g_i}: \Delta \rightarrow \text{rng } g_i \cup K_1$$
,

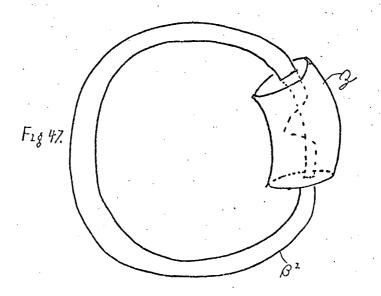
iii) If
$$k_j \cup \Omega$$
 misses rng g_i , then k_j misses rng g_i .

The proof of Lemma One is delayed to \$2, which may be read after Ch IV if desired.

We give a second lemma which is intended to repair a gap which would otherwise appear in the proof in Ch IV. This lemma is quite specialized, but appears here because its proof is just a variation of the proof of Lemma One. As before, the proof is delayed to §2 and may be omitted on a first reading.

Lemma Two. Consider A, Z, ζ_i , k_j , a_j , β^1 , as defined in fig 42. Let Σ be the sphere shown in fig 45. The sphere Σ together with an η - neighbourhood N of Σ lies in Int A; $\beta^1 \subset \text{Int } \Sigma$ - η , and each of a_1 , a_2 , a_3 pierces Σ as shown. Let mappings $g_i: \Delta \to E^3$ be Z-disjoint with $g_i|_{Bd\Delta} = c_i$, where c_i is defined as in Lemma One. Both rng g_i miss the set $\zeta_1 \cup k_1 \cup k_2 \cup k_3$. Let u_{12} , u_{13} be arcs in β^1 which join $a_1 \cap \beta^1$ and $a_2 \cap \beta^1$, $a_1 \cap \beta^1$ and $a_3 \cap \beta^1$ respectively and miss rng g_1 . Let v_{12} , v_{13} be arcs in β^1 which join $a_1 \cap \beta^1$ and $a_2 \cap \beta^1$, $a_2 \cap \beta^1$ and $a_3 \cap \beta^1$ respectively and miss rng g_2 . The arcs u_{12} , u_{13} , v_{12} , v_{13} are not necessarily





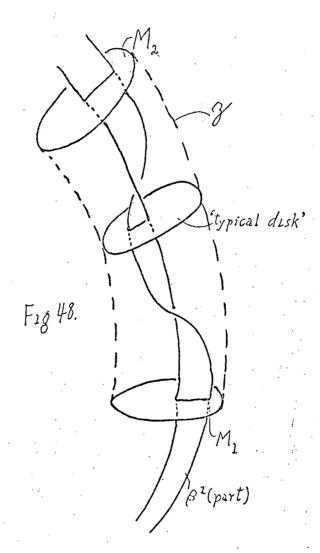
disjoint. Then there exist Z-disjoint mappings $g_i:\Delta \to (\operatorname{rng} g_i - \operatorname{Int}\Sigma) \cup \mathbb{N}$ such that $g_i' = c_i$ on $\operatorname{Bd}\Delta$, and one of $\operatorname{rng} g_1'$, $\operatorname{rng} g_2'$ misses $k_1 \cup k_2 \cup k_3$.

Corollary. One of rng g_1' , rng g_2' misses $\zeta_1 \cup b_1 \cup b_2 \cup b_3 \cup \zeta_2$, and $g_1': \Delta \rightarrow \text{rng } g_1' \cup K_1$.

Although u_{12} lies in an annulus and is joined to Σ by the orderly arcs $a_1 \cap \overline{\text{Int }\Sigma}$, $a_2 \cap \overline{\text{Int }\Sigma}$, the arc $(a_1 \cup u_{12} \cup a_2) \cap \overline{\text{Int }\Sigma}$ may be knotted in $\overline{\text{Int }\Sigma}$, as a few moments experiment will show (an arc ab in a cube K with $ab \subset Bd D = a \cup b$ is knotted if there is no disk $D \subset K$ with $ab \subset Bd D \subset Bd K U ab$). To be knotted the arc must make more than one circuit on the twisted annulus. A similar comment applies to u_{13} , v_{12} , v_{13} .

2. Proof of Lemma One.

(2.1). As a preliminary, we describe an untwisting function $y: \mathbb{E}^3 \to \mathbb{E}^3$ which is onto and one-to-one and which unwinds the twist in β^1 , i.e. $y(\beta^1)$ is the planar annulus shown in fig 46. For well known reasons y cannot be a mapping, but we ensure that y will be discontinuous only on the curved cylindrical surface z shown in fig 47. In fig 47, the end caps of z are called M_1 , M_2 , and the cube $\overline{\text{Int}(z \cup M_1 \cup M_2)}$ is called K. Eventually y will be composed with a mapping whose range misses z. Thus the result of the composition will be a mapping. The function y is defined to be the identity on \mathbb{E}^3 - K and on both M_1 . To define y in Int K: Imagine K to be cut free of the space by means of a cut on z and on M_2 , remaining attached only on M_1 .



K may be thought of as a stack of circular disks of infinitesmal thickness. These disks span the cylinder $\,z\,$ and each meets $\,\beta^{\,1}\,$ in a straight arc. Fig 48 shows M_1 , which is called the initial disk; M_{2} , which is called the final disk; and a 'typical disk' in the stack between M_1 and M_2 . Now apply a twist (which may be thought of as an isotopy of $\ensuremath{\mbox{K}}\xspace)$ to $\ensuremath{\mbox{M}_2}\xspace$ so that $\ensuremath{\mbox{M}_2}\xspace$ rotates once (i.e. through an angle of (2π) in place. When this happens, the disk \mathbf{M}_1 , which is attached to the space, necessarily remains fixed and does not rotate. Each disk intermediate between M_1 and M_2 rotates through an angle which is close to zero for disks close to $\,{\rm M}_{1}^{}\,\,$ and approaches $\,2\pi\,\,$ for disks whose location approaches that of $\,\mathrm{M}_2$. The rotations of the various disks in the stack can be contrived so that $\beta^1 \cap K$ is carried onto the plane which contains $\beta^1 - K$, and so that the final result is homeomorphic to K . In fig 48, the 'typical disk', which is located half-way between $\,{\rm M_1}\,\,$ and $\,{\rm M_2}\,\,$ will rotate through an angle of $\,\pi$. This carries its intersection with β^1 on to the desired plane. Since M_2 has returned to its original position, we can restore the cut at M_2 . Evidently y is one-to-one and continuous in Int K, Ext K, and on $\mathrm{M_1}\ \mathrm{U}\ \mathrm{M_2}$. The fact that we cannot sew up the cut on $\ \mathrm{Z}$ appears in the definition of y as a discontinuity on z . Clearly y carries β^{\perp} into the plane containing $\beta^1 - K$.

- (2.2). We will prove a simpler version of Lemma One to show the general approach.
- (2.21). Let S be a sphere in E^3 having a simply connected neighbourhood N . Let $g: \Delta \to E^3$ so that $g[Bd\Delta] \subset Ext S$. Then there

exists a mapping $g: \Delta \to (\text{rng g - Int S}) \cup \mathbb{N}$ which agrees with g on $Bd\Delta$.

When simplified in this way, (2.21) is insignificant, for there are easier proofs of stronger results, as the reader doubtless sees.

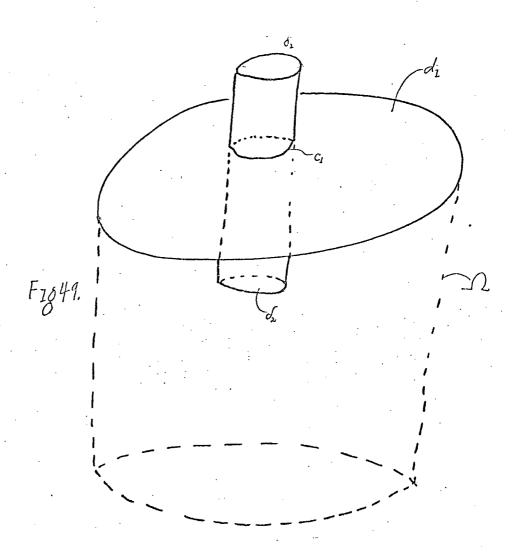
However our proof is intended to show how II(1.2) is used in the proof of Lemma One.

Proof. Apply II(1.2) to obtain disks $W_1, \ldots W_n$ in Δ such that $g^{-1}[\overline{\operatorname{Int S}}] \subset W_1 \cup \ldots \cup W_n$. Since $g^{-1}[\overline{\operatorname{Int S}}]$ misses $\operatorname{Bd}\Delta$, Δ - Int W_1 - ... - Int W_n is a disk with holes and $g^{-1}[\overline{\operatorname{Int S}}] \subset \operatorname{Int W}_1$... Int W_n . If ε in II(1.2) is sufficiently small, then g carries each $\operatorname{Bd}W_r$ into N, for $g[\operatorname{Bd}W_r]$ lies close to, but not in $\overline{\operatorname{Int S}}$ and hence close to S. In II(1.3), take (simply connected) N to be $P_1 = P_2 = \cdots = P_n$ to obtain the mapping $\overline{g} = g|\Delta - \operatorname{Int W}_1 - \cdots - \operatorname{Int W}_n \cup \gamma_1 \cup \gamma_2 \cup \cdots \cup \gamma_n : \Delta \to \operatorname{rng} g \cup N$. Since each point x in Δ lies either in Δ - $\operatorname{Int W}_1 - \cdots - \operatorname{Int W}_n$ in which case $\overline{g}(x) \in E^3 - \overline{\operatorname{Int S}}$, or in some W_r , in which case $\overline{g}(x) \in N$, $\operatorname{rng} \overline{g}$ misses $\operatorname{Int S} - N$. Thus $\operatorname{rng} g$ lies in $(\operatorname{rng} g - \operatorname{Int S}) \cup N$. Finally $\overline{g} = g$ on $\operatorname{Bd}\Delta$ because the two mappings differ only in $W_1 \cup W_2 \cup \cdots \cup W_n$, which misses $\operatorname{Bd}\Delta$.

(2.3). We will now give a formal proof of Lemma One.

Case one: neither $g_i[\Delta]$ meets S. Let $\overline{g}_i = g_i$. Since $\overline{g}_i[\Delta]$ meets Ext S, a connectivity argument shows that $\overline{g}_i[\Delta]$ misses Int S. The rest of the requirements of Lemma One are clear.

In the next two cases we insist that one of the $\mbox{ rng g}_{\mbox{\scriptsize i}}$ touch Ω while the other does not.



Case two: exactly one rng g_1 meets S. The rng g_1 which meets S also meets Ω . Assume that rng g_1 meets $\Omega \subset S$. Let $\overline{g}_2 = g_2$. Evidently i), ii), iii) of Lemma One are true of \overline{g}_2 . Apply the argument of 2.2, taking g in 2.2 to be g_1 , and construct a mapping $\overline{g}_1: \Delta \to (\text{rng } g_1 - \text{Int } S) \cup N$ which agrees with c_1 on $Bd\Delta$. With regard to \overline{g}_1 , i) and ii) are satisfied, and iii) is vacuously satisfied since $g_1[\Delta]$ hits Ω . The \overline{g}_1 are Z-disjoint because we could have taken N small enough to miss rng g_2 . Thus $\emptyset = Z \cap \text{rng } g_1 \cap \text{rng } g_2 = Z \cap (\text{rng } g_1 \cup N) \cap \text{rng } g_2 \supset Z \cap \text{rng } \overline{g}_1 \cap \text{rng } \overline{g}_2$.

Case three: both rng \mathbf{g}_i meet S . One rng \mathbf{g}_i , say rng \mathbf{g}_1 , meets Ω ; the other (rng \mathbf{g}_2) does not. The aim of the proof will be to construct an intermediate pair of mappings \mathbf{g}_i^k such that rng \mathbf{g}_2^k misses S although rng \mathbf{g}_1^k may not. The argument then reduces to an easy variation of either case one or case two.

Outline of proof. a) Choose a component z_1 of $\operatorname{rng} g_2 \cap S$. It is important that, since $\operatorname{rng} g_2$ misses Ω , $\operatorname{rng} g_2 \cap S \subset \operatorname{Int} d_1 \cup \operatorname{Int} d_2$. Using this fact, we construct a circle $c_1 \subset \operatorname{Int} d_1 \cup \operatorname{Int} d_2 - \operatorname{rng} g_1 - \operatorname{rng} g_2$ which encloses (on one of the d_1) points of exactly one of $\operatorname{rng} g_1 \cap S$, $\operatorname{rng} g_2 \cap S$. Although we choose $z_1 \subset \operatorname{rng} g_2 \cap S$, c_1 may turn out to enclose points of $\operatorname{rng} g_1 \cap S$.

b) Assume that $\overline{\operatorname{Int}} \ c_1 \subset \operatorname{Int} \ d_1$. Construct a sphere $\omega \cup \delta_1 \cup \delta_2$ in the shape of a pill-box as shown in fig 49 so that c_1 is the equator of $\omega \cup \delta_1 \cup \delta_2$.

c) An argument like that of case two but using ω υ δ_1 υ δ_2 instead of S yields a pair of Z-disjoint mappings

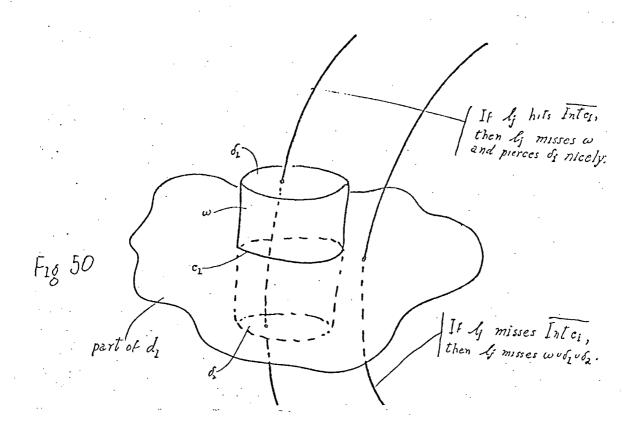
 $g_1^1: \Lambda \to \operatorname{rng} g_1 \cup N$ such that $g_1^1 = c_1$ on $\operatorname{Bd} \Lambda$, Int c_1 misses $\operatorname{rng} g_1^1 \cup \operatorname{rng} g_2^1$; and if $\operatorname{rng} g_1$ misses $\Omega \cup \ell_j$, then so does $\operatorname{rng} g_1^1$. The argument of case two is used virtually as is if c_1 encloses points of $\operatorname{rng} g_1 \cap S$. If c_1 encloses points of $\operatorname{rng} g_2 \cap S$ the argument must be modified somewhat, since the method of case two would not ordinarily ensure that $\operatorname{rng} g_2^1$ would miss all the ℓ_j that $\operatorname{rng} g_2$ misses.

d) It turns out that $\operatorname{rng} \operatorname{g}_2^1$ misses Ω . If some component $\operatorname{z}_2 \subset \operatorname{rng} \operatorname{g}_2^1 \cap \operatorname{S}$ exists in, say, Int d_1 , then we repeat steps a), b) c) to define mappings $\operatorname{g}_1^2: \Delta \to \operatorname{rng} \operatorname{g}_1^1 \cup \operatorname{N}$ with properties analogous to those of the g_1^1 . We continue to construct mappings g_1^3 , g_1^4 , ..., first finding a component $\operatorname{z}_{r+1} \subset \operatorname{rng} \operatorname{g}_2^r \wedge (\operatorname{Int} \operatorname{d}_1 \cup \operatorname{Int} \operatorname{d}_2)$ and then constructing the pair $\operatorname{g}_1^{r+1}: \Delta \to \operatorname{rng} \operatorname{g}_1^r \cup \operatorname{N}$ such that $\operatorname{g}_1^{r+1} = \operatorname{c}_1$ on $\operatorname{Bd}\Delta$ and if $\operatorname{rng} \operatorname{g}_1$ misses $\Omega \cup \ell_1$, then so do $\operatorname{rng} \operatorname{g}_1^1$, $\operatorname{rng} \operatorname{g}_1^r$, $\operatorname{rng} \operatorname{g}_1^r$, $\operatorname{rng} \operatorname{g}_1^r$, $\operatorname{rng} \operatorname{g}_1^r$. We show that $\operatorname{rng} \operatorname{g}_1^r \cap \operatorname{S} \supset \operatorname{rng} \operatorname{g}_1^{r+1} \cap \operatorname{S}$ and that the sequence of mappings ends at a pair g_1^k for which $\operatorname{rng} \operatorname{g}_2^k \cap \operatorname{S} = \emptyset$ although $\operatorname{rng} \operatorname{g}_1^k$ may hit S .

e) The situation now reduces to case one or case two. An additional argument shows that iii) of Lemma One is satisfied.

Details of proof. a) If we assume that $\operatorname{rng} \operatorname{g}_2$ misses Ω but hits S, then there is a component z_1 of $\operatorname{rng} \operatorname{g}_2 \wedge \operatorname{S}$ lying in Int d_1 or Int d_2 , say Int d_1 . By the Zoretti Theorem, there is a circle x_1 in Int d_1 which misses $\operatorname{rng} \operatorname{g}_2$, encloses z_1 , and lies so near to z_1 that it misses $\operatorname{rng} \operatorname{g}_1$ (by 'encloses' we mean 'encloses relative to d_1 '). If x_1 encloses points of just one of $\operatorname{rng} \operatorname{g}_1 \wedge \operatorname{S}$, $\operatorname{rng} \operatorname{g}_2 \wedge \operatorname{S}$,

then let χ_1 be c_1 . We must expect that χ_1 will enclose points of both rng $\mathbf{g}_1 \cap \mathbf{S}$ and rng $\mathbf{g}_2 \cap \mathbf{S}$ (\mathbf{z}_1 could be itself a circle enclosing points of $\operatorname{rng} \ \operatorname{g}_1 \ \operatorname{\sigma}$ S). In this case use the Plane Separation Theorem in Int $\chi_1 \subset \text{Int d}_1$ to construct a circle $\chi_2 \subset \text{Int } \chi_1$ which misses $(g_1[\Delta] \cup g_2[\Delta]) \cap S$ and separates the component z_1 of $g_2[\Delta] \cap S$ from a component z' of $\mathbf{g}_1[\boldsymbol{\Delta}] \circ \mathbf{S}$ in Int \mathbf{x}_1 . Since \mathbf{z}_1 may be in Int χ_1 - Int χ_2 , we cannot predict that Int χ_2 contains points of rng $g_2 \cap S$; but by the Plane Separation Theorem we know that χ_2 encloses points of (rng $\mathbf{g}_1 \ \boldsymbol{\upsilon}$ rng $\mathbf{g}_2) \ \boldsymbol{\upalpha}$ S . If χ_2 encloses points of both rng $g_1 \cap S$ and rng $g_2 \cap S$, then separate Int $\chi_2 \cap (\text{rng } g_1 \cup \text{rng } g_2)$ still further by means of another application of the Plane Separation Theorem. We repeat this procedure, defining circles χ_3 , χ_{L} , ...; χ_{r} being constructed in $% \left(1\right) =\left(1\right) \left(1\right) =\left(1\right) \left(1\right) \left($ both rng $g_i^{r-1} \wedge S$. The following argument shows that the sequence χ_1 , χ_2 , ... must be finite: each annulus $\overline{\text{Int }\chi_r}$ - Int χ_{r+1} contains points of (rng $\mathbf{g_1} \ \cup \ \mathbf{rng} \ \mathbf{g_2}$) $\boldsymbol{\alpha}$ S . Without loss of generality in the construction of the $\chi_{\mathbf{r}}$, we could have replaced the sets $\operatorname{rng} \ \mathbf{g}_{\mathbf{i}} \ \mathbf{n}$ S with 'thickened' sets obtained by covering the $\Re \operatorname{rng} \ g_i \cap S$ by small disks of area σ (from compactness, the thickened rng $\mathbf{g_1}$ o S can be assumed to remain disjoint from the thickened rrng \mathbf{g}_2 \mathbf{o} S) . But since each of the disjoint open annuli must therefore have area at least o, the number of $\,\chi_{\mbox{\scriptsize r}}\,\,$ must be finite and the sequence ends at some $\,\chi_{\mbox{\scriptsize t}}\,\,$. Since χ_{t+1} could be defined if χ_t encloses points of both rng g₁ σ S and rng $g_2 \cap S$, χ_t must enclose points of only one of rng $g_1 \cap S$, rng $\mathbf{g}_2 \, \boldsymbol{n}$ S . Let $\, \boldsymbol{\chi}_{t} \,$ be $\, \mathbf{c}_1 \,$. We repeat that we do not know which of rng $g_1 \cap S$, rng $g_2 \cap S$ is intersected by Int $\chi_t = Int c_1$.



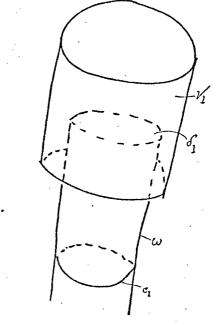


Fig 51.

b) Now assuming d_1 to be horizontal, we build a small sphere in the shape of a pill box consisting of vertical cylinder ω_{0} and end caps δ_{1}, δ_{2} , which are parallel to d_{1} (see fig 49). The cylinder $\,\omega\,$ intersects $\,d_1\,$ only at $\,c_1\,$ and extends equal distances above and below d_1 . Thus Int c_1 (considered as a subset of d_1) lies in Int($\omega \cup \delta_1 \cup \delta_2$). Since c_1 misses rng g_1 υ rng g_2 we build ω so near c_1 that ω misses rng g_1 υ rng g_2 . Fig 50 shows $\dot{\omega} \cup \delta_1 \cup \delta_2$ and part of d_1 . We assume that c_1 has been moved slightly if necessary so as to miss ($\ell_1 \ \lor \ \ell_2 \ \lor \ \ell_3 \ \lor \ \ell_4$) $\land \ d_1$. We also assume that ω misses every ℓ_{i} although this may necessitate making $\,\omega\,\,$ smaller or even curving $\,\omega\,\,$ slightly to follow the curve of ℓ_{i} . The sphere $\omega\,\upsilon\,\delta_{1}\,\upsilon\,\delta_{2}$ is constructed as in fig 50 so that if some ℓ_i meets d_1 then ℓ_i (misses ω and) pierces each Int δ_i just once. Let v_i be the simply connected neighbourhood of δ_i shown in fig 51. We construct v_i so that $\delta_i - \ell_1 - \ell_2 - \ell_3 - \ell_4$ is a deformation retract of v_i - l_1 - l_2 - l_3 - l_4 , and so that v_i misses S \supset Int c_1 and any ℓ_i which misses Int c_1 , i.e. any ℓ_i which misses ω ω δ_1 ω δ_2 . We assume that ω ω δ_1 ω δ_2 ω v_1 ω v_2 has been constructed so that ω υ δ_1 υ δ_2 υ υ υ 1 ies in N, and (since $c_1 \subset Int d_2$) so that $\overline{Int(\omega \cup \delta_1 \cup \delta_2)} \cup v_1 \cup v_2$ misses $\Omega \cup d_2$.

c) Assume for the moment that Int c_1 contains points of rng $g_1 \wedge S$. Because rng g_1 hits Ω , we can ignore iii) in Lemma One as far as \overline{g}_1 is concerned, i.e. the g_1^k which we are about to construct need not miss any ℓ_j . We assume that $\omega \cup \delta_1 \cup \delta_2 \cup \nu_1 \cup \nu_2$ lies so near Int c_1 that $\omega \cup \delta_1 \cup \delta_2 \cup \nu_1 \cup \nu_2$ misses rng g_2 . Since

 ω also misses rng \mathbf{g}_1 , rng \mathbf{g}_1 meets $\omega \ \mathbf{v} \ \delta_1 \ \mathbf{v} \ \delta_2$ only in δ_1 or δ_2 . Apply an argument like that of 2.2 to $\omega_1 \vee \delta_1 \vee \delta_2$, i.e. use II(1.2), taking S in II(1.2) to be $\omega \cup \delta_1 \cup \delta_2$ and N to be $v_1 \cup v_2$, to obtain disks $W_1, \dots W_n$ in Δ such that $g_1^{-1}[\overline{\text{Int}(\omega \cup \delta_1 \cup \delta_2)}]$ lies in Int $W_1 \cup \ldots \cup Int W_n$. Since $gg_1[Bd W_r]$ lies near rng $g_1 \cap (\omega \cup \delta_1 \cup \delta_2)$ as we saw in 2.2, and since ω misses rng g_i , therefore $g_1[Bd\ W_r]$ lies near $\delta_1\ U\ \delta_2$, i.e. in $v_1\ U\ v_2$. Evidently δ_1 each $g_1[Bd\ W_r]$ lies in one v_i . Since each $g_1[Bd\ W_r]$ lies in a simply connected subset of $v_1 \cup v_2$, we can construct the mappings γ_r in II(1.3) and, following the argument of 2.2, define a mapping $g_1^1: \Delta \rightarrow [\operatorname{rng} g_1 - \operatorname{Int}(\omega \cup \delta_1 \cup \delta_2)] \cup v_1 \cup v_2$ such that $g_1^1 = c_1$ on $\operatorname{Bd}\Delta$. Let $g_2 = g_2^1$. Then it is true of both g_i^1 that $g_i^1: \Delta \rightarrow \operatorname{rng} g_i \cup N$, $g_i^1 = c_1$ on BdA, and if rng g_i misses $\Omega \cup \ell_i$ then so does rng g_i^1 (this last property is only true because $\ {\rm rng}\ {\rm g}_1\ \ {\rm hits}\ \ \Omega$ $\ \ {\rm and}\ \ {\rm because}$ g_2 , which can miss some $\Omega \cup \ell_1$, is identical to g_2^1). Additionally Int c_1 on d_1 misses rng g_2 because $g_2 = g_2^1$ and misses rng g_1 because Int c_1 misses $v_1 \cup v_2$. The g_i^1 are Z-disjoint because rng $g_2^1 = \text{rng } g_2$, and rng g_1^1 exceeds rng g_1 only in $v_1 \cup v_2$ which misses rng \mathbf{g}_2^1 . (By 'rng \mathbf{g}_1^1 exceeds rng \mathbf{g}_1 only in N', we mean that rng $g_1 \cup N \supset rng g_1^{\perp}$.)

Suppose that instead of rng $\mathbf{g}_1 \cap \mathbf{S}$, Int \mathbf{c}_1 encloses points of rng $\mathbf{g}_2 \cap \mathbf{S}$. Let $\mathbf{g}_1 = \mathbf{g}_1^1$, and construct $\omega \vee \delta_1 \vee \delta_2 \vee \nu_1 \vee \nu_2$ as in b), this time so that $\omega \vee \delta_1 \vee \delta_2 \vee \nu_1 \vee \nu_2$ misses rng \mathbf{g}_1 and ω misses rng \mathbf{g}_2 . Constructing \mathbf{g}_2^1 is harder than constructing \mathbf{g}_1^1 as we did in the last paragraph for we must ensure that rng \mathbf{g}_2^1 misses any set $\Omega \vee \ell_j$ that rng \mathbf{g}_2 misses. We must take careful account of the various ℓ_j . Some ℓ_j are not missed by rng \mathbf{g}_2 and can be ignored.

Some ℓ_1 are missed by rng g_2 but do not meet $\overline{\text{Int } c_1}$; we note that ω v δ_1 v δ_2 v v v v has been constructed so that any ℓ_1 which misses Int c_1 also misses $\omega \cup \delta_1 \cup \delta_2 \cup v_1 \cup v_2$. With this precaution, it is safe to ignore those ℓ_1 which miss $\operatorname{rng} \mathsf{g}_2$ and also miss $\overline{\operatorname{Int}\, \operatorname{c}_1}$. In the remainder of this paragraph we will assume that ℓ_1 and ℓ_2 are those ℓ_1 which miss $\operatorname{rng} \, \operatorname{g}_2$ and hit $\overline{\operatorname{Int} \, \operatorname{c}_1}$. We think that the procedure in the general case that some subset ℓ_j , ℓ_{j_2} , ... ℓ_{j_s} of ℓ_1 , ℓ_2 , ℓ_3 , ℓ_4 misses rng g_2 and hits Int c_1 will be evident. We proceed to define g_2^1 using II(1.2) and II(1.3) as before. The only difficulty occurs when we wish to shrink $g_{2|Bd \ W_{m}}$ to a point so as to define γ_r . It was easy to shrink $g_1|_{Bd\ W_r}$ a point in one component, say v_1 , of v_1 v_2 in the course of defining $\mathbf{g}_{1}^{\mathbf{I}}$. But in this case we must shrink $\mathbf{g}_{2\,|\,\mathrm{Bd}\,\,\mathrm{W}_{r}}$ to a point in \mathbf{v}_{1} - ℓ_{1} - ℓ_{2} ; otherwise rng γ_{r} and hence rng g_{2}^{1} will hit $\ell_{1} \cup \ell_{2}$. The reason that $g_{2|Bd|W_{-}}$ will shrink to a point on $v_1 - \ell_1 - \ell_2$ is that $g_2[W_r]$ misses $\ell_1 \ \upsilon \ \ell_2$ (because rng \mathbf{g}_2 does) and can be assumed to miss z in fig 47 without loss of generality. There is a retract R (though π not a deformation retract) of E 3 - z - ℓ_1 - ℓ_2 onto δ_1 - ℓ_1 - ℓ_2 . Additionally it turns out that R restricted to v_1 - ℓ_1 - ℓ_2 is a deformation retract of v_1 - ℓ_1 - ℓ_2 onto δ_1 - ℓ_1 - ℓ_2 . This means that $Rg_{2|Bd W_r}$ is homotopic to $g_{2|Bd W_r}$ in v_1 - ℓ_1 - ℓ_2 ; and since $\text{Rg}_{2|\text{Bd W}_r}$ shrinks to a point in $\text{Rg}_2[\text{W}_r] \subset \text{V}_1 - \ell_1 - \ell_2$, therefore g 2|Bd w r shrinks to a point in v 1 - $^{\ell}$ 1 - $^{\ell}$ 2 as required. We delay the description of the retract R and the proof that $g_2[Bd\ W_r]$ can miss z until the end of this proof. Except for the use of $\ensuremath{\mathtt{R}}$ to make $g_{2|Bd W_r}$ shrink to a point, the construction of g_2^1 is like that of g_1^1 ,

and we have $g_2^1: \Delta \to [\operatorname{rng} g_2 - \operatorname{Int}(\omega \cup \delta_1 \cup \delta_2)] \cup (\nu_1 \cup \nu_2 - \ell_1 - \ell_2)$. The g_1^1 are Z-disjoint by an argument like that in the previous paragraph, and $g_1^1: \Delta \to \operatorname{rng} g_1 \cup N$, $g_1^1 = c_1$ on Bd Δ as before. We know that $\operatorname{rng} g_2^1$ misses Ω because $\operatorname{rng} g_2^1$ exceeds $\operatorname{rng} g_2$ only in $\nu_1 \cup \nu_2$ which is remote from Ω . If ℓ_j misses $\operatorname{rng} g_2$, then either ℓ_j meets $\operatorname{Int} c_1$, in which case ℓ_j misses $\operatorname{rng} g_2^1$ because ℓ_j is one of ℓ_1 , ℓ_2 above; or else ℓ_j misses $\operatorname{Int} c_1$, in which case ℓ_j misses $\operatorname{rng} g_2^1$ because ℓ_j is remote from $\nu_1 \cup \nu_2$.

Note that the g_i^1 satisfy the hypothesis of Lemma One. The important difference between g_i and g_i^1 is that Int $c_1 \cap \operatorname{rng} g_i^1 = \emptyset$ whereas Int c_1 hits one of $\operatorname{rng} g_1$, $\operatorname{rng} g_2$. Since $\operatorname{rng} g_i^1 \subset \operatorname{rng} g_i \cup v_1 \cup v_2$, $\operatorname{rng} g_i^1 \cap S \subset \operatorname{rng} g_i \cap S$ (because $v_1 \cup v_2$ misses S). Evidently we can write $(\operatorname{rng} g_1^1 \cup \operatorname{rng} g_2^1) \cap S \subset (\operatorname{rng} g_1 \cup \operatorname{rng} g_2) \cap S$ where the inclusion is proper.

d) Since the g_1^1 satisfy the hypothesis of Lemma One, we look for a component z_2 of $\operatorname{rng} g_2^1 \cap S$ in Int $d_1 \cup \operatorname{Int} d_2$ and repeat a), b), c) to obtain a circle $c_2 \subset \operatorname{Int} d_1 \cup \operatorname{Int} d_2$ and Z-disjoint mappings $g_1^2 : \Delta + \operatorname{rng} g_1^1 \cup N$ with $g_1^2 = c_1$ on $\operatorname{Bd}\Delta$, and if $\operatorname{rng} g_1$ misses $\Omega \cup \ell_1$, then $\operatorname{rng} g_1^1$ and $\operatorname{rng} g_2^2$ also miss $\Omega \cup \ell_1$. Furthermore $(\operatorname{rng} g_1^2 \cup \operatorname{rng} g_2^2) \cap S \subset (\operatorname{rng} g_1 \cup \operatorname{rng} g_2) \cap S \subset (\operatorname{rng} g_1 \cup \operatorname{rng} g_2) \cap S$ both inclusions being proper. We can continue in this way, defining mappings g_1^3 , g_1^4 , ... and components $z_3 \subset \operatorname{rng} g_2^2 \cap (\operatorname{Int} d_1 \cup \operatorname{Int} d_2)$, $z_4 \subset \operatorname{rng} g_2^3 \cap (\operatorname{Int} d_1 \cup \operatorname{Int} d_2)$, ... so that $g_1^r : \Delta + \operatorname{rng} g_1^{r-1} \cup N$, $g_1^r = c_1$ on $\operatorname{Bd}\Delta$, and if $\operatorname{rng} g_1$ misses $\Omega \cup \ell_1$, then so does $\operatorname{rng} g_1^r$.

Furthermore (rng $g_1^r \cup \text{rng } g_2^r$) \cap S \subset (rng $g_1^{r-1} \cup \text{rng } g_2^{r-1}$) \cap S where the inclusion is proper. An argument from compactness like that used in a) to show that the number of χ_r was finite can be used to show that there must be a final pair of Z-disjoint mappings g_1^k . Since g_1^{k+1} could be defined if z_{k+1} existed in rng $g_2^k \cap$ S, therefore rng g_2^k must miss S.

Since rng $g_1^r \subset rng \ g_1^{r-1} \cup N$, evidently $g_1^k : \Delta \to rng \ g_1^{k-1} \cup N \subset rng \ g_1^{k-2} \cup N \subset \ldots \subset rng \ g_1 \cup N$; and $g_1^k = c_1$ on Bd Δ , while if $\Omega \cup \ell_j$ misses some rng g_1 , then $\Omega \cup \ell_j$ misses rng g_1^k . Since rng g_2^k misses S, the argument reduces to either Case one or Case two. In the course of the argument of Case one or two, g_2^k , whose range already misses S, will be set equal to \overline{g}_2 . This means that \overline{g}_2 has the properties of g_2^k ; thus i), ii), iii) of Lemma One are true for \overline{g}_2 . The argument of either Case one or of Case two will now construct a new $\overline{g}_1 : \Delta \to rng \ g_1^k \cup N \subset rng \ g_1 \cup N$ with $\overline{g}_1 = c_1$ on Bd Δ , so that \overline{g}_1 , \overline{g}_2 are Z-disjoint. This proves i), ii) of Lemma One for \overline{g}_1 , while iii) is vacuously true by the Case three assumption.

Case four: Exactly one rng $\mathbf{g_i}$ meets S but misses Ω . The reader will find that the method of case three works here almost word or word if it is assumed that the rng $\mathbf{g_i}$ which hits S is rng $\mathbf{g_2}$. When we arrive at the point in Case three where $\mathbf{g_i^k}$ is defined, we can let $\mathbf{g_i^k}$ be $\mathbf{g_i}$ immediately (or go to Case one). Actually the retract R works on S in Case two just as well as on $\omega \cup \delta_1 \cup \delta_2$ in Case three. Thus a quick proof is possible by adapting Case two.

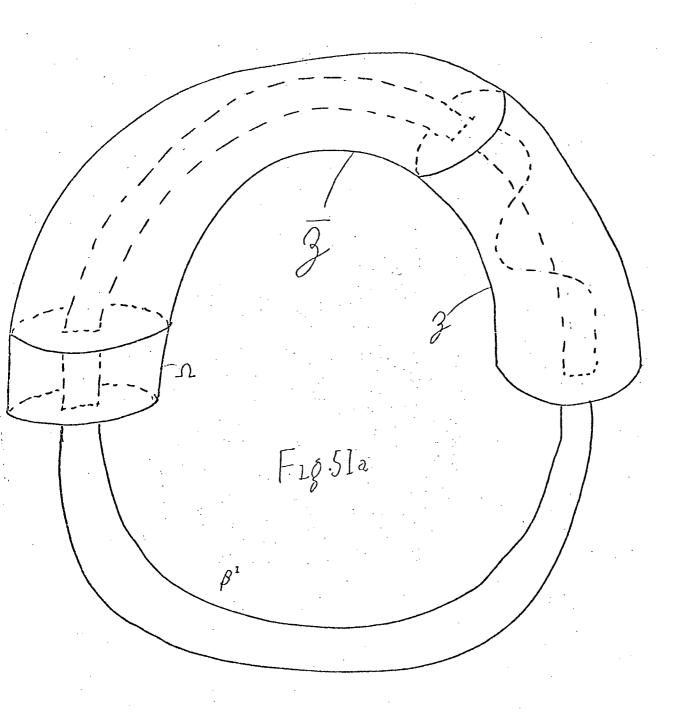
Case five: Both rng $\mathbf{g}_{\mathbf{i}}$ hit S; both rng $\mathbf{g}_{\hat{\mathbf{i}}}$ mîss Ω .

Proceed as in Case three to the point where the g_i^k are defined, allowing for the fact that iii) in Lemma One applies to both \overline{g}_i rather than only to \overline{g}_2 as in Case three (thus one may have to use the retract R to construct both g_i^r , whereas in Case three, R was used only to construct g_2^r). When the g_i^k are defined, the argument reduces to Case four or Case one.

Case six: Both rng g_i hit S; both rng g_i hit Ω . This case is not used in the applications of Lemma One, which always require every set $\Omega \cup \ell_j$ to miss one rng g_i . It is not hard to prove Case six using the ideas of the other cases.

(2.4). The retract R.

This retract was used in 2.3 Case three c). We will show how to define R: $E^3-z-\ell_1-\ell_2 \to \delta_1-\ell_1-\ell_2$; the definition of R when δ_1 is replaced by δ_2 is similar. Strictly speaking, the proof of Lemma One requires a retraction onto $d_1-\ell_{j_1}-\ldots-\ell_{j_s}$, where $\ell_{j_1},\ldots,\ell_{j_s}$ are some subset of $\ell_1,\ \ell_2,\ \ell_3,\ \ell_4$; however we continue the assumption in 2.3 Case three c) that rng g_2 misses $\ell_1\cup\ell_2$. Assume that the unique boundary component of β^1 which is a planar circle lies on the Y-Z plane and that the centre of this circle is the origin. The idea is that if we untwist β^1 by means of g_1 , all the circles $g[\ell_j]$ will be nice circles on the Y-Z plane with centre the origin. We assume further that δ_1 lies on the left-hand X-Y half-plane. We describe R in terms of several mappings which are applied in sequence to $E^3-z-\ell_1-\ell_2$. Each mapping will leave



 δ_1 - ℓ_1 - ℓ_2 fixed, and the last will be onto δ_1 - ℓ_1 - ℓ_2 .

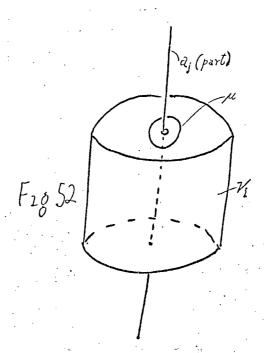
First: untwist β^1 by applying the mapping $y_{|E}^3 - z - \ell_1 - \ell_2$ (y) becomes a mapping by restricting it so that the domain misses the 'bad' set z). Each $y[\ell_1]$ is a plane circle with centre the origin. Second: using the symmetry of $E^3 - y(\ell_1 \cup \ell_2)$ across the X - Z plane, reflect the right-hand half-space minus $y(\ell_1 \cup \ell_2)$ onto the left-hand half-space minus $y(\ell_1 \cup \ell_2)$. This reflection carries $\mathbf{E}^3 - \mathbf{z} - y(\ell_1 \cup \ell_2)$ into those points in E^3 with non-positive coordinates which do not lie on $\;\ell_1\; {\it v}\; \ell_2\;$. Third: retract the left-hand half-space minus $y(\ell_1 \ v \ \ell_2)$ (which is the same as the left-hand half-space minus $\ \ell_1 \ v \ \ell_2$) onto the left-hand X - Y half-plane minus $\ensuremath{\ell_1}\ensuremath{\ensuremath{\ell_2}}\ensuremath{\ensuremath{\ell_2}}$. This is easy because the remaining parts of $\ \ell_1^{},\ \ell_2^{}$ are nice semicircles with centre the origin; one could imagine the X - Z plane hinged along the X axis. Using this hinge, topple the upper half of the X - Z plane onto the left-hand X - Y plane; simultaneously bring the lower half of the X - Z plane up to meet the $1\tilde{\mathbf{e}}\mathsf{f}\mathsf{t}\mathsf{-hand}$ X - Y plane. These movements define a (deformation) retract which crushes the left-hand half space minus $\ell_1 \cup \ell_2$ onto the left-hand X - Y half-plane minus $\ell_1 \cup \ell_2$. Finally retract the left-hand X - Y half-plane minus $\ensuremath{\ell_1} \ensuremath{\ensuremath{\upsilon}} \ensuremath{\ensuremath{\ell_2}}$ onto δ_1 - ℓ_1 - ℓ_2 . The four successive mappings define R . Note that R acts on $\,\nu_{1}^{}\,\,$ as a deformation retract (this was used in §2.3 Case three c)).

Finally we will show that $g_2[W_r]$ in 2.3 Case three c) can be assumed to miss the curved cylinder z. Since $g_2[\Delta]$ misses Ω , by the case three assumption, we can construct a mapping $\hat{g}_2:\Delta \to E^3$ whose range misses the curved cylinder \overline{z} in fig 51a and which agrees

with \mathbf{g}_2 on every point of Δ which maps under \mathbf{g}_2 outside of a small neighbourhood of \overline{z} . In fig 51a, \overline{z} is constructed so as to contain $\Omega \vee \overline{z}$ and to miss $\ell_1 \cup \ell_2$. Thus rng $\hat{\mathbf{g}}_2$ can be assumed to miss $\ell_1 \cup \ell_2$. The sets \mathbf{v}_1 , \mathbf{v}_2 (see 2.3 Case Three b)) miss Ω and could have been constructed so as to miss all of \overline{z} . Hence we assume that $\hat{\mathbf{g}}_2 = \mathbf{g}_2$ on Bd \mathbf{W}_r since $\mathbf{g}_2[\mathrm{Bd}\ \mathbf{W}_r] \subset \mathbf{v}_1$ by the definition of \mathbf{W}_r . We do not intend that $\hat{\mathbf{g}}_2$ should replace \mathbf{g}_2 since \mathbf{g}_1 , $\hat{\mathbf{g}}_2$ may not be Z-disjoint; but if $\mathbf{g}_2[\mathbf{W}_r]$ hits z, we can apply the retract R not to $\mathbf{g}_2|_{\mathbf{W}_r}$ but to $\hat{\mathbf{g}}_2|_{\mathbf{W}_r}$ and use the fact that $\mathbf{g}_2|_{\mathbf{Bd}\ \mathbf{W}_r}$ shrinks to a point in \mathbf{v}_1 - ℓ_1 - ℓ_2 iff $\hat{\mathbf{g}}_2|_{\mathbf{Bd}\ \mathbf{W}_r}$ does. Essentially the same argument applies to the construction of the other mappings in the sequence \mathbf{g}_2 , $\mathbf{g}_2^{(1)}$, $\mathbf{g}_2^{(2)}$, ... \square .

Proof of Lemma Two.

We will modify the argument of §2 so as to serve as a proof of Lemma Two. We assume familiarity with §2 in what follows. Modifying the argument of §2 to fit fig 45 presents a small and a large difficulty. The small problem is that we cannot build pillboxes according to the nice picture in fig 49, where c_1 is planar and the δ_1 can be considered to be horizontal while ω is vertical. It will be appreciated that the problem is more apparent than real; we have room to construct $\Sigma \subset \operatorname{Int} A$ with some obvious smoothness conditions so that if c is any circle on Σ and $\operatorname{Int} c$ is defined, then a sphere $\omega \cup \delta_1 \cup \delta_2$ can be constructed together with neighbourhoods v_1 , v_2 so that $\omega \cup \delta_1 \cup \delta_2 \cup v_1 \cup v_2$ behaves like the corresponding set in §2, i.e. $\omega \cap \Sigma = c$, δ_1 , δ_2 are disks in $\Sigma = \Sigma$ which meet ω only at its



two boundary components, while the ν_i are simply connected neighbourhoods of the δ_i which miss Σ . Furthermore, if a_j hits $\overline{\text{Int }c}$, then a_j pierces each Int δ_i just once and misses ω ; while if a_j misses $\overline{\text{Int }c}$, then a_j misses $\overline{\text{Int }(\omega \ \upsilon \ \delta_1 \ \upsilon \ \delta_2)} \ \upsilon \ \nu_1 \ \upsilon \ \nu_2$. We require that when Int c and hence δ_i hits just one a_j , then ν_i - a_j is homeomorphic to the structure shown in fig 52. This requirement is easy to manage; for Σ can be made to meet the a_j near β^1 , where (according to the definition on Ch II) the a_j are straight and parallel and perpendicular to β^1 (in fact it is easy to make $a_j \cap N$ a straight arc perpendicular to Σ).

The hard problem is that in Lemma Two we cannot use the retract R, which was crucial to the proof of iii) of Lemma One. The reason is that we permit the arcs $(a_1 \lor u_{12} \lor a_2) \cap \overline{\text{Int }\Sigma}$, etc. to be knotted, and in general deviate from the specialized geometry in fig 19. Recall that R was used to show that certain mappings $g_{\mathbf{i} \mid \text{Bd } W_{\mathbf{r}}}$ will shrink to a point in $v_{\mathbf{i}} - \ell_{1} - \ell_{2} - \ell_{3} - \ell_{4}$. Instead of R, we use the following easy but very weak result.

(3.1). Let ν_i be the usual neighbourhood of δ_i . Let ν_i intersect only a_1 as shown in fig 52. Let $f: Bd\Delta \to \nu_i - a_1$. Let $F: \Delta \to E^3$ so that f = F on $Bd\Delta$ and rng F misses a simple closed curve L such that $L \supset a_1$. The curve L may be knotted. Then f shrinks to a point in $\nu_i - a_1$. A similar result is true if a_2 or a_3 replaces a_1 .

Proof. Let μ be the small circle shown in fig 52. Then μ can be considered to represent the sole generator y of $\nu_{\hat{1}}$ - $a_{\hat{1}}$,

and also (by consulting, say, the definition in [6 Ch VI]) a generator of the Wirtinger presentation of E^3-L (we specify the particular presentation only to be sure that μ does not represent a trivial generator). If $i:\pi_1(\nu_1-a_1)\to\pi_1(E^3-L)$ is the inclusion homomorphism, then, with a change of basepoint, $f\in y^m$ for some integer m, and i(y) is an element of $\pi_1(E^3-L)$. Then $f\in i(y^m)=(i(y))^m$, which is the identity of $\pi_1(E^3-L)$ because f shrinks to a point in E^3-L . Since i(y) is a non-trivial element (in fact a Wirtinger generator) of $\pi_1(E^3-L)$, either m=0 or i(y) is an element of finite order. It is known ([7, (31.9)]) that the fundamental group of the complement of a knot has no element of finite order; therefore m-0, and f represents the identity g^0 in $\pi_1(\nu_1-a_1)$ \square .

 $a_{\hat{3}}$ say $a_{\hat{1}}$, then we use (3.3) instead of the retract R to shrink the various mappings $~{\rm g}_{\rm i\,|\,Bd~W_{\rm r}}~$ to a point in $~\nu_1~{\it U}~\nu_2~$ as was done in §2, Case three. Thus if rng g_i hits Int c_1 , rng g_i^1 misses a_1 because Int c_1 meets only one of a_1 , a_2 , a_3 ; while if rng g_i misses Int c₁, rng g_i misses a₁ because g_i = g_i. And by the usual arguments, a_2 and a_3 , which are remote from $\omega \cup \delta_1 \cup \delta_2$, continue to miss both rng g_i^1 . In applying (3.1), we let L be u_{12} or v_{12} (assuming that Int c_1 hits a_1) depending as rng \mathbf{g}_1 or rng \mathbf{g}_2 hits Int \mathbf{c}_1 . If Int \mathbf{c}_1 hits \mathbf{a}_2 only, let L be u_{12} or v_{12} again; if Int c_1 hits a_3 only, let L be u_{13} or v_{13} . Unfortunately, as the reader doubtless sees, if c_1 encloses more than one of $a_1 \cap \Sigma$, $a_2 \cap \Sigma$, $a_3 \cap \Sigma$, then the present argument fails (because the argument with (3.1) is weaker than the original argument in §2 which used the retract R), and g_{i}^{1} cannot be constructed so that rng g_i^1 misses all of a_1 , a_2 , a_3 . The trick of proving Lemma Two is to apply the argument of \$2 so that none of $\mbox{Int } c_1$, Int c_2 , ... ever hits more than one of a_1 , a_2 , a_3 . By extending the above ideas to further pairs g_i^2 , g_i^3 , ... and using methods from §2, we can prove

(3.2). In the context of Lemma Two, let D be a disk such that $D \subset \Sigma$, and let Bd D miss rng $g_1 \cup rng g_2 \cup a_1 \cup a_2 \cup a_3$. Then there exist circles c_1 , c_2 , ... c_m in Int D and Z-disjoint mappings g_1^1 , g_1^2 , ... g_1^m such that $g_1^r: \Delta \to rng g_1^{r-1} \cup (n-\Sigma)$, $g_1^1 = g_1^2 = \ldots = g_1^m = c_1$ on Bd Δ , c_r encloses (relative to D) points of just one of rng $g_1^{r-1} \cap \Sigma$, rng $g_2^{r-1} \cap \Sigma$, and rng g_1^m misses D. If, additionally, each c_r can be constructed so that c_r encloses just one of $a_1 \cap \Sigma$, $a_2 \cap \Sigma$,

 $a_3 \wedge \Sigma$, then rng g_1^m can be constructed so as to miss $a_1 \cup a_2 \cup a_3$.

Proof. (3.2) is proved in the same way as Lemma One. We can ignore Cases two and three in the proof of Lemma One because the fact that Bd D misses rng $\mathbf{g}_1 \cup \mathbf{rng} \ \mathbf{g}_2$ evidently takes the place of the condition in Lemma One that Ω misses rng $\mathbf{g}_1 \cup \mathbf{rng} \ \mathbf{g}_2$. Clearly we cannot have rng \mathbf{g}_1^m miss Int Σ in this version of the argument because D is a proper subset of Σ . The only part of the proof which does not have an exact counterpart in §2 is the statement that rng $\mathbf{g}_1^r \subset \mathbf{rng} \ \mathbf{g}_1^{r-1} \cup (n-\Sigma)$. The reason that rng $\mathbf{g}_1^r = \mathbf{g}_1^r \subset \mathbf{rng} \ \mathbf{g}_1^{r-1} \cup (n-\Sigma)$ is a usual, that $\mathbf{v}_1 \cup \mathbf{v}_2$ is remote from Σ .

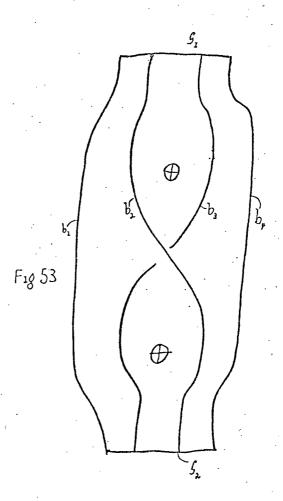
(3.3). Corollary. If, additionally to the hypothesis of (3.2), D misses rng g_1 , then rng g_1^m misses $a_1 \cup a_2 \cup a_3$ regardless of the number of a_1 , a_2 , a_3 which are hit by the Int c_r . Similarly, if rng g_2 misses D, then rng g_2^m misses $a_1 \cup a_2 \cup a_3$.

Proof. According to the argument of §2, if rng g_i misses D, then we let g_i = g_i^m immediately.

We will now give the proof of Lemma Two. The following question does not look like a simplification at first glance: Do there exist Z-disjoint mappings $f_i: \Delta \to (\operatorname{rng} g_i \cup n) - a_1 - a_2 - a_3$ with $f_i = c_i$ on $\operatorname{Bd}\Delta$, and a decomposition of Σ into disks D_1 , D_2 so that $\operatorname{D}_1 \cup \operatorname{D}_2 = \Sigma$ and $\operatorname{D}_1 \cap \operatorname{D}_2 = \operatorname{Bd} \operatorname{D}_1 = \operatorname{Bd} \operatorname{D}_2$, and so that Int D_2 misses one $\operatorname{rng} f_i$, say $\operatorname{rng} f_1$, and hits a_2 and a_3 ; while Int D_1 hits a_1 ?

Case one: the mappings f_i exist as described above. Look at the decomposition $\Sigma = D_1 \cup D_2$. Apply (3.2) to D_1 to convert the f_i to Z-disjoint mappings $\overline{f}_i: \Delta \to \operatorname{rng} f_i \cup (n-\Sigma)$ with $\overline{f}_i = c_i$ on Bd Δ and such that $\operatorname{rng} \overline{f}_i$ misses a_1 and D_1 . In (3.2), the condition that Int c_r hit at most one of a_1 , a_2 , a_3 is satisfied because a_2 , a_3 , miss $D_1 \supset \operatorname{Int} c_r$. Furthermore $\operatorname{rng} \overline{f}_i$ misses a_2 and a_3 by the usual argument. Now apply (3.3) to D_2 , to replace the \overline{f}_i with Z-disjoint mappings $g_i:\Delta \to \operatorname{rng} \overline{f}_i \cup (n-\Sigma)$ with $g_i|_{Bd\Delta} = c_i$. We know that $\operatorname{rng} \overline{f}_1$ misses D_2 since $\operatorname{rng} f_1$ does, and \overline{f}_i evidently satisfies the hypothesis of (3.2) and (3.3). By (3.3), $\operatorname{rng} g_i$ misses $a_1 \cup a_2 \cup a_3$, although $\operatorname{rng} g_2$ probably does not. Since $\operatorname{rng} g_i$ now misses all of Σ , $\operatorname{rng} g_i$ misses $\operatorname{Int} \Sigma$ except perhaps in n.

Case two: no mappings f_i exist as described. Let $d \in \Sigma$ be a disk pierced by a_1 which is small enough to miss both rng g_i and $a_2 \cup a_3$. Let $D = \overline{\Sigma - d}$. Using (3.2), construct a sequence of circles and mappings c_1 , g_1^1 , c_2 , g_1^2 , c_3 , g_1^3 , ... as described in (3.2), ending in the construction of Z-disjoint mappings $g_1^m = g_1' : \Delta \to \operatorname{rng} g_1 \cup (n - \Sigma)$ with $g_1' = c_1$ on $\operatorname{Bd}\Delta$ and such that $\operatorname{rng} g_1'$ misses D and $a_2 \cup a_3$. We know that every c_r encloses at most one of $a_1 \cap \Sigma$, $a_2 \cap \Sigma$, $a_3 \cap \Sigma$, as required by (3.2): for otherwise g_1^{r-1} , $\overline{\operatorname{Int} c_r}$, Σ - $\overline{\operatorname{Int} c_r}$ satisfy the definition of f_1 , D_1 , D_2 given above, which means that (since g^{r-1} exists) c_r contradicts the Case two assumption. Evidently $\operatorname{rng} g_1'$ misses not only $D \cup a_2 \cup a_3$, but also $d \cup a_1$, so that $\operatorname{rng} g_1'$ misses $a_1 \cup a_2 \cup a_3$ and all of Σ , etc $\overline{\Omega}$.



CHAPTER FOUR. GENERALIZATION OF A THEOREM OF BING: MAIN PROOF.

1. We will use Lemma One, Lemma Two and I§5 to prove II(2.2). The organization of the proof is much like that of [12 §7] and we depend on the reader's familiarity with [12] for orientation (although a detailed reading is required only of the section called 'Part II of Proof' in [12 §7]). As in [12 §7], we first give a (somewhat altered) definition of Property Q, then induce Property Q through the steps of the dogbone construction. This argument occupies most of the length of this chapter. As in [12 §7], it follows immediately (and for more or less the same reasons) that some big element of the decomposition hits both singular disks $f_{\cdot}[\Delta]$ in II(2.2).

Still following [12], we will not present a formal induction, but will show that if A has Property Q, then so does $A_1 \cup A_2 \cup A_3 \cup A_4 = a_1$ (Bing proves that one A_j has Property Q; our version of Property Q is only useful when applied to a_1 , a_2 , a_3 , ..., of Ex 2 in §2). The proof of this is divided into Part I and Part II as in [12 §7]. In Part I, we look at the set $\zeta_1 \cup b_1 \cup b_2 \cup b_3 \cup b_4 \cup \zeta_2$ (see fig 43) which serves a purpose like that of the set $\bigcup_j \operatorname{pq}_j r_j s$ in [12, fig 2]. We show that the f_1 in II(2.2) can be replaced by mappings pq_1 such that each $\operatorname{rng} \operatorname{q}_1$ misses one b_j and both ζ_1 . We call the set $\zeta_1 \cup b_1 \cup b_2 \cup b_3 \cup b_4 \cup \zeta_2$ the cradle of A_j , and later represent it as in fig 53, which preserves the embedding of $\zeta_1 \cup b_1 \cup b_2 \cup b_3 \cup b_4 \cup \zeta_2$ in A_j . In [12], $\bigcup_j \operatorname{pq}_j r_j s$ behaves like the cradle of A_j in that each $\operatorname{pq}_j r_j s$ misses one of the disks D_j in II(2.1). In Part II of our proof we follow [12] very closely and require a detailed reading

of the corresponding part of [12, Th 10]. There are a few alterations; these are required by the fact that some homotopies are replaced by isotopies.

2. Properties P and Q.

We will define a Property P on double ended lassos $\ell \cup a \cup m$ with respect to closed sets Y_1 , Y_2 . The lasso $\ell \cup a \cup m$ consists of circles ℓ and m connected by an arc a. In Ch II we often specified constructions only up to homotopy (e.g. the intersecting principal paths of Ch II). The consequence was that we ignored singularities in these constructions. In this chapter, this practice is emphatically not allowed; in particular, in the lasso $\ell \cup a \cup m$, the circles ℓ and m are disjoint simple closed curves and a meets $\ell \cup m$ only at its end points. One of the things that make the present chapter harder than Ch II is that geometric constructions have to be moved isotopically, whereas in Ch II homotopy was good enough.

Properties P and Q are defined in terms of their negatives, which we write Property $^{\circ}P$ and Property $^{\circ}Q$. A double ended lasso ℓ U a V m has Property $^{\circ}P$ with respect to closed sets Y_1 , Y_2 iff one of the following two conditions obtains.

We intend that Property $^{P}(b)$ should be symmetric, i.e. a \cup ℓ may miss $Y_1 \cup Y_2$ and the point y may be in m-a. Regardless of whether $\ell \cup a \cup m$ has Property $^{P}(a)$ or Property $^{P}(b)$, each of ℓ and m has Property $^{P}(a)$ as defined in I §5 for circles with base point (the base points here are taken to be $\ell \cap a$, $m \cap a$). This statement, which is important, is easily checked. Evidently $\ell \cup a \cup m$ may have both Property $^{P}(a)$ and Property $^{P}(b)$.

Property $^{\circ}P$ is the negative of Bing's Property P in [12]. It is easy to see that our Property $^{\circ}P$ implies the negative of Bing's property, i.e. our Property $^{\circ}P$ implies that if $\mathbf{x}_1 \in \ell$ and $\mathbf{x}_2 \in \mathbf{m}$ and $\ell \cup \mathbf{a} \cup \mathbf{m}$ has Property $^{\circ}P$ (by our definition), then there is an arc in $\ell \cup \mathbf{a} \cup \mathbf{m}$ with end points \mathbf{x}_1 , \mathbf{x}_2 which misses one of \mathbf{Y}_1 , \mathbf{Y}_2 . We will neither use nor prove the complete equivalence of the two definitions here, although a proof will be found to be straightforward.

Property $^{Q}_{Z, c_{1}, c_{2}}$ is defined on dogbones. If a dogbone X has Property $^{Q}_{Z, c_{1}, c_{2}}$, this means roughly that the centre of X has Property $^{Q}_{Z, c_{1}, c_{2}}$, this means roughly that the centre of X has Property $^{Q}_{Z, c_{1}, c_{2}}$ and for $i = 1, 2, c_{1}$: $^{B}_{A} + ^{B}_{A} - ^{B}_{A} - ^{B}_{A}$. Then X has Property $^{Q}_{Z, c_{1}, c_{2}}$ iff there exist Z-disjoint mappings $^{G}_{1}$, $^{G}_{2}$ such that $^{G}_{1}$: $^{A}_{1} + ^{B}_{1}$, $^{G}_{2}$ on $^{B}_{A}$, and the centre of X has Property $^{Q}_{Z, c_{1}, c_{2}}$ with respect to $^{G}_{1}$, $^{G}_{2}$ with the obvious meaning. We define X to have Property $^{Q}_{Z, c_{1}, c_{2}}$ iff X fails to have Property $^{Q}_{Z, c_{1}, c_{2}}$ (i.e. with respect to every qualified pair of mappings $^{G}_{1}$). Note that a statement like 'X has Property $^{Q}_{Z, c_{1}, c_{2}}$

with respect to g_1 , g_2 ' means very little.

Example 1). Suppose Z=X=A and c_1 , c_2 are the two circles shown in fig 28. Then A has Property Q_Z , c_1 , c_2 . For if c_1 (c_2) shrinks to a point, it must hit the upper (lower) eye ℓ (m) of A. Thus if f_i is an extension of c_i to all of Δ , then $\operatorname{rng} f_1$ hits ℓ and $\operatorname{rng} f_2$ hits m. This 'kills' Property P for $k=\ell \cup a \cup m$ with respect to $\operatorname{rng} f_1$, $\operatorname{rng} f_2$, since Property P would require either that one $f_i[\Delta]$ miss both ℓ and m or that one of ℓ or m miss both $f_i[\Delta]$.

Example 2). Let $Z = A_1$; c_1 , c_2 as in Ex. 1). Then A_1 has Property ${}^{^{1}\!\!Q}_{Z,c_1}$, c_2 ; for the c_1 can shrink to a point so as to miss Z and A_1 . We emphasize that 'X has Property ${}^{^{1}\!\!Q}_{Z,c_1}$, c_2 does not imply that c_1 , c_2 link the eyes of X.

Evidently if X has Property Q_Z , c_1 , c_2 and f_1 , f_2 are any Z-disjoint mappings of Δ into E^3 with $f_1 = c_1$ on Bd Δ , then the centre of X fails to have Property P(a) with respect to $f_1[\Delta]$, $f_2[\Delta]$, and consequently both $f_1[\Delta]$ and $f_2[\Delta]$ meet (the centre of) X. This suggests that the obvious way to attack the proof of II(2.2) is to let Z = A and let c_1 , c_2 be the c_1 in II(2.2), and we will eventually do this. But it turns out that in this case there is no sequence $A \supset A_j \supset A_{jk} \supset \ldots$ such that each of A, A_j , A_{jk} , ... has Property Q_A , c_1 , c_2 , with the c_1 defined as in II(2.2); in fact every dogbone $X \neq A$ has Property Q_A , c_1 , c_2 . We overcome this difficulty with the next definition.

A set $\{X_1, \ldots X_m\}$ of dogbones has Property Q_Z, c_1, c_2 iff each X_r , $r=1, 2, \ldots$ m has Property Q_Z, c_1, c_2 with respect to the same pair of mappings f_1 , f_2 , and the same triple Z, c_1, c_2 . If $\{X_1, \ldots X_m\}$ fails to have Property Q_Z, c_1, c_2 , then we will say that $\{X_1, \ldots X_m\}$ has Property Q_Z, c_1, c_2 . If the set of components of some a_s has Property Q_Z, c_1, c_2 and if $g_1: \Delta \to E^3$ is an extension of c_1 , i=1,2, and the g_1 are Z-disjoint; then some component X of a_s fails to have Property Q_Z, c_1, c_2 with respect to the g_1 . As we saw earlier, this means that both $g_1[\Delta]$ meet X. We will say that a_s has Property Q_Z, c_1, c_2 iff the set of components of a_s has Property Q_Z, c_1, c_2 . Eventually we will show that each of a_1, a_2, a_3, \ldots has Property Q_Z, c_1, c_2 .

- 3. We now give our version of [12, Th 10].
- (3.1). Let $Z \supset A$ and c_1 , c_2 by any circles whatever in $E^3 Z$. In particular, the c_1 do not necessarily link the eyes of A. Then if $\{A_1, A_2, A_3, A_4\}$ has Property Q_{Z,c_1} , c_2 , so has A.

We remark that in [12], the proof of Th 10 does not use the fact that Bd D_1 , Bd D_2 (in fig 1 of [12]) link the eyes of A, even though a short proof of [12, Th 10] can be constructed along the lines of the second paragraph of [12 §7]. The reason is that in later applications of the argument of the proof of [12, Th 10] (which is a disguised induction step) to, say, A_1 and A_{11} , A_{12} , A_{13} , A_{14} , the Bd D_1 do not in fact link the eyes of A_1 . For a similar reason we state (3.1) for very general circles C_1 rather than the C_1 in fig 28. We assume that

Z, c_1 , c_2 have been chosen once and for all before the proof of (3.1) begins, and will now write Property Q for Property Q Z, c_1 , c_2 . We will not refer to Bing's Property Q again in this paper. We will continue the convention in Ch III that i=1, 2, and j=1, 2, 3, 4.

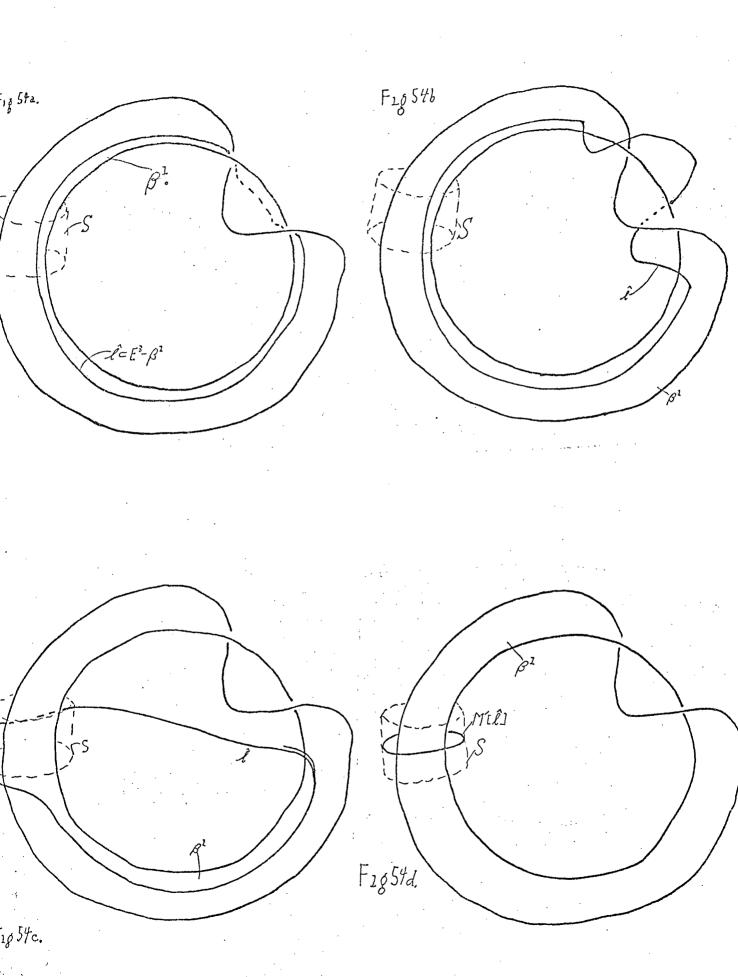
Proof of (3.1): Part I

In this part of the proof we assume that $\{A_1, A_2, A_3, A_4\}$ has Property ${}^{}\!Q$ with respect to mappings g_1 , g_2 and show that the g_i can be replaced by Z-disjoint mappings $g_1^{\prime}: \Delta \rightarrow E^3$ with $g_1^{\prime} = c_1$ on $Bd\Delta$ and with the property that in the cradle $c_1 \cup b_1 \cup b_3 \cup b_4 \cup c_2$ (see fig 43), each b_j misses one rng g_1^{\prime} while $c_1 \cup c_2$ misses both. By the definition of Property $c_1^{}\!Q$, each $c_2^{}\!$ has Property $c_2^{}\!P$ with respect to the rng $c_1^{}\!R$. Look at $c_1^{}\!R$ $c_1^{}\!R$ $c_2^{}\!R$ and recall the definition of bridging in I §5. The construction of the $c_1^{}\!R$ divides into three cases depending on the way that the sets rng $c_1^{}\!R$ $c_1^{}\!R$ bridges $c_1^{}\!R$ bridge $c_1^{}\!R$ and $c_1^{}\!R$. If rng $c_1^{}\!R$ $c_1^{}\!R$ or rng $c_2^{}\!R$ $c_1^{}\!R$ bridges $c_1^{}\!R$ but not both, then we say that $c_1^{}\!R$ is bridged once by rng $c_1^{}\!R$ or rng $c_2^{}\!R$ respectively. If both sets rng $c_1^{}\!R$ $c_1^{}\!R$ rng $c_2^{}\!R$ $c_1^{}\!R$ bridge $c_1^{}\!R$ then $c_1^{}\!R$ is said to be bridged twice. The bridging of $c_1^{}\!R$ is defined anologously. The three cases (not exclusive) are

Case one. Each k has Property $\sim P(a)$; neither of β^1 , β_1 is bridged twice.

Case two. Some k have Property $^{\circ}P(b)$; neither of $^{\beta}$, $^{\beta}$ is bridged twice.

Case three. One of β^1 , β_1 is bridged twice.



These cases are clearly exhaustive (taking 'one' in case three to mean 'at least one'; however the reader has probably noticed that if one of β^1 , β_1 is bridged twice, the other cannot be bridged even once).

Case one. Since each k_i misses one rng g_i , this case suggests an immediate application of Lemma One. It is easily seen that the hypothesis of Lemma One is satisfied except for the fact that the rng $\mathbf{g}_{\underline{\mathbf{1}}}$ may hit Ω . If this happens, we alter the $\,\mathbf{g}_{\underline{\mathbf{1}}}\,$ by means of the following argument: assume that \mathbf{k}_1 misses $\mathrm{rng}\ \mathbf{g}_1$ and \mathbf{k}_4 misses rng g_2 (if another pair of k_i miss the rng g_i or if all four miss the same rng $\mathbf{g}_{\mathbf{i}}$, the method is similar or easier). Since $\ell_{\mathbf{i}}$ misses rng g_1 , there is a circle $\ell \subset E^3 - \beta^1$ which lies near ℓ_1 and approximates it so that $\,\ell\,$ misses rng ${\bf g}_1$. We imagine $\,\ell\,$ sliding on the surface of the twisted band $\,\beta^{\,1}\,$ and eventually coming to rest directly over ℓ_{4} . Although we use the term 'slide', we intend that ℓ stays close to but does not touch $\ensuremath{\beta^1}$. By sliding ℓ on the side of $\ensuremath{\beta^1}$ which is free of the arcs a_i , we are assured that ℓ can move without touching the a. This shows that there is a homeomorphism M of E^3 onto itself such that M is fixed on $E^3 - K_1$ and on $\beta^1 \cup a_1 \cup a_2 \cup a_3 \cup a_4$; and carries ℓ to a position directly over $\mathtt{M}[\ell_4] = \ell_4$. Clearly $\mathtt{M}[\ell]$ misses rng Mg $_1$, ℓ_4 misses rng Mg $_2$. Construct a small annulus α so that its boundary components are M[ℓ] and ℓ_{4} . This can be done so that α misses k_1 , k_2 , k_3 and k_4 - ℓ_4 . By Th 5 (in Ch I), Int α contains a simple closed curve $\hat{\ell}$ which bounds no disk in α and which misses both rng Mg . Figs 54a, ..., d show how $\hat{\ell}$ may be moved to the location of an equator of S without hitting $\beta^1 \cup k_1 \cup k_2 \cup k_3 \cup k_4$.

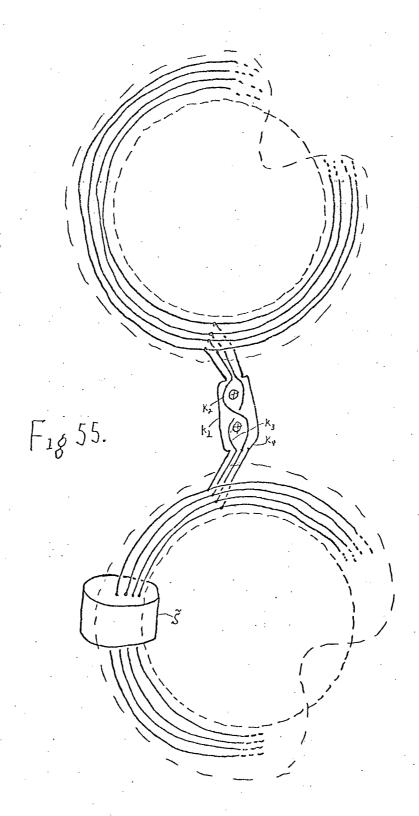
This shows that there is a homeomorphism M' of E^3 onto itself which fixes β^1 , every k_1 , $E^3 - K_1$, and carries $\hat{\ell}$ onto the location $M'[\ell]$ shown in fig 54d. Evidently $M'[\hat{\ell}]$ misses both $M'Mg_{\hat{i}}[\Delta]$; and in fact we can assume that all of Ω misses both M^Mg_i[Δ], since otherwise an obvious homeomorphism can be used to push the ${
m M^2Mg}_{i}$ away from Ω . Note that the M'Mg continue to be Z-disjoint and M'Mg = c on Bd Δ , while each k_1 has Property $^{\circ}P(a)$ with respect to the M $^{\circ}Mg_1[\Delta]$ because both $\, \text{M}' \,$ and $\, \text{M} \,$ are fixed on each $\, k_{\underline{1}} \,$. We can now apply Lemma One to construct Z-disjoint mappings \overline{g}_i with $\overline{g}_i = c_i$ on BdA, such that $\operatorname{rng} \overline{g}_i \subset \operatorname{rng} \operatorname{M'Mg}_i \cup \operatorname{K}_1$ and (since Ω misses both $\operatorname{rng} \operatorname{M'Mg}_i$) rng g_i misses every k_i that rng g_i misses. Since $\zeta_1 \subset$ Int S-N, both rng \overline{g}_i miss ζ_1 . Since M and M' are fixed outside of K_1 , rng \overline{g}_i < rng g_i \vee K_1 . Now apply a result like Lemma One to 'the β_1 end' of $\bigcup_{i=1}^{n} k_i$ to construct Z-disjoint mappings g_i such that $g_i = c_i$ on Bd Δ , rng g'_i \subset rng g'_i \cup K_2, and rng g'_i misses ζ_2 as well as any k_{i} that rng \overline{g}_{i} misses. It may be necessary to alter the \overline{g}_{i} with homeomorphisms which act like M, M' above, in order to make rng g_i miss those k_i which rng g_i misses. Evidently rng g_i misses both ζ_{i} and every k_{i} that rng g_{i} misses. Since each k_{i} misses one rng g_i , the cradle of A has the required property.

Case two. In this case we allow some of the k_j to have Property $^{\circ}P(b)$ with respect to the rng g_i . We reduce this case to Case one by converting the k_j with Property $^{\circ}P(b)$ to Property $^{\circ}P(a)$ or, more accurately, we will define mappings $G_{i0} = g_i$, $^{\circ}G_{i1}$, G_{i2} , $^{\circ}G_{i3}$, $^{\circ}G_{i4}$ with the usual properties such that k_j has Property $^{\circ}P(a)$ with respect to rng G_{ij} , rng G_{2j} , and in fact each k_j misses one rng G_{i4} .

The argument then reduces to Case one.

We will show how $G_{i,1}$ is constructed and indicate the construction of the other $G_{i,i}$. If k_1 has Property ${}^{\diamond}P(a)$ with respect to the rng g_i , then let $g_i = G_{i0} = G_{i1}$. If k_1 has Property P(b) with respect to rng G_{10} , rng G_{20} , we assume that $a_1 \cup m_1$ misses rng $^{\rm G}_{
m 1o}$ $^{
m U}$ rng $^{\rm G}_{
m 2o}$ and that ℓ_1 has Property $^{
m VP}({
m b})$ (as defined in I §5 for a circles with basepoint $\ell_1 \cap a_1$), since otherwise we simply 'turn the picture upsidedown'. Now by Th 6 or Th 7 (in I §5), since at most one of rng $G_{10} \cap \beta^1$, rng $G_{20} \cap \beta^1$, bridges β^1 there is a circle $\ell \subset Int \beta^1$ which bounds no disk in β^1 , contains the base point $\ell_1 \cap a_1$ and misses one of the $\operatorname{rng} G_{io}$, say, $\operatorname{rng} G_{1o}$. We now have a centre (or at least a double ended lasso) with Property $^{\circ}P(a)$ since ℓ $^{\circ}$ $^{\circ}$ u $^{\circ}$ umisses rng \mathbf{G}_{10} ; but ℓ is likely to be a ver disorderly circle and among other delinquencies, probably hits $k_2 \cup k_3 \cup k_4$ (which means that ℓ \prime \cup $\mathbf{a_1}$ \cup $\mathbf{m_1}$ can't be used in Lemma One (the construction of R in Lemma One absolutely requires disjoint $\,\ell_{\, extstyle i}^{\,}$) . We get disjoint loops and a picture like fig 44 by the following procedure which recalls the manipulation of ℓ in Case one. Let λ be a simple closed curve which lies near and approximates $\ensuremath{\ell}'$ but misses $\ensuremath{\beta}^1$. A short straight arc a connects $\ell_1 \cap a_1$ to a base point on λ so that a meets λ at only one point. We can assume that $\lambda\ \upsilon$ a misses rng ${\tt G}_{10}$, and we now regard $\lambda \cup a \cup a_1 \cup m_1$ as a double ended lasso which misses rng $G_{10} = \text{rng g}_1$. Now slide λ over β^1 , keeping the base point fixed, so that the final position of λ is directly over ℓ_1 . As before, we choose the 'right' side of β^1 to slide λ on so that λ will miss $a_1 \cup a_2 \cup a_3 \cup a_4$.

We now have a double ended lasso which looks like k_1 except that the upper loop rides near but not on β^{1} , and it remains only to telescope a υ λ so that a collapses and λ moves to the location of ℓ_1 . We conclude that there is a homeomorphism M^{II} of E^{3} onto itself which fixes $E^3 - K_1$, k_2 , k_3 , k_4 , and carries $\lambda \cup a \cup a_1 \cup m_1$ onto k_1 . Evidently $M''G_{10}$ has the required properties of G_{11} . Clearly k_1 misses G_{11} , and since rng $G_{11} \cap k_2 = \text{rng } G_{10} \cap k_2 = k_2$ continues to have Property $^{\circ}P$ with respect to the rng G_{i1} , and a similar argument applies to k_3 , k_4 . Since M is not fixed on β^1 , we must ask how the rng \boldsymbol{G}_{i1} bridge $\boldsymbol{\beta}^1$. It is clear that $\boldsymbol{\beta}^1$ is bridged at most once by the rng \mathbf{G}_{11} , since ℓ_1 , which misses rng \mathbf{G}_{11} , separates the boundary components of β^1 ; and of course β_1 is bridged by the rng G_{11} just as it was bridged by the rng G_{10} . Thus neither of β^1 , β_1 is bridged twice by the $\operatorname{rng} \mathsf{G}_{11}$. If k_2 misses one of $\operatorname{rng} \mathsf{G}_{11}$, $\operatorname{rng} \mathsf{G}_{21}$, then let $G_{i,1} = G_{i,2}$, Otherwise k_3 has Property $^{\circ}P(b)$ with respect to the rng G_{i1} , and we construct G_{i2} so that rng $G_{12} \cap (k_1 \cup k_3 \cup k_4) = \text{rng } G_{11} \cap (k_1 \cup k_3 \cup k_4), \text{ and } k_2 \text{ misses one}$ of the rng G_{12} (note that we may have to work at the lower end of the figure; the fact that neither band β^1 or β_1 is bridged twice by the rng G_{11} is used in the second application of Th 6 or Th 7). Evidently $\mathbf{k_1}$ misses one of the rng $\mathbf{G_{i2}}$. Proceeding in the same way, we define G_{13} so that k_3 misses one rng G_{13} , and since G_{13} can be constructed so that rng $G_{13} \cap (k_1 \cup k_2) = \text{rng } G_{12} \cap (k_1 \cup k_2)$ each of k_1, k_2 misses one rng G_{13} . Finally define G_{14} so that each k_1 misses one rng G_{14} . Evidently the G_{14} can be used in the argument of Case one to construct the g_1' . When altering the g_1 to G_{11} , G_{11} to G_{12} ,



etc., we preserve 'Z-disjointness' because we adjust only points in Z . Similarly each $G_{ij}=c_{ij}$ on $Bd\Delta$.

Case three. In this case we know only that one of β^1 , β_1 , say β^1 is bridged twice. It is easy to see that if β^1 is bridged that some ℓ_{j_0} misses, say, rng g_1 ; then by I(1.7), rng $g_1 \cap \beta^1$ cannot bridge β^1 , so that the number of bridges is at most one. But if each k_i has Property $\circ P(b)$, then in every case, m_i must miss rng $\mathbf{g}_1 \cup \mathbf{rng} \ \mathbf{g}_2$ and ℓ_i must have Property $^{\mathrm{o}}\mathbf{P}(\mathbf{b})$. For evidently if any $\ell_i \subset E^3$ - rng g_1 - rng g_2 , then there can be no bridges at all. We are thus led to the conclusion that when case three holds, there is just one possible configuration (assuming that β^1 is bridged twice): β^1 is bridged twice, β_1 is bridged not even once, and each k, has Property $^{P}(b)$ with respect to rng g_1 , rng g_2 , with $m_1 \cup a_1 \subset E^3 - \text{rng } g_1 - \text{rng } g_2$. Except for the fact that ζ_2 may not miss rng $g_1 \cup rng g_2$, the picture begins to resemble fig 45, (though we still must construct the arcs u_{13} , etc.). We first alter the g_{i} so that ζ_2 misses rng \mathbf{g}_1 υ rng \mathbf{g}_2 . This is done just as in Case one. Fig 55 shows the $k_{\mbox{\scriptsize 1}}$ and a sphere $\overset{\circ}{S}$ placed in the usual way with respect to ζ_2 and the m_i . In fig 55, the k_i do not have Property ${}^{\circ}P(a)$, so that we use Lemma One itself and not the corollary. Using the method of Case one, construct Z-disjoint mappings $g_1: \Delta \rightarrow (\text{rng } g_1 \cup K_2) - \zeta_2$ such that $\operatorname{rng} \overline{g}_{i}$ misses every m_{i} that $\operatorname{rng} g_{i}$ misses. This simply means that $\operatorname{rng} \overline{g}_1 \cup \operatorname{rng} \overline{g}_2$ misses each m_1 . An examination of the method of Case one shows that if rng g_1 misses a_1 , so does rng g_2 ;

thus $\operatorname{rng} \overline{g}_1$ $\operatorname{rng} \overline{g}_2$ misses all four $\operatorname{m}_j \cup \operatorname{a}_j$. We also know that each point of ℓ_j which misses $\operatorname{rng} g_1$ also misses $\operatorname{rng} g_1$; this means that ℓ_j has Property $\operatorname{P}(b)$ with respect to $\operatorname{rng} g_1$, $\operatorname{rng} g_2$. Therefore the four k_j have Property P with respect to the $\operatorname{rng} g_1$. On the other hand, the fact that the inclusion $\operatorname{rng} g_1 \subset (\operatorname{rng} g_1 \cup K_2) - \zeta_2$ may be proper means that the number of bridges on β^1 with respect to $\operatorname{rng} g_1 \cap \beta^1$, $\operatorname{rng} g_2 \cap \beta^1$ may not be two, but may be one or zero. If this happens, then, since the number of bridges on β_1 with respect to $\operatorname{rng} g_1 \cap \beta_1$, $\operatorname{rng} g_2 \cap \beta_1$ is zero (because of the presence of, say, $\operatorname{m_1}^{\subset E^3}$ - $\operatorname{rng} g_1$ - $\operatorname{rng} g_2$, using a previous argument) we have reduced the situation to either Case one or Case two, i.e. we have each k_j with Property P with respect to the $\operatorname{rng} g_1$ and neither β^1 nor β_1 is bridged twice. However in the 'worst case', β^1 continues to be bridged twice.

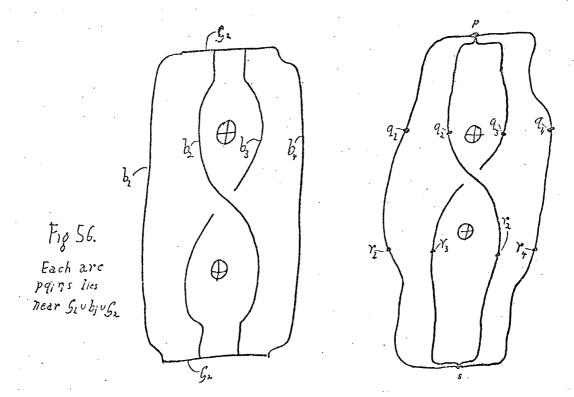
Lemma Two. The hypothesis of Lemma Two is satisfied except that we must construct u_{12} , u_{13} , v_{12} , v_{13} . Since $\operatorname{rng} \overline{g}_2 \cap \beta^1$ bridges β^1 , there is a component Q of $\operatorname{rng} \overline{g}_2 \cap \beta^1$ which connects the boundary components of β^1 . Q is compact and misses $\operatorname{rng} \overline{g}_1$. By the definition of Property $\operatorname{P}(b)$, Q meets a continuum $e_1 \subset \ell_1$ and a continuum e_2 in ℓ_2 such that e_i contains $a_i \cap \ell_i$ and misses one of $\operatorname{rng} \overline{g}_1$, $\operatorname{rng} \overline{g}_2$. Since e_i hits $Q \subset \operatorname{rng} \overline{g}_2$, e_i must miss $\operatorname{rng} \overline{g}_1$. Since the whole continuum $e_1 \cup Q \cup e_2$ misses $\operatorname{rng} \overline{g}_2$, we use $\operatorname{I}(2.5)$ to construct an arc u_{12} which joins $\ell_1 \cap a_1$ and $\ell_2 \cap a_2$ in ℓ_3 and misses $\operatorname{rng} \overline{g}_3$. The constructions of u_{13} , v_{12} , v_{13} are similar.

Now by Lemma Two there are Z-disjoint mappings $g_{\hat{i}}: \Delta \to (\operatorname{rng} \overline{g}_{\hat{i}} - \operatorname{Int} \Sigma) \cup n$ (where Σ , n are the sets described in Lemma Two) with $g_{\hat{i}} = \overline{g}_{\hat{i}} = c_{\hat{i}}$ on $\operatorname{Bd}\Delta$, and such that one $\operatorname{rng} g_{\hat{i}}$, say $\operatorname{rng} g_{\hat{i}}$, misses $b_1 \cup b_2 \cup b_3$ while both $\operatorname{rng} g_{\hat{i}}$ miss $c_1 \cup c_2$. Evidently $\operatorname{rng} g_{\hat{i}} \subset \operatorname{rng} g_{\hat{i}} \cup K_1$.

In the argument of Case three we did not succeed in constructing the g_1' so that $\zeta_1 \cup \zeta_2$ misses rng $g_1' \cup$ rng g_2' and each b_j misses one rng g_1' ; instead $\zeta_1 \cup \zeta_2$ misses both rng g_1' and three b_j miss the same rng g_1' . In Part II of the proof of (3.1), it turns out that it is sufficient to define the g_1' so that three b_j miss the same rng g_1' (the same thing happens in the proof of [12, Th 10]). With some additional complication, it is possible to improve the argument of Lemma Two so as to yield the usual result, i.e. to construct g_1' so that each b_j misses one rng g_1' ; however we omit this argument.

We have now completed the three cases of the proof of Part I of (3.1). Note that in each Case, we constructed g_i so that rng $g_i \subset \operatorname{rng} g_i \cup K_1 \cup K_2$. Thus we can write $\operatorname{rng} g_i \subset \operatorname{rng} g_i \cup A$. This will be important when we apply the argument of (3.1) to the components of a_2 , a_3 , etc. To summarize the situation: if $\{A_1, A_2, A_3, A_4\}$ has Property Q with respect to g_1 , g_2 , then there exist Z-disjoint mappings $g_i : \Delta \to E^3$ such that

 $g_i' = c_i$ on $Bd\Delta$, $rng g_i' \subset rng g_i \cup A$, $if c_1 \cup c_2 \cup b_1 \cup b_2 \cup b_3 \cup b_4$ is the cradle of A, then both $rng g_i' = c_i$ or three b_i miss the same $rng g_i'$.

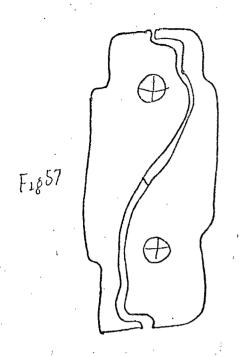


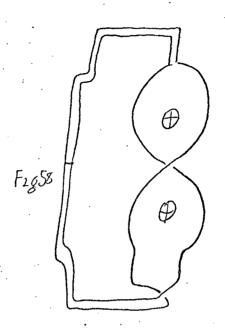
Part II of the proof of (3.1).

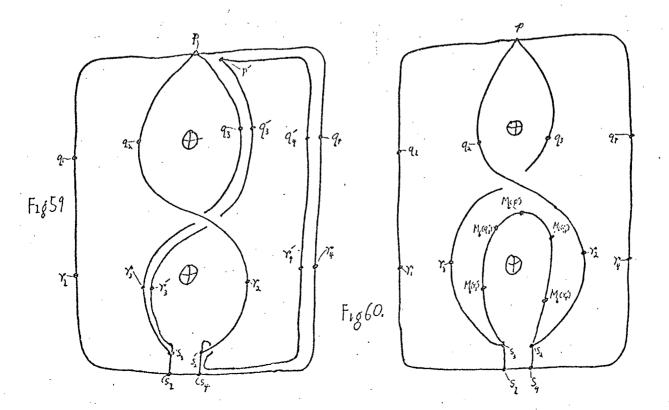
We remind the reader that we are proving a result much like Bing's Th 10 of [12], which is also divided into a Part I and Part II.

Our Part II is very similar to Part II in Bing's proof and we absolutely require familiarity in detail with Bing's Part II (this is only a matter of half a page). We think it likely that the reader sees from the proof in [12], how to complete Part II here, and instead of a formal proof, we will give what amounts to a gloss on Bing's method, plus a few comments required by the fact that our Property Q is not quite identical to Bing's.

We begin by replacing $\zeta_1 \cup b_1 \cup \dots \cup b_4 \cup \zeta_2$ by the figure Upq ris shown in fig 56. This can be done so that either each arc pq_ir_is misses one $rng g_i$ or three pq_ir_is miss the same $rng g_i$. Our terminology is now like that of [12] except that rng g_i^* replaces D, in [12]. We follow the division into cases found in [12]. We will not prove that the three cases given in [12] exhaust the possibilities, but remark for plausibility that the case division ... 1) Three pq.r.s miss one rng g₁, 2) pq₁r₁s plus pq₂r₂s misses rng g₁, pq₃r₃s plus pq_4r_4s misses rng g_2 , 3) pq_1r_1s plus pq_4r_4s misses rng g_1 , $pq_2^r r_2^s$ plus $pq_3^r r_3^s$ misses $rng g_2^* \dots$ seems at first glance to ignore the possibility: pq_1r_1s plus pq_3r_3s misses $rng g_1$, pq_2r_2s plus pq_4r_4s misses rng g_2^{\prime} . However this last variation is just Case Two with the diagram inverted. We will now describe how Bing's Part II can be altered to show that there exist Z-disjoint mappings $F_i: \Delta \to E^3$ such that $F_1 = c_1$ on $Bd\Delta$ and the centre of A has Property $\circ P$ with respect to $\operatorname{rng} F_1$, $\operatorname{rng} F_2$.

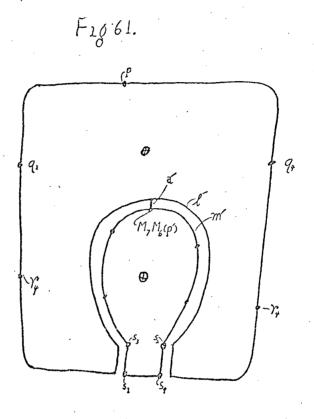






Case One: any three of $\operatorname{pq}_{j}r_{j}s$ (j=1,2,3,4) miss the same $\operatorname{rng} g_{i_0}$. If $\operatorname{pq}_{2}r_{2}s$ is an arc which fails to miss $\operatorname{rng} g_{i_0}$, then the structure shown in fig 57 lies near $\operatorname{pq}_{1}r_{1}s$ $\operatorname{U} \operatorname{pq}_{3}r_{3}s$ $\operatorname{U} \operatorname{pq}_{4}r_{4}s$ and misses $\operatorname{rng} g_{i_0}^*$. The structure in fig 57 can be moved to the position of the centre k of A by a homeomorphism M_{5} which fixes E^3 - A. Evidently $\operatorname{M}_{5}g_{1}^*$, $\operatorname{M}_{5}g_{2}^*$ are the required F_{1} , F_{2} . If $\operatorname{pq}_{4}r_{4}s$ is an arc which fails to miss $\operatorname{rng} g_{i_0}^*$, then one uses the structure in fig 58 which lies near $\operatorname{pq}_{1}r_{1}s$ $\operatorname{U} \operatorname{pq}_{2}r_{2}s$ $\operatorname{U} \operatorname{pq}_{3}r_{3}s$ and misses $\operatorname{rng} g_{i_0}^*$. If $\operatorname{pq}_{1}r_{1}s$ or $\operatorname{pq}_{3}r_{3}s$ fail to miss $\operatorname{rng} g_{i_0}^*$, the method is like one of those already given. If all four $\operatorname{pq}_{j}r_{j}^*r_{j}^*s$ miss $\operatorname{rng} g_{i_0}^*$, then 'forget' one of them.

Case Two. pq1r1s plus pq2r2s misses rng g1, pq3r3s plus pq_4r_4s misses $rng g_2$. We replace $\bigcup_i pq_ir_is$ with the more complicated construction in fig 59. In fig 59, s has been replaced by s_1 , s_2 , s_3 , s_4 which lie near s so that the s_1 and arcs s_1s_3 , s_1s_4 , s_2s_4 misserng g_1 U rng g_2 . Abusing the notation slightly, we have arcs $pq_jr_js_j$ with $pq_1r_1s_1 \cup pq_2r_2s_2 \subset E^3 - rng g_1$. $pq_3^rq_3^sq_3$ $pq_4^rq_4^sq_6$ E^3 - rng g_2^s . We build two new arcs: $p^*q_3^*r_3^*s_3$, which lies near $pq_3r_3s_3$ and misses rng g_2 , and $p'q_4'r_4's_2$ which lies near $pq_4r_4s_4$ and also misses $rng~g_2^\prime$. Apply a move M_6 which carries $s_3 r_3 q_3 p_4 r_4 s_4$ to the location shown in fig 60 and fixes $pq_1r_1s_1$, s_1s_3 , s_2s_4 , and s_1s_4 . Look at a disk in A bounded by the circle $pq_1r_1s_1s_3M_6(r_3)M_6(q_3)M_6(p_5)M_6(q_4)M_6(r_4)s_2r_2q_2p$. We will call this disk T and assume that it is just the obvious disk suggested by the figure. Thus T misses all but the end points of $pq_4r_4s_4s_2$. Later we will need the fact that T can be constructed so as also to miss all but the end points of $pq_3r_3s_3$ (in Case 3). There is an arc



 λ C T with end points s₂ and s₃ which misses both rng M₆g₁ because arc $s_3^{M_6}(r_3)^{M_6}(q_3)^{M_6}(p_4)^{M_6}(q_4)^{M_6}(r_4)^{S_2}$ misses rng $m_6^{S_2}$ and arc $s_3s_1r_1q_1pq_2r_2s_2$ misses rng $M_6g_1^2$. (We will now begin to abbreviate our arc nomenclature). Define a move M_7 which moves λ to the position of arc $s_3^{M_6}(p)s_2$ and fixes each $pq_1^{r_1}s_1^{r_2}$ $\mathbf{s_{3}s_{1}s_{4}s_{2}}$. Although we do not know the location of $\,\lambda\,$ in $\,$ T, $\,$ this can be done by means of the A - move defined in I §3. Evidently rng $M_7M_6g_1$ misses $pq_1r_1s_1$, rng $M_7M_6g_2$ misses $pq_4r_4s_4$, and both rng $M_7M_6g_1$ miss the circle $s_1s_3M_6(p)s_2s_4s_1$. Fig 61 shows $\mathtt{pq_1r_1s_1} \quad \textit{U} \quad \mathtt{pq_4r_4s_4} \ \textit{U} \ \ \mathtt{s_1s_3M_6(p')s_2s_4s_1} \quad \text{replaced by a set} \quad \textit{\ell'} \ \textit{U} \ \ \texttt{a'} \ \textit{U} \ \ \texttt{m'}$ which lies very near the first set so that $m' \cup a'$ misses both rng $M_7M_6g_i$, and ℓ has Property $\sim P(b)$ with respect to the rng $M_7M_6g_i$ (it is easy to give ℓ this property since much of ℓ can coincide with $s_1^r q_1^p q_4^r q_5^4$). Evidently $\ell' u$ a'u m' can be moved to the by a move M_8 , then the centre of A has Property $^{\circ}P(b)$ with respect to the rng $^{M}8^{M}7^{M}6^{g_{i}}$, which we define to be the required ^{F}i .

Case three. pq_1r_1s plus pq_4r_4s misses $rng\ g_1'$, pq_2r_2s plus pq_3r_3s misses $rng\ g_2'$. The mechanism of this case resembles that of Case two. We repeat the construction in fig 59 and define M_6 precisely as in Case two, so that we arrive once more at fig 60. However, since the $rng\ g_1$ are related differently to the various parts of the figure, we have this time: $s_3s_1s_4s_2\subset E^3-rng\ M_6g_1-rng\ M_6g_2$ as usual, but $pq_1r_1s_1\cup pq_4r_4s_4\cup M_6(p^*)M_6(r_4')s_2\subset E^3-rng\ M_6g_1$, $pq_2r_2s_2\cup pq_3r_3s_3\cup M_6(p^*)M_6(r_3')s_3\subset E^3-rng\ M_6g_2$. In this case we

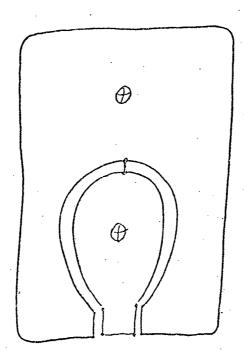


Fig 62.

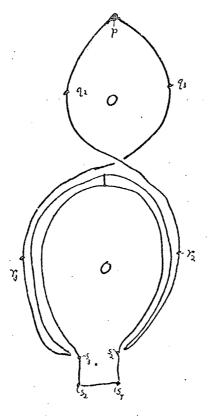


Fig 63.

must use a fact that we stated but did not completely use in Case 2, viz. that M_6 fixes all four pq_ir_is . We assume tht the disk T is placed so as to miss $pq_3r_3s_3$. We use Th 4 from I §4 at this point; at the analogous place in [12], Th 7 of [12] is used. By Th 4, since $s_3^{M_6}$ (q3)M₆(p1) and pq2r2s2 miss rng M₆g2 and M₆(p1)M₆(q4)s2 and $s_3 s_1 r_1 q_1 p$ miss rng $M_6 g_1$, there is an arc $\overline{\lambda} \subset T$ with end points $\mathbf{s}_3, \ \mathbf{s}_2, \ \text{such that} \ \overline{\lambda} \ \text{misses either} \ \text{rng} \ \mathbf{M}_6^{} \mathbf{g}_1^{} \ \text{or} \ \text{rng} \ \mathbf{M}_6^{} \mathbf{g}_2^{}$. Apply a move M_7 , similar to M_7 , to move $\overline{\lambda}$ to $s_3^{M_6}(p^2)s_2$. This can be done by an A - move as before; but some care should be taken so that M_7 fixes every pq₁r₁s₁ (as well as, of course $s_3s_1s_4s_2$); the reader might first prefer to move $pq_3r_3s_3$ to a new location where it cannot interfere with the collar of T used in the A - move. The proof is now completed along the lines of the previous cases, using the fact that if $s_3^{M}_6(p)s_2$ misses rng $M_7^{M}_6s_1$, then the set shown in fig 62 lying near $s_1 s_3 M(p^2) s_2 s_4 s_1 \cup s_1 r_1 q_1 p q_4 r_4 s_4$ misses rng $M_7 M_6 g_1$; while if $s_3^{M_6}(p')s_2$ misses rng $M_7^{M_6}g_2$, then the set in fig 63 lying near $s_1 s_3 M_6 (p^2) s_2 s_4 s_1 v s_3 r_3 q_3 p q_2 r_2 s_2$ misses rng $m_7 M_6 g_2$. This completes part II of the proof of III(2.1) .

Corollary to III(2.1). If $\{A_1, \ldots A_4\}$ has Property $\sim Q$ with respect to mappings g_1 , g_2 , then A has Property $\sim Q$ with respect to mappings F_1 , F_2 such that rng $F_1 \subset \operatorname{rng} g_1 \cup A$.

Proof. We know that rng g_1 c rng $g_1 \cup A$. And all the moves given in Part II of the proof of III(2.1) can be defined so as to fix $E^3 - A D$.

4. Proof of II(2.2).

We have now shown that if $\{A_1, \ldots A_4\}$ has Property ${\bf Q}$ then A has Property $\circ Q$. Our argument now diverges somewhat from Bing's in [12]. Suppose that $f_i: \Delta \to E^3$ such that $f_i|_{Bd\Delta} = c_i$ and the f_i are Z-disjoint (we continue to take Z, c_1 , c_2 to be assigned arbitrarily according to the remark at the beginning of §3). We will show that if A has Property Q, then each of a_1 , a_2 , a_3 , ... has Property Q; the proof of II(2.2) follows directly from this fact. If A has Property Q, then by (3.1), $\{A_1, \ldots, A_L\}$ has Property Q. i.e. a_1 has Property Q . (3.1) does not imply that some A_i has Property Q for the luminous reason that each A_{i} has Property $\sim Q_{i}$, as the argument in §1 Ex 2 shows. However we can show that a_2 has Property Q by adapting the argument of the proof of (3.1) to show that for each A_j , if $\{A_{j1}, \dots A_{j4}\}$ has Property \mathcal{Q} then so does A_j . This is easy to do since the proof is simply restated in terms of images under the embedding $\,h_{\,\mathbf{i}}\,$ of various subsets of A . Occasionally in the proof of III(2.1) we constructed arcs which were perpendicular to certain surfaces. While h, does not preserve this property, the reader will appreciate that we used such constructions for topological purposes, e.g. to make one arc lie along another, or to miss certain subsets, and these properties are preserved by h_j . We do not re-define Z, c_1 , c_2 of course, since we intend to show that the same Property $\mathbf{Q}_{\mathbf{Z_i},\mathbf{C_i}}$, \mathbf{c}_2 is possessed by each of a_1, a_2, a_3, \ldots . We originally defined Z to contain A so that we have $Z \supset A_i$ as required. We intend of course to let Z = A eventually. To show that a_2 has Property Q, assume that

the set of components of a_2 has Property ${}^{\diamond}Q$ with respect to qualified mappings g_1 , g_2 . Apply a result like the corollary of (3.1) to $\{A_{11}, \ldots A_{14}\}$ to obtain Z-disjoint mapping $F_{i1}: \Delta \rightarrow \text{rng g}_i \cup A_1$ such that $F_{i1} = c_i$ on BdA, and A_1 has Property Q with respect to the $\mathbf{F}_{\mathbf{1}\mathbf{1}}$. We can see that since rng $\mathbf{F}_{\mathbf{1}}$ does not exceed rng $\mathbf{f}_{\mathbf{1}}$ $E^3 - A_1$, the dogbones A_{21} , ... A_{24} , A_{31} , ... A_{34} , A_{41} , ... A_{44} tinue to have Property $\circ Q$ with respect to the F_{i1} , for as we saw earlier, possession of Property ${}^{\circ}P$ depends on the fact that rng g $_{\hat{t}}$ misses certain continua in various dogbones, and this property is inherited by rng Fi. at least for dogbones in E^3 - A_1 . Construct Z-disjoint mappings $F_{12}:\Delta \rightarrow$ rng F_{i1} U A_2 such that $F_{i2} = c_i$ on Bd Δ and A_2 has Property Q with respect to the F_{i2} . Once again, degbones in $E^3 - A_2$ which have Property Q with respect to the Fill continue to have Property Q with respect to the File. This means that not only $^{A}_{2}$, but $^{A}_{1}$, $^{A}_{31}$, \cdots $^{A}_{34}$, $^{A}_{41}$, \cdots $^{A}_{44}$ have Property ${}^{\sim}\!Q$ with respect to the ${\rm F}_{\rm i\,2}$. Evidently we can continue in this way and finally derive Z-disjoint mappings $F_{i4}: \Delta \rightarrow E^3$ which agree with c_1 on Bd Δ and with respect to which, all of A_1 , ... A_{Δ} have Property ${\scriptstyle \sim} {\rm Q}$. Assume that the set of components of a_3 have Property ${\scriptstyle \sim} {\rm Q}$. Then an argument like that of (3.1) Corollary can be applied to each A_{ik} in (perhaps lexicographic) order to show eventually that a_2 has Property ${\sim}$ Q . If A has Property Q, then by induction, a_1 , a_2 have Property ${\tt Q}$ and ${\tt a}_{\tt Q}$ must also have Property ${\tt Q}$. We think that it is now evident how to proceed in the case that $m = 4, 5, \ldots$

We will show how the induction argument above implies II(2.2). If the f_i in the hypothesis have ranges that intersect in A , then II(2.2) is true; thus we consider only the case that rng $f_1 \cap rng f_2 \cap A = \emptyset$,

i.e. the case that the f_i are A - disjoint. In the preceding argument we showed that for a fixed choice of Z, c_1 , c_2 , if A has Property Q_{Z,c_1} , c_2 , then so does each a_m . If Z=A and c_1 , c_2 are the c_i in II(2.2) then $A \subseteq Z \subseteq E^3$ - rng c_1 - rng c_2 as required, and A has Property Q by an argument like that of \$1 Ex 1. By the induction argument, every a_m has Property Q_{Z,c_1} , c_2 . As we saw earlier, this means that both rng f_i hit some component of a_m for m however large.

Finally we will show that both rng f. must hit a big element Λ of the dogbone decomposition $\,G$. Let $\,\hat{G}\,$ be the set of all elements of the dogbone construction (i.e. all components of a_1 , a_2 , a_3 , ...) which meet both $\operatorname{rng} \, \mathbf{f}_1$ and $\operatorname{rng} \, \mathbf{f}_2$. Evidently $\hat{\mathbf{G}}$ is infinite, for by the arguments of this chapter, each $a_{\rm m}$ must contain an element of \hat{G} . Clearly one of A_1 , ... A_4 must contain an infinite subset of \hat{G} , for the four A_{i} contain all of \hat{G} . If A_{i} contains an infinite subset of \hat{G} , then one of A_{j1} , ... A_{j4} , say A_{jk} , contains an infinite subset of \hat{G} . There is a sequence $A \supset A_j \supset A_{jk} \supset A_{jkl} \ldots$ each of which contains infinitely many dogbones which meet both $\operatorname{rng} f_i$. Obviously each member of the sequence meets both $\operatorname{rng} f_i$, and the intersection A \cap A $_j$ \cap A $_{jk}$ \cap A $_{jk\ell}$ \cap ... meets both rng f $_i$. One can also use the dogbone metric to show that if the images of the $\operatorname{rng} f_{\mathbf{i}}$ are disjoint in \mathcal{D} , then there is a neighbourhood system of the points of \mathcal{D} consisting of small 3-cells around the small points and images of dogbones about the big points such that no neighbourhood of diameter smaller than ϵ (in the dogbone metric) meets both images of the $\mbox{ rng } f_{\hat{\textbf{1}}}$. This implies that some a_m has Property $\sim Q$, cf. proof of Th 12 of [12] \square .

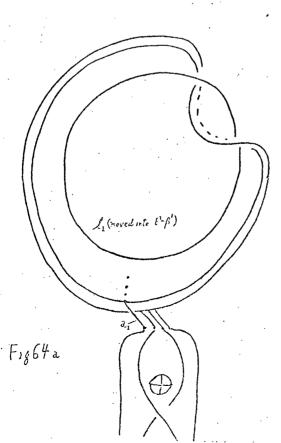
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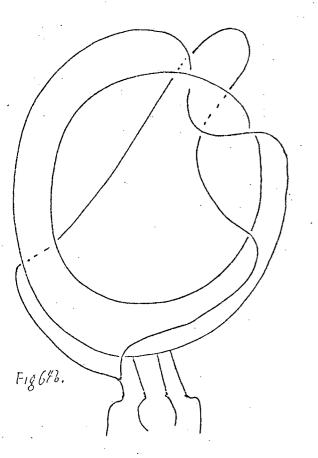
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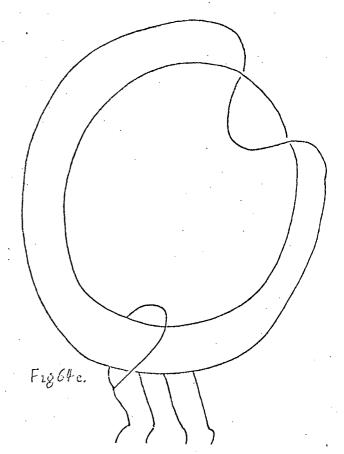
Appendix.

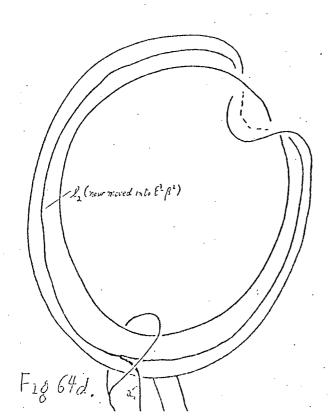
Figs 642,... j show how to deform the upper part of fig 19 so that it Looks like the upper part of Fig 1 of [12].



Push l, off \beta^1. Move a, so that a, passes through \beta^1 and now approaches \beta^1 from the upper side, while a, a, a, a, continue to approach from the lower side.



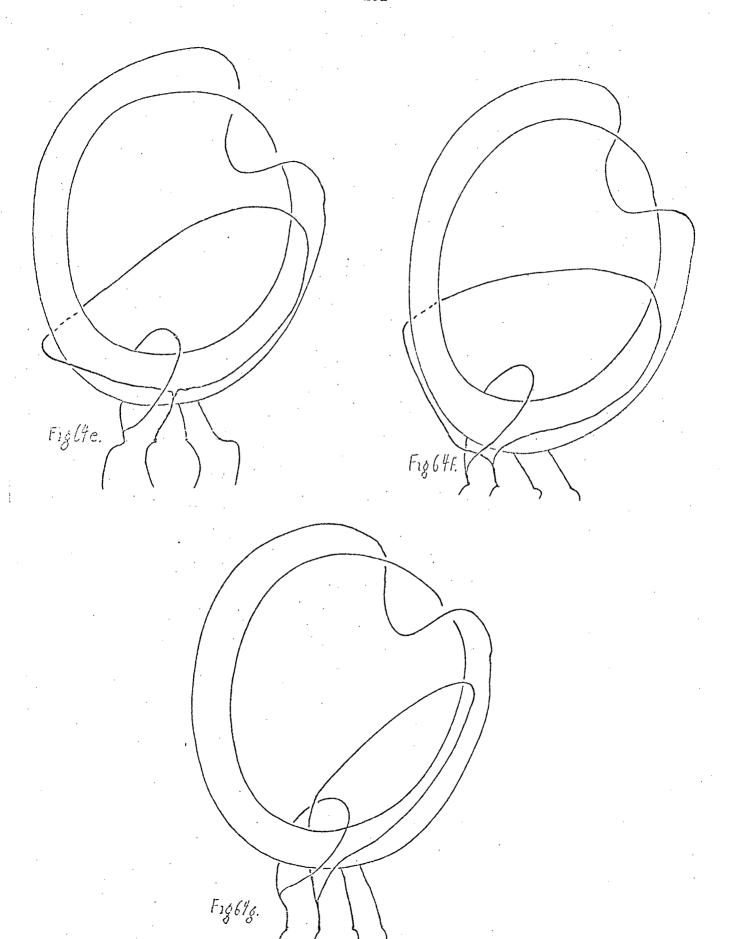


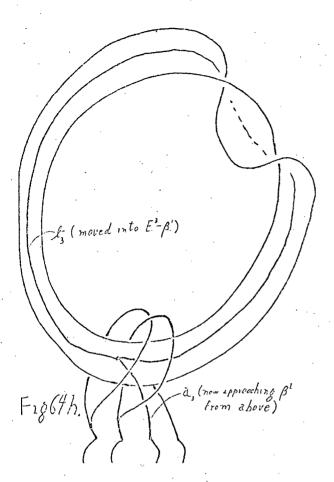


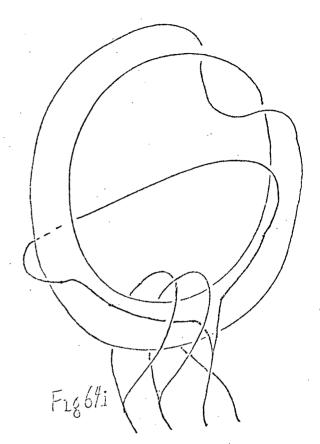
Push off le trom B.

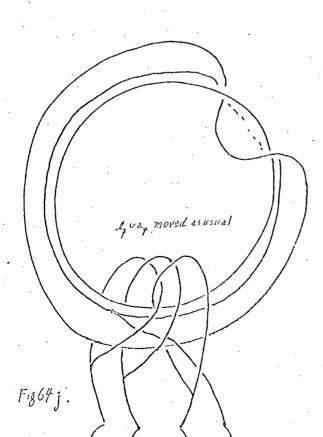
Bring as through B'

so that as new approaches
B' from above.









Etc.