

**DELAY-THROUGHPUT ANALYSIS OF INTER- AND  
INTRA-MAN VOICE AND DATA INTEGRATED  
TRAFFIC IN IEEE 802.6 MAN BASED PCN**

by

**WILLIAM Y. L. WONG**

**B.Sc.(EE), University of Manitoba, 1992**

**A THESIS SUBMITTED IN PARTIAL FULFILLMENT OF  
THE REQUIREMENTS FOR THE DEGREE OF  
MASTER OF APPLIED SCIENCE**

in

**THE FACULTY OF GRADUATE STUDIES  
DEPARTMENT OF ELECTRICAL ENGINEERING**

We accept this thesis as conforming  
to the required standard

**THE UNIVERSITY OF BRITISH COLUMBIA**

**February 1995**

**© William Y. L. Wong, 1995**

In presenting this thesis in partial fulfilment of the requirements for an advanced degree at the University of British Columbia, I agree that the Library shall make it freely available for reference and study. I further agree that permission for extensive copying of this thesis for scholarly purposes may be granted by the head of my department or by his or her representatives. It is understood that copying or publication of this thesis for financial gain shall not be allowed without my written permission.

Department of ELECTRICAL ENGINEERING,

The University of British Columbia  
Vancouver, Canada

Date Feb 27, 1995

## Abstract

This thesis focuses on a proposed IEEE 802.6 Metropolitan Area Network (MAN) based distributed communication architecture for supporting the Personal Communication Services (PCS). Future-generation telecommunications systems are intended to combine all kinds of networks (wireline or wireless) and services (data, voice, video, graphics, etc.) into one single universal personal communications system.

An IEEE 802.6 MAN-based distributed Personal Communication Network (PCN) architecture is introduced in this thesis, and network performance analysis based on this architecture is presented. Isochronous (voice) traffic will be transported by high priority queue-arbitrate (QA) slots instead of by using the more conventional pre-arbitrated (PA) slots. Consequently, the highest priority QA slots will be used to transport signalling traffic while the lowest priority QA slots will be used to transport data traffic. Since analytical models are not available, simulation models have been developed to evaluate the transmission delay of voice and data packets for both inter- and intra-MAN traffic. From these simulation models, nodal transmission delay is found to be depended on the physical location of the transmitting node and traffic levels within the network. Inter-MAN traffic nodal transmission delay is additionally dependent on the connection scheme of homogeneous bridges which interconnect different MANs.

With CCITT Q.931 used as the signalling protocol for call setup, the QA access for isochronous traffic (QAAIT) call setup and clearing is simpler and faster than the PA access for isochronous traffic (PAAIT), QA access avoids specific channel allocation and call clearing. Although QAAIT has its advantages in call setup and clearing, its use under light traffic load only is advised. Under heavy traffic loads PAAIT is advised, since the later case provides bounded packet transmission delay at all times by controlling network call access.

# Table of Contents

Abstract . . . . .	ii
List of Figures . . . . .	vi
List of Tables . . . . .	ix
Acknowledgment . . . . .	x
Chapter 1      Introduction . . . . .	1
Section 1.1      Background and Motivation . . . . .	1
Section 1.2      PCN Architecture Based on IEEE 802.6 MAN . . . . .	2
Section 1.3      Outline of The Thesis . . . . .	5
Chapter 2      Performance Analysis of IEEE 802.6 MAN . . . . .	6
Section 2.1      The IEEE 802.6 MAN . . . . .	6
2.1.1      The IEEE 802.6 MAN architecture . . . . .	6
2.1.2      The distributed queue access protocol . . . . .	8
Section 2.2      Slot Reuse Schemes . . . . .	11
2.2.1      Destination Release . . . . .	11
2.2.2      Previous Slot Information (PSI) . . . . .	11
Section 2.3      Modelling and Simulation of the DQDB Protocol . . . . .	13
2.3.1      DQDB protocol model . . . . .	13
2.3.2      Simulation models . . . . .	14
2.3.3      Simulation results . . . . .	15

Chapter 3	Delay Analysis of QA Access for Integrated Intra-MAN Voice and Data Traffic . . . . .	19
Section 3.1	Voice-Data Integration . . . . .	19
3.1.1	Voice clipping . . . . .	20
Section 3.2	Simulation Models . . . . .	22
Section 3.3	Simulation Results . . . . .	24
3.3.1	Simultaneous voice calls arrival without silence deletion . . . . .	24
3.3.2	Exponentially distributed voice call arrivals without silence deletion . . . . .	28
3.3.3	Exponentially distributed voice call arrivals with silence deletion	32
3.3.4	Exponentially distributed voice calls arrival for different voice coding rate with silence deletion . . . . .	34
Chapter 4	Delay Analysis of QA Access for Inter- and Intra-MAN Voice and Data Integrated Traffic within a Backbone MAN Network . . . . .	42
Section 4.1	Homogeneous Bridges Connection Scheme . . . . .	43
Section 4.2	Simulation Model . . . . .	46
Section 4.3	Preliminary Simulation Results . . . . .	46
Section 4.4	Statistical Results for The Six-MAN Network . . . . .	50
Section 4.5	General Results of The Inter- and Intra-MAN Traffic within a Backbone MAN Network . . . . .	56
Section 4.6	Call Setup Delay . . . . .	61
4.6.1	Signalling protocol . . . . .	61

4.6.2	Call setup procedure . . . . .	63
4.6.3	Performance analysis . . . . .	63
Chapter 5	Summary and Conclusions . . . . .	65
Section 5.1	Results Summary . . . . .	65
Section 5.2	Suggestions for Future Work . . . . .	67
Appendix A	Source Code . . . . .	69
References . . . . .		104

## List of Figures

Figure 1. IEEE 802.6 MAN based distributed PCN architecture . . . . .	3
Figure 2. MAN interconnection within a backbone MAN . . . . .	4
Figure 3. DQDB dual-bus configuration . . . . .	6
Figure 4. Frame and slot format in DQDB MAN . . . . .	7
Figure 5. Uniform destination distribution for data packets in a DQDB MAN . . . . .	16
Figure 6. Geometric destination distribution for data packets in a DQDB MAN . . . . .	17
Figure 7. Active voice station model . . . . .	20
Figure 8. Voice clipping probability vs. number of ready-to-transmit voice stations . . . . .	22
Figure 9. Mean voice packet waiting delay vs. QA voice Throughput . . . . .	25
Figure 10. Percentage of voice packets discarded vs. QA voice Throughput . . . . .	26
Figure 11. Mean data packet waiting delay vs. QA data Throughput for various voice throughput . . . . .	27
Figure 12. Mean packets waiting delay vs. nodes locations . . . . .	27
Figure 13. Mean voice packet waiting delay vs. QA voice Throughput . . . . .	29
Figure 14. Percentage of discarded voice packet vs. QA voice Throughput . . . . .	29
Figure 15. Mean data packet waiting delay vs. QA data Throughput for various voice Throughput . . . . .	30
Figure 16. Mean packet waiting delay vs. node location . . . . .	31
Figure 17. Percentage of discarded voice packets vs. number of voice calls per node . . . . .	34
Figure 18. Mean data packet waiting delay vs. QA data Throughput for various number of active voice calls with voice coding rate of 64 kbps . . . . .	35

Figure 19. Mean data packet waiting delay vs. QA data Throughput for various number of active voice calls with voice coding rate of 32 kbps . . . . .	36
Figure 20. Mean data packet waiting delay vs. QA data Throughput for various number of active voice calls with voice coding rate of 16 kbps . . . . .	36
Figure 21. Mean data packet waiting delay vs. QA data Throughput for various number of active voice calls with voice coding rate of 4 kbps . . . . .	37
Figure 22. Mean data packet waiting delay vs. QA data Throughput with different numbers of active voice calls per node for various voice coding rates . . . . .	38
Figure 23. Relation between voice coding rate and number of active voice calls activated .	40
Figure 24. Geographical location of access MANs within a backbone MAN . . . . .	42
Figure 25. Illustration of three homogeneous bridges connection . . . . .	43
Figure 26. Illustration of four homogeneous bridges connection . . . . .	44
Figure 27. Physical representation of a backbone MAN network . . . . .	45
Figure 28. Mean inter-MAN data packet transmission delay vs. total QA data traffic for 32 kbps voice coding rate . . . . .	47
Figure 29. Mean intra-MAN data packet transmission delay vs. total QA data traffic for 32 kbps voice coding rate . . . . .	48
Figure 30. Mean inter-MAN data packet transmission delay vs. total QA data traffic for 4 kbps voice coding rate . . . . .	48
Figure 31. Mean intra-MAN data packet transmission delay vs. total QA data traffic for 4 kbps voice coding rate . . . . .	49
Figure 32. Mean inter-MAN packets waiting delay over homogeneous bridges . . . . .	52
Figure 33. Mean inter-MAN packet transmission delay via different bridges . . . . .	53

Figure 34. Mean intra-MAN packets transmission delay vs. nodes locations . . . . .	54
Figure 35. Mean inter-MAN data packet transmission delay vs. nodes locations . . . . .	55
Figure 36. Mean inter-MAN voice packet transmission delay vs. nodes locations . . . . .	56
Figure 37. Mean inter- and intra-MAN data packet transmission delay vs. total data traffic for 75% inter-MAN traffic . . . . .	57
Figure 38. Mean inter- and intra-MAN data packet transmission delay vs. total data traffic for 50% inter-MAN traffic . . . . .	58
Figure 39. Mean inter- and intra-MAN data packet transmission delay vs. total data traffic for 25% inter-MAN traffic . . . . .	58
Figure 40. Mean inter-MAN voice packet transmission delay vs. total QA voice traffic for various levels of inter-MAN traffic . . . . .	59
Figure 41. Mean intra-MAN voice packet transmission delay vs. total QA voice traffic for various levels of inter-MAN traffic . . . . .	59
Figure 42. Call setup procedure of Q.931 protocol for PA access. For QA access, signalling involving the Bandwidth Manager and VCI is eliminated. . . . .	62

## List of Tables

Table 1. Slot Access Control Field coding . . . . .	8
Table 2. Behavior of nodes in Figure 3 . . . . .	11
Table 3. Example of PSI . . . . .	12
Table 4. Results summary of simultaneous and exponential voice calls arrival without silence detection . . . . .	31
Table 5. Comparison between models with and without silence deletion . . . . .	33
Table 6. QA voice Throughput achieved for various voice coding rates and number of active voice calls . . . . .	39
Table 7. Delay comparison for different voice coding rate voice stations . . . . .	41
Table 8. Mean inter- and intra-MAN voice packet transmission delay for 50 voice calls per node . . . . .	50
Table 9. A typical simulation results for the six-MAN network . . . . .	51
Table 10. Data packet transmission delay and throughput achieved for different voice throughput levels . . . . .	60
Table 11. Mean inter-MAN call setup delay vs. background PA throughput over 2 transitional MANs . . . . .	64
Table 12. Mean intra-MAN call setup delay vs. background PA throughput . . . . .	64

## **Acknowledgment**

I would like to take this opportunity to thank my research supervisor, Dr. R. W. Donaldson, for his guidance, suggestions and support in producing this thesis. I would also like to thank my family and friends for their encouragement and support in my graduate studies. Special thanks to my colleagues and staff in the department for their kindness and help during my studies in UBC. At last, but not least, thanks to the Canadian Institute of Telecommunications Research for continued support of this research.

# **Chapter 1. Introduction**

## **Section 1.1 Background and Motivation**

The B-ISDN ( Broadband Integrated Services Digital Network ) concept supports telecommunications involving universally available integrated broadband communications networks and services. The B-ISDN is intended to provide an integrated network capable of gathering information from variety sources and delivering it to designated destinations regardless of the service type. The B-ISDN is designed to be capable of handling all kinds of services, including data, voice, video and graphics.

Personal Communication Services ( PCS's ) are expected to have a continuing growth via B-ISDN in the coming decades. The PCS concept enables people to call other people irrespective of the geographical location of either party, using only callee's personal identification number ( PIN ). PCS access to the B-ISDN is via both wireline and wireless terminals [1]. The Personal Communication Network ( PCN ) which supports PCS's must be capable of the following [2]:

- carrying many types of traffic;
- being integrated with current and future wireline communication networks;
- serving a mass market in urban areas;
- operating efficiently in sparsely populated areas;
- operating indoors, outdoors, and in vehicles;
- serving terminals which may move at high speed;

The many distinctive features and potential demands required for the PCN suggests that the centralized communication architectures may not be suitable for implementing the PCN. The

predicted growth in mobile subscribers requires that cell dimensions be reduced substantially from the size currently used in mobile radio system, to improve spectral efficiency through frequency reuses. Reduction in cell radii will increase handoff rates. With the BTS antenna in these microcells situated a few meters above ground, a 20–30 dB drop in signal level can occur at street corners when a mobile terminal ( MT ) loses line-of-sight with the BTS. As a result of the increased processing load arising from the growing demand for wireless access, together with the necessity of a fast handoff due to the propagation effects discussed above, current centralized network and control architectures are probably unsuitable for future PCN's.

A distributed communication architecture based on the IEEE 802.6 MAN ( Metropolitan Area Network ) has been proposed for implementing the PCN. The proposed architecture of IEEE 802.6 MAN will be discussed in the next section, and all the analyses in this thesis are based on this architecture. Not surprisingly, MANs are expected to be a significant element in the evolution towards B-ISDN.

## **Section 1.2**

### **PCN Architecture Based on IEEE 802.6 MAN**

The proposed IEEE 802.6 MAN based distributed PCN architecture is shown in Figure 1.

The MAN has two 155 Mbps buses, one for each direction of transmission. Several base stations ( BSs ), Local Area Networks ( LANs ), and Private Branch Exchange ( PBXs ) are connected to a MAN node and are subnetworks of the MAN. Each MAN node functions as an interface to enable subnetworks to access the local MAN. Consequently, information between subnetworks can be exchanged within the local MAN, which then serves as a distributed switch. Interconnection of adjacent MANs enables coverage of an entire Metropolitan Area ( MA ) as shown in Figure 2. Initially, the relatively few MANs deployed in the MA would be

interconnected via point-to-point homogeneous bridges. As the deployed number of MANs

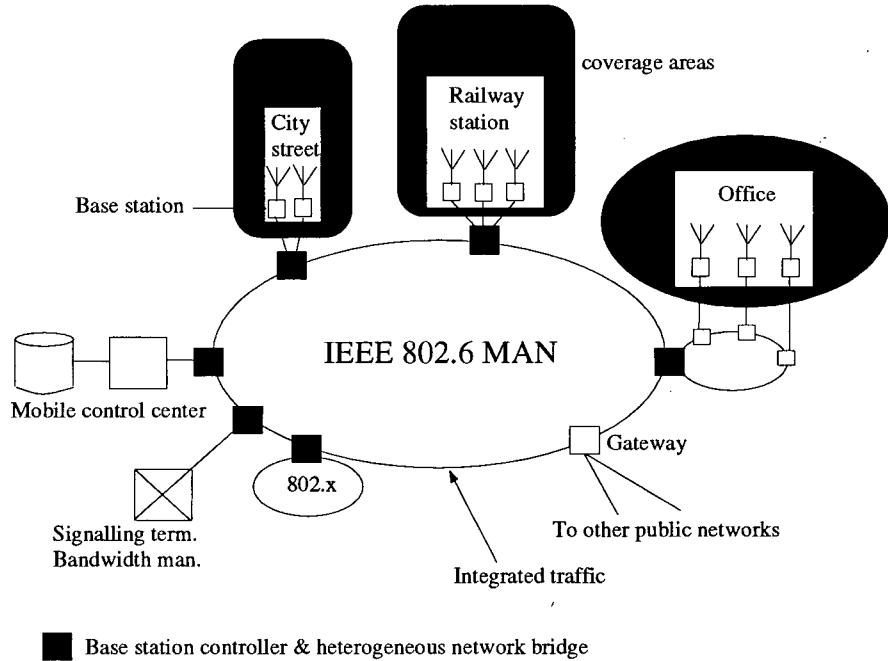


Figure 1. IEEE 802.6 MAN based distributed PCN architecture

rises, network manageability would necessitate the use of a centralized higher-speed multiport bridge — a B-ISDN switch at a Central Office ( CO ). Since the existing connection-oriented telecommunication networks will continue to exist for many years, an Interworking Unit ( IWU ) at the CO should have the functionality to support both ATM and Synchronous Transfer Mode ( STM ) switching fabrics. Additionally, a backbone MAN could interconnect these access MANs using homogeneous bridges to form a cluster ( see Figure 2 ). An IWU would then connect the backbone MAN(s) in a MA. If neighboring access MANs within a cluster are also bridged, call set up and inter-MAN handoff would be simplified, and reliability against bridge failure would be enhanced [1].

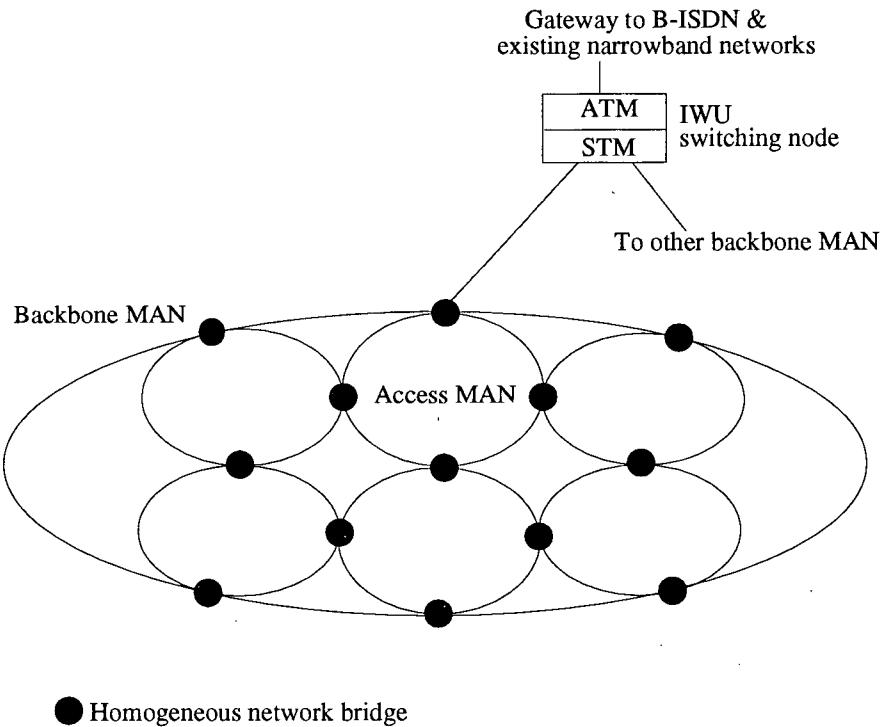


Figure 2. MAN interconnection within a backbone MAN [1]

The IEEE 802.6 MAN provides two modes of access control to the dual bus. These are Queue arbitrated ( QA ) and Pre-Arbitrated ( PA ), which use QA and PA slots for access, respectively. QA access is controlled by the Distributed Queueing protocol ( details of the protocol are provided in chapter 2 ) and would be used typically to provide non-isochronous services, such as data or graphics. PA access would be used to provide isochronous services, such as real-time voice or video [3]. Although allocation of PA slots involves some overhead, PA access provides a constant packet transmission delay at any time. On the other hand, QA access provides a variable packet transmission delay with a simpler access method. In this thesis, QA access is used to provide all services in order to achieve full statistical multiplexing.

The main advantage of using the 802.6 MAN as the base-site interconnection network is that not only can it readily interface to the current wireline public network facilities, but it

is also capable of implementing the Asynchronous Transfer Mode ( ATM ) for interworking between B-ISDNs and MAN-based PCNs. Since the cell format of ATM is similar to 802.6 MAN, the interworking between two networks is simplified and the transition to ATM will be relatively easy and inexpensive. In that case, the use of the 802.6 MAN as a transit network for ATM-based B-ISDN appears not only as a stepping stone in the evolution, but also as a more permanent solution [4].

### **Section 1.3 Outline of The Thesis**

In chapter 2, the IEEE 802.6 MAN with distributed queue dual bus ( DQDB ) operation is described. The DQDB protocol with QA access will be evaluated along with two slot reuse schemes, *Destination Release and Previous Slot Information*. These reuse schemes are solely for data packet transmission. Different types and levels of user traffic are considered.

Transmission of voice and data integrated traffic using QA access on a single MAN is described and analyzed in chapter 3. Voice and data packet waiting delay is evaluated for different user traffic levels and different speech encoding/decoding rate at voice stations.

In chapter 4, transmission delay of voice and data integrated traffic using QA access for Inter- and Intra-MAN will be evaluated. Various level of inter-MAN traffic, different user traffic levels and different speech encoding/decoding rates at voice stations is also considered in the analysis. For completeness, comparison of call setup delay between QA access isochronous traffic and PA access isochronous traffic is presented.

A summary of the main results is presented in Chapter 5, together with suggestions for further work.

# Chapter 2. Performance Analysis of IEEE 802.6 MAN

In order to understand the working of IEEE 802.6 MAN, the MAN architecture and protocol mechanisms are reviewed in this chapter. The general performance of the DQDB protocol using QA access with and without slot reuse schemes, including Destination Release and PSI, are compared using simulation models designed solely for data packet transmission.

## Section 2.1 The IEEE 802.6 MAN

### 2.1.1 The IEEE 802.6 MAN architecture

The DQDB medium access control ( MAC ) protocol has been standardized by the IEEE 802.6 committee as part of its MAN standard. The DQDB protocol is intended for use with a dual-bus configuration as illustrated in Figure 3.

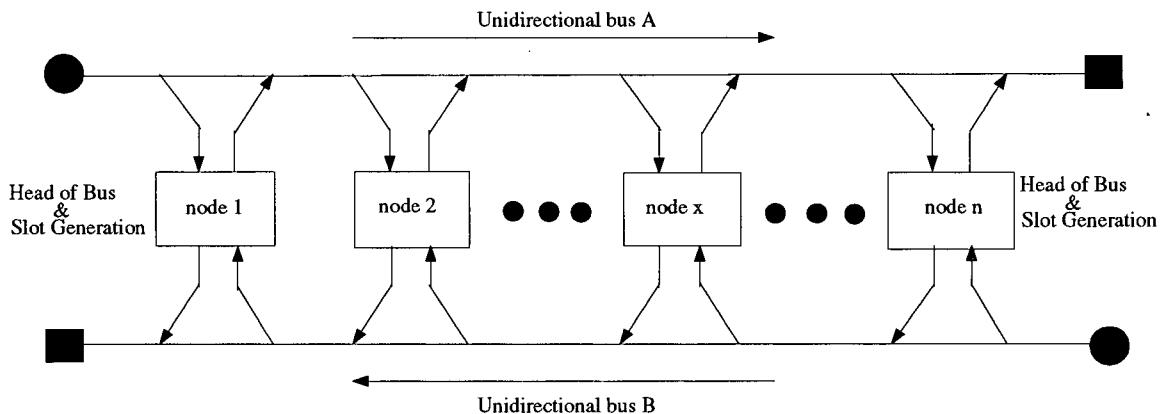


Figure 3. DQDB dual-bus configuration

As shown in Figure 3, each node is connected to two unidirectional buses which transmit in opposite directions. A node accesses the right-moving bus to send packets to nodes on its right and the left-moving bus for nodes on its left. Two head nodes are responsible for generating

slots on two buses. Transmission on each bus consists of a steady-stream of fixed-size slots with a length of 53 octets. The 53 octet length was chosen for compatibility with ATM. Nodes read and copy data from the slots and gain access to the bus by writing to empty slots. Streams of slots are controlled by a 125  $\mu$ sec-clock. Two head stations generate multiple slots to the shared medium every 125  $\mu$ sec; the number of slots generated per clock cycle depends on the physical data rate. Figure 4 shows the frame and slot format for the DQDB MAN.

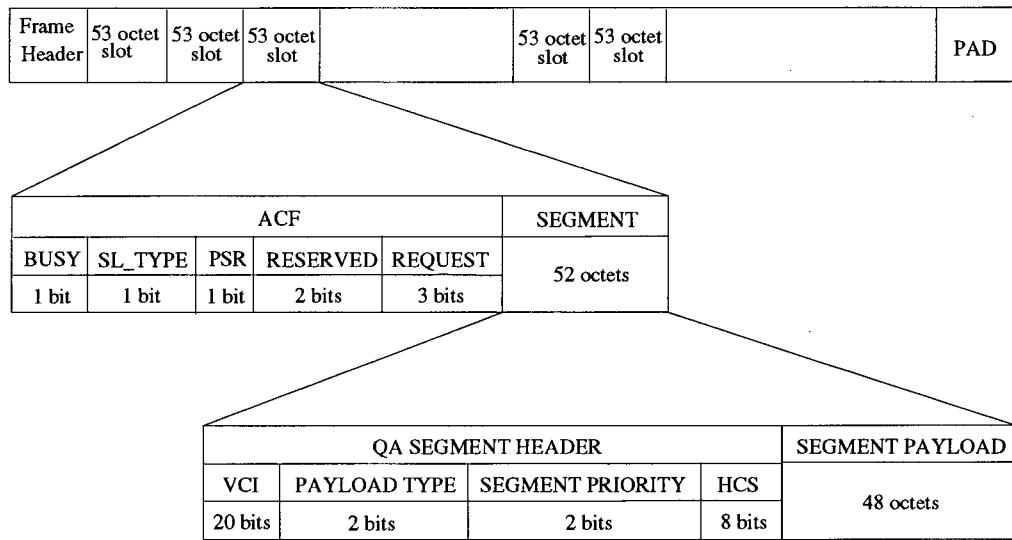


Figure 4. Frame and slot format in DQDB MAN

Each 53-octet slot contains a 1-octet Access Control Field ( ACF ) and a 52-octet segment. The ACF contains five fields that control access to slots. The BUSY bit indicates whether or not the slot contains information ( BUSY=1, NOT BUSY=0 ). The SL\_TYPE bit indicates whether the slot is a QA slot ( SL\_TYPE=0 ) or a Pre-Arbitrated ( PA ) slot ( SL\_TYPE=1 ). Table 1 shows the slot state with respect to the combinations of BUSY and SL\_TYPE. The PSR bit indicates whether the segment in the previous slot may be cleared ( PSR=1 ) or not ( PSR=0 ). The REQUEST field contains three REQ bits which are used in the operation of the three priority

level distributed queue access mechanism. The remaining 52 octets of the slot are divided into a 4-octet QA segment header and a 48-octet payload. The QA segment header contains four fields. The 20-bit Virtual Channel Identifier ( VCI ) provides a means to identify the virtual channel to which the QA segment belongs. A single VCI space is shared by all services. The VCI value corresponding to all bits being set to one is the default value of the connectionless MAC service provided by the MAC Convergence Function. All nodes shall use the default connectionless VCI for both transmission and reception of QA segments. The 2-bit Payload\_Type field indicates the nature of the data to be transferred. The 2-bit Segment\_Priority field is reserved for future use with Multiport Bridging. The 8-bit Header Check Sequence ( HCS ) field provides for detection of errors and correction of single-bit errors in the QA segment header [3].

BUSY	SL_TYPE	Slot State
0	0	Empty QA slot
0	1	Reserved
1	0	Busy QA slot
1	1	PA slot

Table 1. Slot Access Control Field coding [3]

Normally, PA slots are used to carry isochronous traffic ( e.g. voice, video ) while QA slots carry data traffic. For PA slots, some arbitration function must be used to allocate a dedicated sequence of slots for each stream of data and this function has not yet been standardized. In this thesis QA access only will be employed for all traffic types.

### **2.1.2 The distributed queue access protocol**

The 802.6 medium access control technique is the distributed-queue dual-bus ( DQDB ) protocol. The DQDB layer is independent of the physical layer. Therefore, a variety of DQDB

networks can be implemented using the same access layer but operating at different rates over different transmission system [5]. Three common transmission systems for DQDB MAN are as follows:

- ANSI DS3: transmits at 44.736 Mbps.
- ANSI SONET ( CCITT SDH ): transmits at 155.52 Mbps and above.
- CCITT G.703: transmits at 34.368 Mbps and 139.264 Mbps.

In this thesis, we will consider the ANSI SONET transmission system as the foundation of our DQDB protocol analysis. Distributed queueing is a media access protocol that controls the access of QA slots on the DQDB bus. The DQDB protocol is well suited for handling bursty traffic.

The operation of the DQDB protocol involves use of two control fields: the BUSY bit indicates whether or not a slot is used, and the REQUEST field indicates when a segment has been queued for access. Each node, by counting the number of requests it receives and unused slots that pass, can determine the number of segments queued ahead of it. Therefore, whenever a node has a segment for transmission, the node uses this count to determine its position in the distributed queue. This counting operation establishes a single queue across the subnetwork of segments queued for access to each bus [3]. The behavior of each node, refer to Figure 3, is summarized in Table 2 ( in this table, we are referring to A as the transmitting bus and B as the request bus ).

(a) Behavior of Node 1:

(i) At time when it is ready to issue the next QA slot on bus A

	No preceeding requests outstanding	One or more preceeding requests outstanding
Node 1 has no data to send	Issue a free QA slot	Issue a free QA slot and reduce preceeding requests by one
Node 1 has a QA segment to send	Insert data into the QA slot; any following requests become preceeding request	Issue a free QA slot and reduce preceeding requests by one

(ii) At time when Node 1 receives the next QA slot on bus B

	Incoming slot contains a request	Incoming slot does not contain a request
Node 1 has no data to send	Add 1 to preceeding requests	—
Node 1 has a segment to send	Add 1 to following requests	—

(b) Behavior of Node x:

(i) At time when it observes a free QA slot on bus A

	No preceeding request outstanding	One or more preceeding request outstanding
Node x has no data to send	Let free slot pass	Let free slot pass and reduce preceeding request by one
Node x has a QA segment to send and has issued a request on bus B	Inset data into the QA slot; any following requests become preceeding requests	Let free slot pass and reduce preceeding request by one

(ii) At time when Node x observes a QA slot on bus B

	Incoming slot contains a request	Incoming slot does not contain a request
Node x does not have an outstanding request	Add 1 to preceeding requests	—
Node x has a QA segment to send and has issued a request for that segment	Add 1 to following requests	—
Node x has a QA segment to send and has not issued a request for that segment	Add 1 to preceeding requests	Insert request into the passing slot

Table 2. Behavior of nodes in Figure 3 [5]

## Section 2.2

### Slot Reuse Schemes

In order to investigate the efficiency of QA access to the DQDB MAN, two slot reuse schemes, *Destination Release and Previous Slot Information (PSI)*, together with the absence of slot reuse are compared.

#### 2.2.1 Destination Release

In Destination Release slot reuse, every destination station releases slots for further reuse by downstream stations. This scheme offers maximum possible usage of capacity. However, the complexity of the receiver hardware is increased, the access protocol becomes more complicated and the delay at every station increases.

#### 2.2.2 Previous Slot Information (PSI)

PSI, introduced by Sharon and Segall [6], is a more adaptive scheme and less efficient than Destination Release, but will perform slot reuse without delay and without increasing the

hardware complexity. PSI utilizes the fact that a station is always listening to the bus in order to find messages destined to itself, and that every station reads the destination of every passing slot. PSI is based on a comparison of the destination addresses of consecutive slots. The scheme employs a reuse flag R in the slot's ACF ( Access Control Field ). Flag R is set by a transmitting node in any slot whose destination is beyond the destination of the packet that was received and retransmitted in the previous slot. If a busy slot, B=1, is received, a station  $\alpha$  performs reuse if both the destination of the previous received slot is located before  $\alpha$  and R=0. The setting R=0 implies that the destination of the current slot is located before the destination of the previous slot. If station  $\alpha$  reuses the slot and if the destination of the current slot is beyond the destination of the previous slot, R is set to 1.

A simple example of PSI is given in Table 3. This example indicates the value of the Destination Address ( DA ) and the R bit of four consecutive slots travelling to stations connected to the bus in the order  $\alpha, \beta, \gamma, \delta$ ; station  $\alpha$  transmits to station  $\beta$ , and station  $\gamma$  transmits to  $\delta$  starting with the second slot.

Slot	0		1		2		3	
	Stns	DA	R bit	DA	R bit	DA	R bit	DA
$\alpha$	NULL	0	$\beta$	1	$\beta$	0	$\beta$	0
$\beta$	NULL	0	$\beta$	1	$\beta$	0	$\beta$	0
$\gamma$	NULL	0	$\beta$	1	$\delta$	1	$\delta$	0
$\delta$	NULL	0	$\beta$	1	$\delta$	1	$\delta$	0

Table 3. Example of PSI [6]

Since slot 0 remains empty, its DA is represented as NULL and its R bit remains 0. Station  $\alpha$  sets the R bit of slot 1 since the NULL address of slot 0 is considered upstream to all stations, and  $\alpha$  transmits in slot 1 to  $\beta$ . Since the R bit is set, slot 1 will never be reused. When  $\alpha$

transmits in slot 2, it does not set the R bit because it does not transmit to a station beyond the destination of slot 1. Station  $\gamma$  knows that slot 1 is destined to  $\beta$  which is before itself, and so observing that the R bit of slot 2 is not set, it can reuse slot 2 to transmit data to  $\delta$ . But it sets the R bit because it transmits to station  $\delta$  which is beyond  $\beta$ . Thus slot 2 can never be reused again. Finally, both stations  $\alpha$  and  $\gamma$  transmit in slot 3 and neither transmits to a station beyond their last transmitted packet destinations,  $\beta$  and  $\delta$  respectively, so R bit is not set.

## **Section 2.3**

### **Modelling and Simulation of the DQDB Protocol**

#### **2.3.1 DQDB protocol model**

The DQDB protocol serves a distributed collection of queues; usual service order in each queue is first-in, first-out ( FIFO ). Since the access control mechanisms of two buses are identical we discuss the access control of bus A only. Consider that node  $\alpha$  has a data packet to send to node  $\beta$  which is downstream of node  $\alpha$ . Node  $\alpha$  sends a request on Bus B to the head node of Bus A to obtain an empty slot on Bus A for packet transmission.

At each node, a queue of each priority level is formed and two counters, ReQuest Counter ( RQC ) and CountDown Counter ( CDC ) for each queue are employed to track the status of the distributed queue of Bus A. The RQC is used to track the number of slots requested by the downstream stations while the CDC is used for transmission queueing. For example, when a station sees the request bit of a slot set on Bus B, then station increments its RQC by one. After a station sends its request, the value of its RQC is transferred to its CDC and its RQC is reset to zero. The CDC is now decremented for each passing idle slot on Bus A until it reaches zero, when the station is then allowed to transmit in the next idle slot. While awaiting transmission on Bus A, the RQC continues to increment for any new request made by downstream stations

on Bus B. Evidently, each station will only have one and only one packet of information in the system queue at any time, because a new request is not allowed until the current data packet has been transmitted.

### **2.3.2 Simulation models**

In this preliminary analysis, we only consider the transmission of one-way single-priority level traffic ( i.e. data traffic ) on a DQDB MAN. Three models are developed for performance comparison: (1) standard DQDB MAN, (2) DQDB MAN with PSI slot reuse, and (3) DQDB MAN with Destination Release slot reuse. All simulation models are written in **SIMSCRIPT II.5** [7].

The simulation model for standard DQDB MAN is built exactly as described in the last section. However, modifications are needed for PSI-DQDB MAN and Destination Release-DQDB MAN models.

In addition to the standard DQDB MAN model, the PSI-DQDB MAN model introduces a R bit and a reuse bit in the ACF in each slot, and two memory units to keep track of the previous slot address received ( PSAR ) and the previous slot address transmitted ( PSAT ) in each node. The two memories, PSAR and PSAT, are used to control the setting of R bit in the current slot and determine whether or not the current slot can be reused. Each node continuously updates the two memories as each slot passes that node. If a node has not transmitted in the previous slot, its PSAR will equal its PSAT. PSAR may not be equal PSAT if a node has transmitted in an empty slot or has reused a slot. The R bit may be set only if a node transmits either in an empty slot or by reusing a slot, as described earlier in section 2.2.2.

Implementing the Destination Release DQDB MAN model is simpler. The BUSY bit in the ACF is reset to 0 once a node copies the information from a QA slot which is addressed to that

node. Apart from these differences, the following assumptions are common to all three models:

- i. The data rate on each bus is 155 Mbps
- ii. Slot length is 53 octets (424 bits)
- iii. Slot transmission time is  $2.73 \mu\text{s}$
- iv. Nodes are evenly distributed along the bus
- v. Successive nodes are separated by 3 slots
- vi. The number of buffers available to each node is 30
- vii. All messages are 1 slot in duration
- viii. Inter-arrival time of data packets at each node is exponentially distributed
- ix. Each node has the same Poisson arrival rate
- x. The number of nodes, excluding two head stations, is 100 ( 50 nodes per bus )
- xi. Simulation time equals to 50,000 slots

Since PSI performance is highly dependent on the traffic destination distribution, two bursty data traffic models have been simulated for completeness. The first one is the uniform destination traffic distribution; the second is the geometric destination traffic distribution. The probability function,  $P(x)$ , of a geometric distribution is as follows:

$$P(x) = \theta(1 - \theta)^{x-1} \quad x = 1, 2, 3, \dots$$

where  $\theta$  is the degree of traffic locality [8]. We selected  $\theta$  to be 0.9, to simulate traffic with a high degree of locality.

### **2.3.3 Simulation results**

Simulation results plotted in Figures 5 and 6 show mean packet delay versus throughput. Mean packet delay is defined as the elapsed time from when a packet enters the station buffer until that packet is successfully transmitted. Figure 5 shows the uniform destination distribution results while Figure 6 shows the geometric destination distribution results.

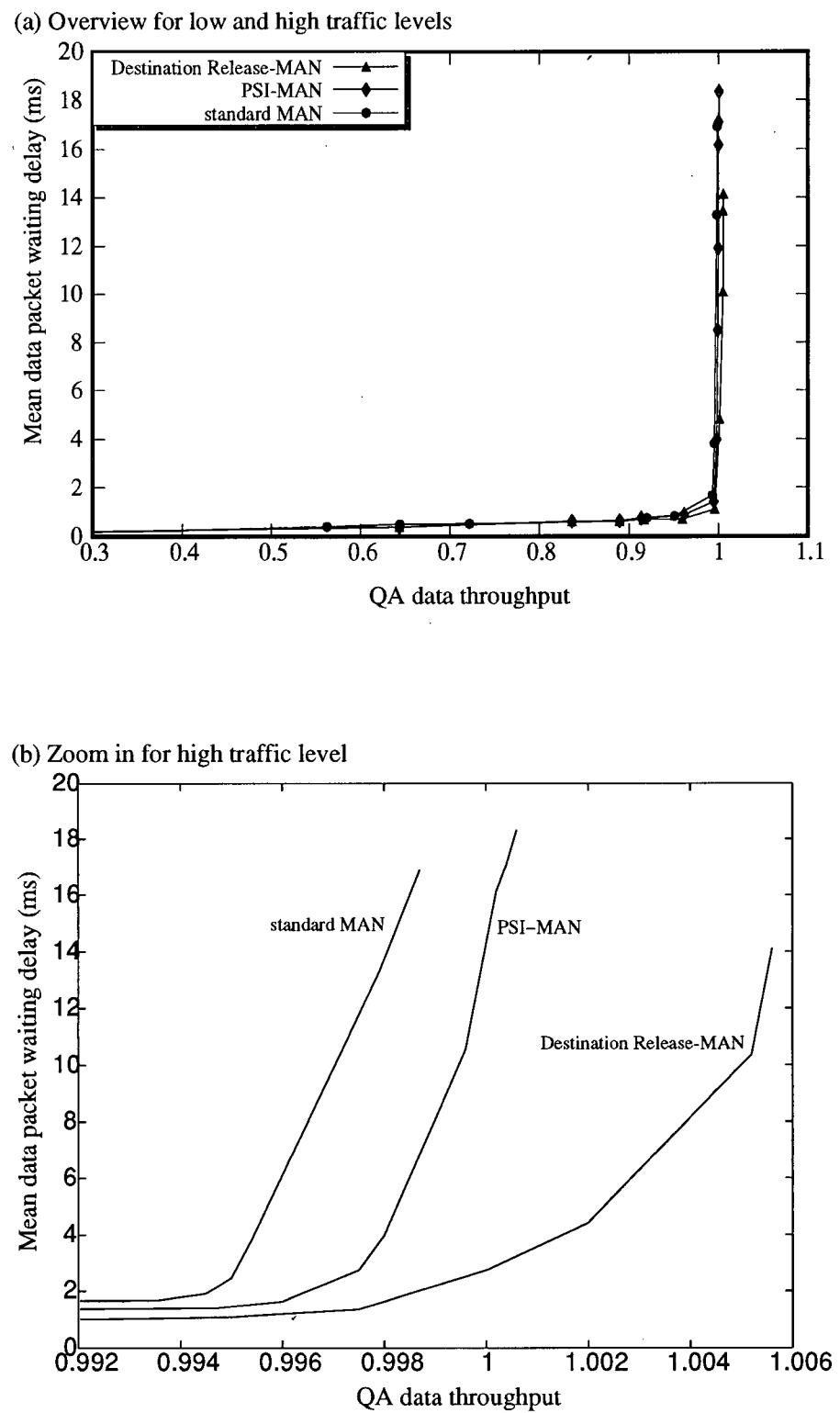


Figure 5. Uniform destination distribution for data packets in a DQDB MAN

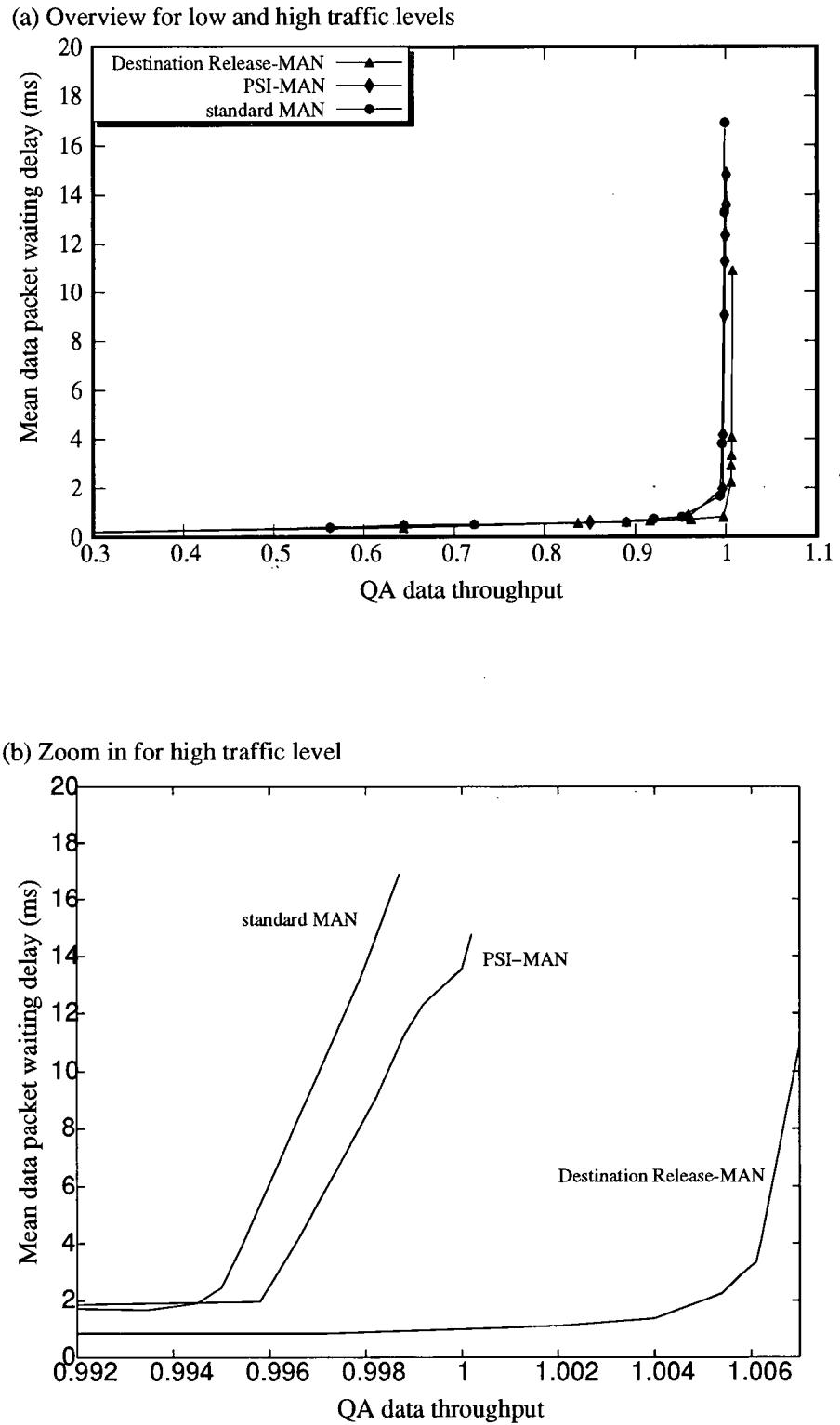


Figure 6. Geometric destination distribution for data packets in a DQDB MAN

From Figures 5a and 6a, one observes that all three simulation models have approximately the same mean packet delay at traffic levels below 0.95, for both uniform and geometric traffic. However, under heavy traffic ( above 0.95 ), Figures 5b and 6b show that the PSI slot reuse scheme shows some improvement in the packet waiting delay with respect to the QA throughput. The Destination Release slot reuse scheme shows even a better improvement in the packet waiting delay with respect to the QA data throughput than the PSI.

At a given delay value, the two slot reuse schemes do not result in much improvement in throughput ( i.e. less than 10 percent ). The PSI slot reuse scheme shows less improvement in geometric destination traffic than in uniform destination traffic , once a station transmitted its current packet, say in slot 1, slot 2 is either occupied or idle while passing the station. When the address memories are updated prior to arriving at slot 2, the updated addresses have a very high probability at being less than the station's address because the upstream stations' packet addresses are highly localized. Therefore the probability of having a set reuse bit, R, in the slot where the station transmitted its next packet is increased considerably for geometrically distributed destination traffic. This result occurs because a station is allowed only one packet in the DQDB system queue at any time. If a station were not limited to having at most one packet in the system queue, the performance of PSI in geometrically distributed destination traffic would be better than in uniformly distributed destination traffic.

The results in Figures 5 and 6 suggest that it is not worthwhile to implement either slot reuse schemes DQDB MAN, since either one complicates the implementation and could introduce processing delay.

# **Chapter 3. Delay Analysis of QA Access for Integrated Intra-MAN Voice and Data Traffic**

Voice packets have different transmission requirements from data packets. To preserve the integrity of a conversation, voice packets must be delivered within a maximum delay time comparable to the period between generation of consecutive voice packets. Some loss of voice packet is tolerable. On the other hand, data packets can undergo longer and variable delay but no loss [9]. Base on these properties of voice and data packet transmission requirements, a system model is built for voice and data integrated traffic in a IEEE 802.6 MAN. The model allows for use of different values of key parameters such as speech digitizing rate, voice stations activation time and bandwidth available for QA access.

## **Section 3.1 Voice-Data Integration**

It is assumed that a fixed number of voice calls is activated at each node for each simulation, and these calls will remain active throughout the simulation. Each voice call independently alternates between a talkspurt period and a silence period; these are independent exponentially distributed random variables with mean  $1/\alpha=1.5$  s and  $1/\mu=2.25$  s, respectively. Therefore, the probability,  $P(t \rightarrow s)$ , that a voice call which changes from talkspurt to silence in  $f$  seconds is

$$P(t \rightarrow s) = 1 - e^{-\alpha f}$$

and from silence to talkspurt is

$$P(s \rightarrow t) = 1 - e^{-\mu f}$$

Figure 7 shows the transition probabilities of the associated two-state Markovian model.

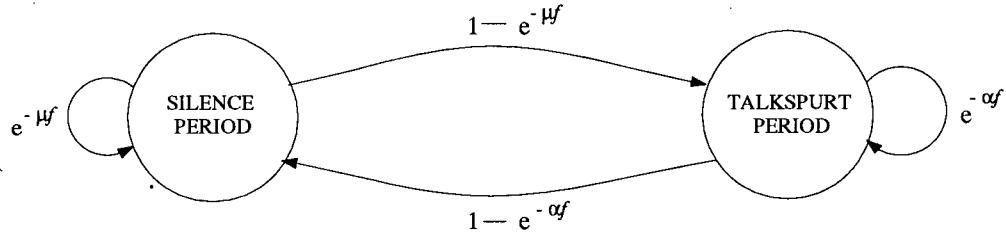


Figure 7. Active voice station model [9]

One sees that the transition probabilities of the above two-state Markovian model are highly depended on  $f$ . Figure 7 shows that the transitional probabilities increase exponentially as the time variable  $f$  increases and conversely for the non-transitional probabilities. The longer the time a call remains in one state, the greater the tendency to change. Since voice packets can tolerate losses, we assume that if a voice packet cannot be transmitted within a constant time period  $t$ , the voice packet will be discarded. Our choice of  $t=5.5$  ms is to ensure compatibility with the air interface which has 5.5 ms frame length.

### 3.1.1 Voice clipping

Voice clipping occurs only if the number of ready-to-transmit voice stations within a frame period exceeds the number of slots/frame . For each active voice call in the model, we assume that talkspurt and silence periods alternate independently with exponential probability density functions  $f(\tau)$  and  $g(\tau)$  respectively, where  $\alpha$  and  $\mu$  are constants and

$$f(\tau) = \alpha e^{-\alpha\tau}$$

$$g(\tau) = \mu e^{-\mu\tau}$$

From these two probability density functions, we can determine the probability,  $P(t)$ , that an active voice station is in talkspurt mode at any arbitrary time instant. With the help of  $P(t)$ , we can then determine the voice clipping probability for a specific transmission system.

Since functions  $f$  and  $g$  are independent, the joint probability density  $h(\tau_1, \tau_2) = f(\tau_1)g(\tau_2)$  and

$$\begin{aligned}
 P(t) &= P(\tau_1 > \tau_2) = \int_0^\infty \int_{\tau_2}^\infty h(\tau_1, \tau_2) d\tau_1 d\tau_2 \\
 &= \int_0^\infty \int_{\tau_2}^\infty \alpha \mu e^{-(\alpha \tau_1 + \mu \tau_2)} d\tau_1 d\tau_2 \\
 &= \alpha \mu \int_0^\infty e^{-\mu \tau_2} \left[ \int_{\tau_2}^\infty e^{-\alpha \tau_1} d\tau_1 \right] d\tau_2 \\
 &= \mu \int_0^\infty e^{-(\mu + \alpha) \tau_2} d\tau_2 \\
 &= \frac{\mu}{\alpha + \mu}
 \end{aligned}$$

Voice clipping occurs whenever more stations are in talkspurt mode than there are QA slots available. As a result, the voice clipping probability,  $P_c$ , for  $N$  active voice stations and  $k$ , slots/frame, available is as follows:

$$P_c = \begin{cases} \sum_{T=k+1}^N \binom{N}{T} P(t)^T (1 - P(t))^{N-T} & , N > k \\ 0 & , N \leq k \end{cases}$$

where  $\binom{N}{T} P(t)^T (1 - P(t))^{N-T}$  is the probability of having  $T$  out of  $N$  number of voice stations in talkspurt mode.

Figure 8 shows the voice clipping probability versus the number of ready-to-transmit voice stations. For the values of constants listed earlier,  $P(t)=0.4$ , in which case  $k=2 \times 2010$  slots/frame/2 buses for a 5.5 ms frame length of 424 bit per slot at 155 Mbps transmission rate.

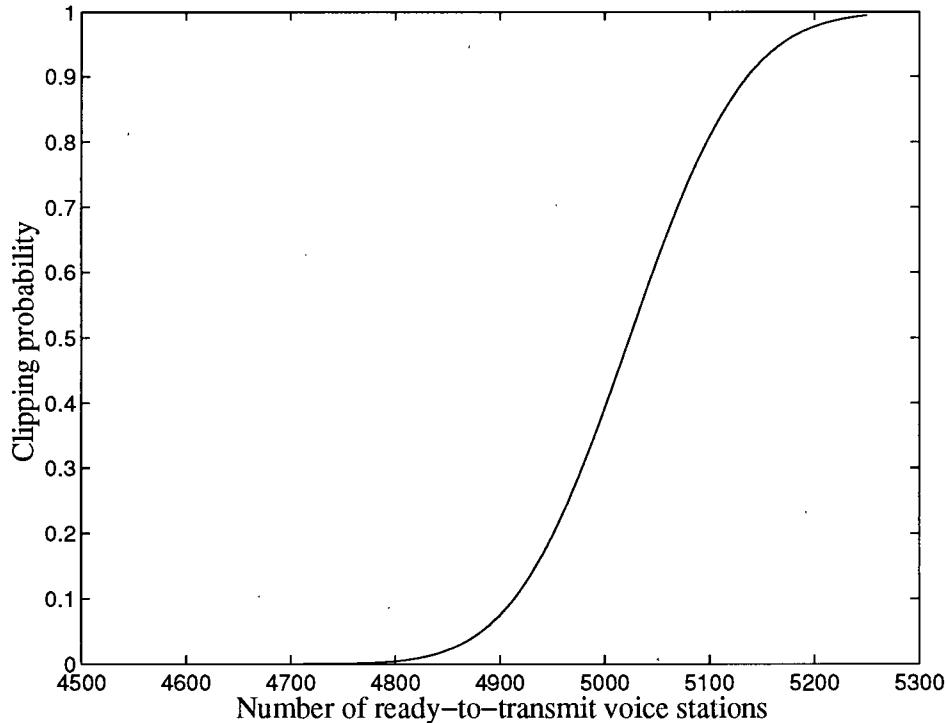


Figure 8. Voice clipping probability vs. number of ready-to-transmit voice stations

## Section 3.2 Simulation Models

Voice and data traffic are assigned access levels priority 1 and 0, respectively. Priority 2, the highest level, is reserved for control signalling. Three priority levels are established by operating three queues, one for each level. Within each level, the performance characteristics described above are maintained. Waiting packets acquire access as soon as capacity becomes available; however, access priority is always given to packets in higher level queues.

Since voice packets are discarded if not be transmitted within 5.5 ms, the implementation of a DQDB counter mechanism will be slightly more complicated than for data only networks. According to the counters mechanism discussed in section 2.3.1, each node tracks the number of preceding and following requests for its downstream stations in each priority level. Therefore, if a voice packet at one particular node is discarded, its priority 1 countdown counter and request transmission for the next packet cannot just simply be reset. For example, let us consider two nodes,  $\alpha$  and  $\beta$ , with  $\beta$  located downstream from  $\alpha$ . Assume that there is only one voice packet in node  $\alpha$  and several voice packets in node  $\beta$  awaiting transmission. Consider node  $\alpha$  at the instant when no following request (  $RQC=0$  ) is received for voice traffic from downstream stations and the priority level 1 CDC is reset because its last voice packet is discarded; then its next transmission request will be the request for the data packet which located at the head of the data packet queue, since node  $\alpha$  does not have any voice packet awaiting transmission. Possibly, node  $\alpha$  may transmit its requested data packet while node  $\beta$  still has voice packets waiting for transmission since the priority 1 level CDC is reset and RQC is equal to zero in node  $\alpha$ . Consequently, the priority accessing scheme is violated. In order to preserve priority accessing, preceding requests, if any, will be added to the following requests ( i.e. the value of CDC is added to RQC ) for each discarded packet in each priority level before the next packet makes request for transmission. Hence, the problem of the above example will be solved.

In our model, all voice call destinations are predetermined and are uniformly distributed among destination nodes. Since voice packet delay is our primary concern, all voice call connections are assumed to be established at the beginning of any simulation. Moreover, each voice call will remain active throughout each simulation run. Except for the differences in parameter values, the following assumptions apply to all simulation models. These models are

developed for two way traffic for more realistic analysis; packets are sent on bus A or B, as appropriate:

- i. The transmission rate on each bus is 155 Mbps
- ii. Slot length is 53 octets ( 424 bits )
- iii. Slot transmission time is  $2.73 \mu s$
- iv. Nodes are evenly distributed along the bus
- v. Successive nodes are separated by 3 slots
- vi. Twenty buffers are available at each node for data packets
- vii. All packet lengths, including voice and data, equal 1 slot in duration
- viii. Inter-arrival time of data packet at each node is exponentially distributed
- ix. Each node has the same Poisson arrival rate
- x. The number of nodes including two head stations equals 50
- xi. Each node has the same number of active voice calls
- xii. Priority is given to voice packets
- xiii. Any voice packet that cannot be transmitted within 5.5 ms will be discarded

Simulation models are written in **SIMSCRIPT II.5** [7]. Statistics are compiled after an initial period of 10,000 slots. Analysis is performed for various values of the following parameters: (1) speech digitizing rate of voice stations, (2) percentage of system bandwidth available for QA access, and (3) activation time for voice calls.

## **Section 3.3**

### **Simulation Results**

#### **3.3.1 Simultaneous voice calls arrival without silence deletion**

The simulation model used here is based on the assumption that all voice calls arrive simultaneously at the beginning of the frame and remain in the talkspurt period throughout the simulation ( without silence deletion ). The voice coding rate is 64 kbps; thus each voice station generates a voice packet every 6.0 ms.

Figure 9 shows the mean waiting delay<sup>1</sup> for a voice packet versus the QA voice throughput<sup>2</sup>. Since all voice packets enter MAN nodes at the same time, an increase in M, the number of voice packets, transmitted results initially in a proportional increase in the mean waiting delay. As M increases further, the mean voice packet waiting delay remains at 3 ms because further increase in the number of voice calls ( voice packets ) will only increase the number of discarded voice packets; these are packets cannot be transmitted within the 5.5 ms time limit as shown in Figure 10. There will not be any significant change in either the waiting delay or throughput as M increases beyond some limit. Figure 10 shows that the number of voice calls that the network can accommodate is about 3600 within an acceptable value of 2% discarded voice packets under the assumption of simultaneous voice calls arrival without silence deletion.

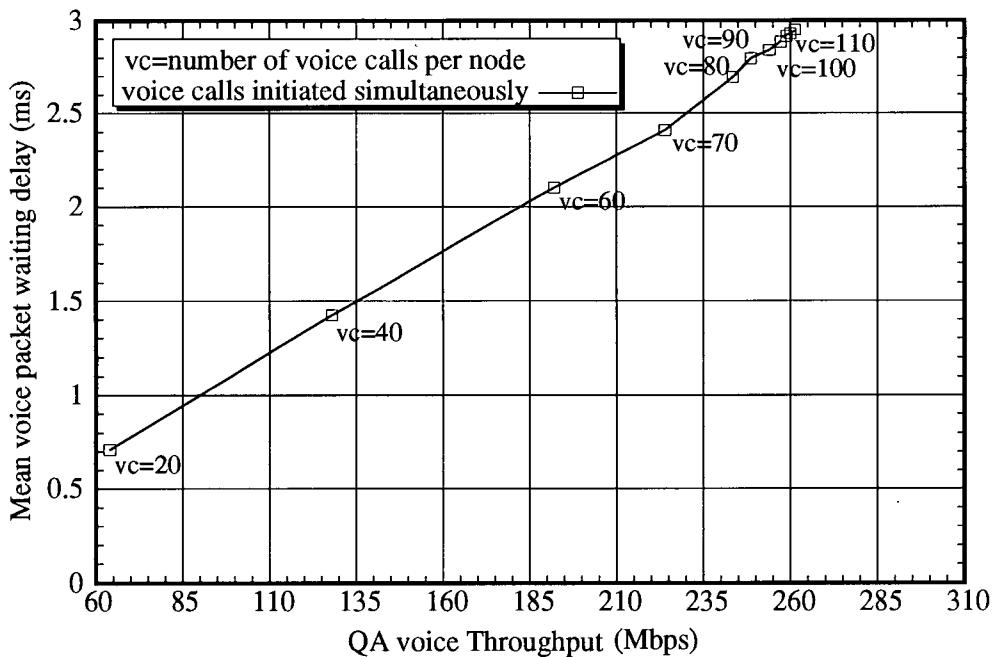


Figure 9. Mean voice packet waiting delay vs. QA voice Throughput

<sup>1</sup> waiting delay is calculated from the time a packet enter a node until it is transmitted.

<sup>2</sup> the voice throughput is calculated only for the segment payload excluding the overheads.

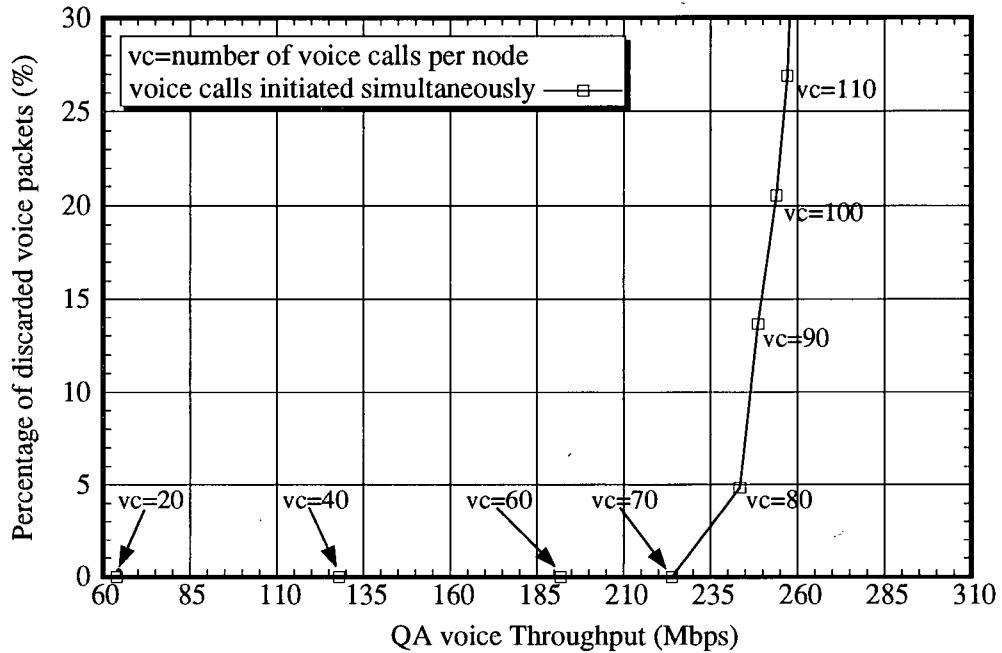


Figure 10. Percentage of voice packets discarded vs. QA voice Throughput

Figure 11 shows the mean data packet waiting delay versus the QA data throughput for different voice throughput levels. Voice calls are initiated simultaneously at the beginning of the simulation. As expected, data packet waiting delay increases as the network's voice throughput increases for the same data throughput.

Figure 12 shows a typical example of the mean waiting delay of data and voice packets with respect to the nodes' locations. Node 1 and 50 are the head of bus (HB). One sees that both voice and data packet mean waiting delay varies among the nodes. Nodes located near HBs have less delay than nodes located near the buses midpoints.

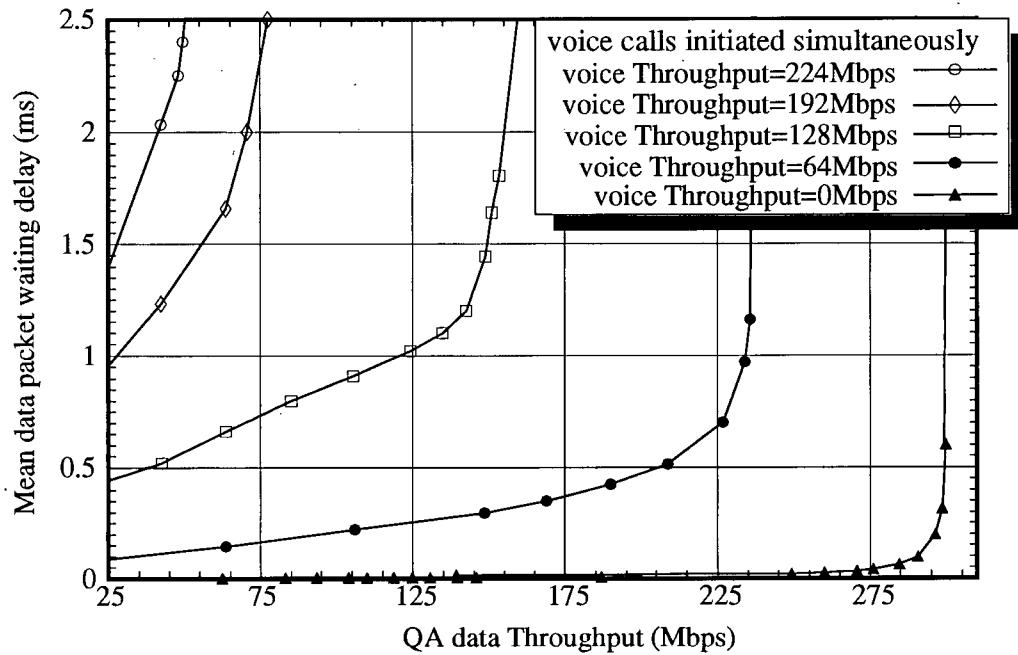


Figure 11. Mean data packet waiting delay vs. QA data Throughput for various voice throughput

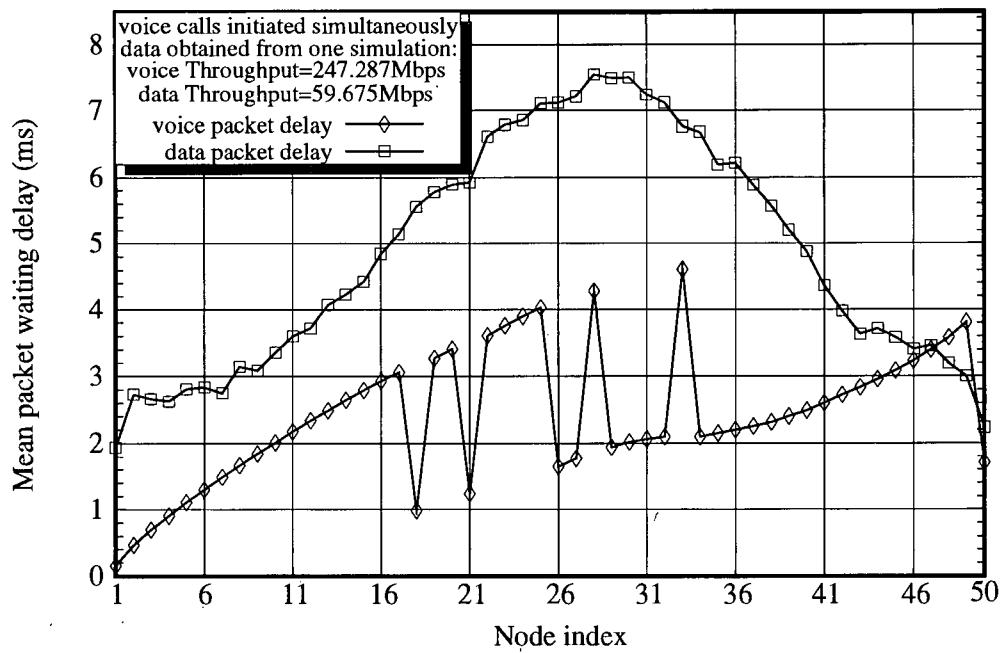


Figure 12. Mean packets waiting delay vs. nodes locations

Slots are generated from HBs travel along the buses unidirectionally; nodes located near HBs ( upstream nodes ) therefore will have a better chance for bus access than will nodes located at the bus midpoint.

### **3.3.2 Exponentially distributed voice call arrivals without silence deletion**

The simulation model used here is based on the same assumptions as in the previous section except that the voice call arrival times are Poisson. Time separation occurs between successive voice calls, and improvements in both voice and data waiting delay are expected as a result.

Figure 13 shows the mean voice packet waiting delay versus the QA voice throughput for both simultaneous and exponential voice call arrival models. The Poisson arrivals result in large improvements in mean waiting delay relative to the delay for simultaneous arrivals. Also, the QA voice throughput achieved under Poisson arrivals is higher than the simultaneous voice call arrival model because more voice packets are transmitted within the 5.5 ms time limit as the number of voice calls increases. The mean voice packet waiting delay limits at 5.5 ms under Poisson arrivals, as expected.

Figure 14 shows the percent discarded voice packets versus throughput. One sees that the number of voice calls that the network can accommodate is approximately 4200 within an acceptable 2% discarded rate. Under Poisson arrivals, the network can accommodate approximately 600 more voice calls than with the simultaneous call arrival model.

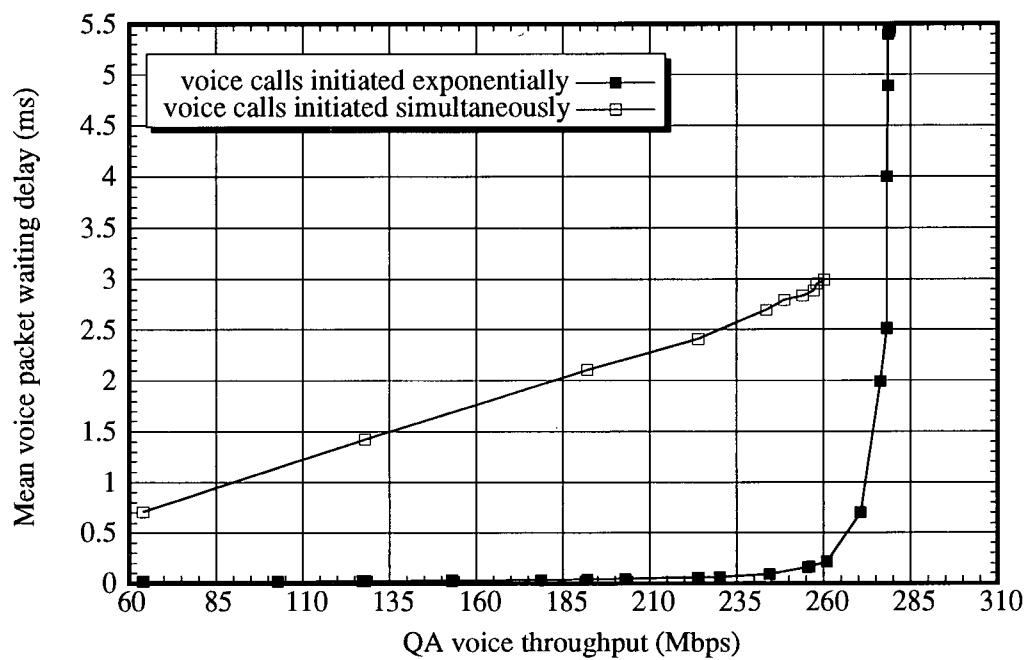


Figure 13. Mean voice packet waiting delay vs. QA voice Throughput

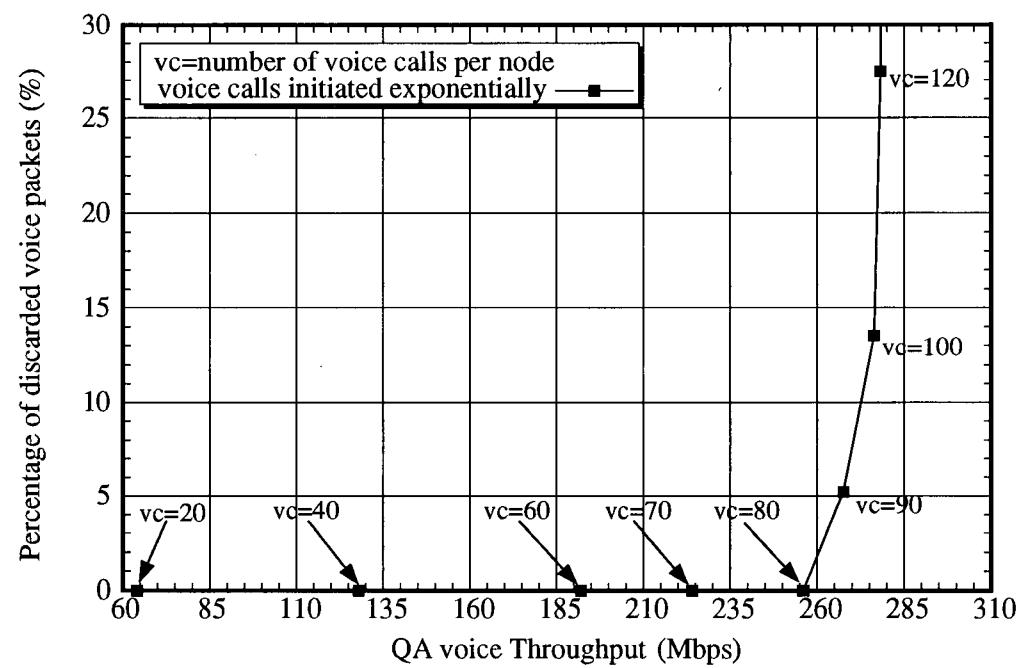


Figure 14. Percentage of discarded voice packet vs. QA voice Throughput

Figure 15 shows the mean data packet waiting delay versus data throughput for different voice throughput levels for both simultaneous and exponential voice call arrival models. The mean waiting delay is lower for the exponentially distributed voice call arrivals. However, both models result in the same ( asymptotic ) maximum data throughput levels.

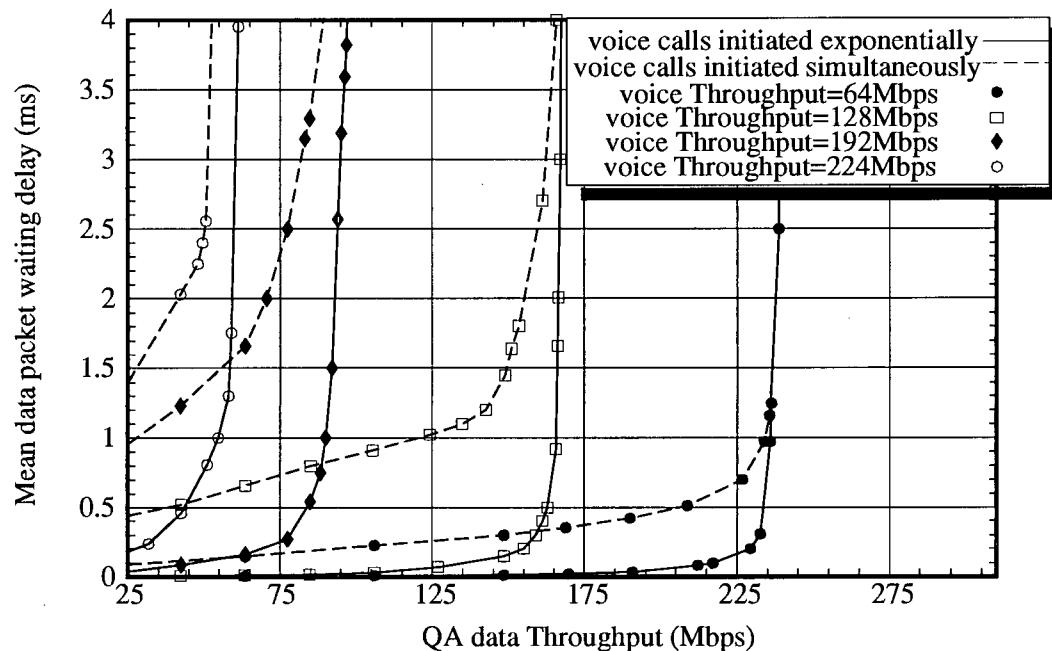


Figure 15. Mean data packet waiting delay vs. QA data Throughput for various voice Throughput

Figure 16 shows that the mean packet waiting delay varies among nodes. Upstream nodes have more rapid access to the bus than do nodes located in the middle of the bus.

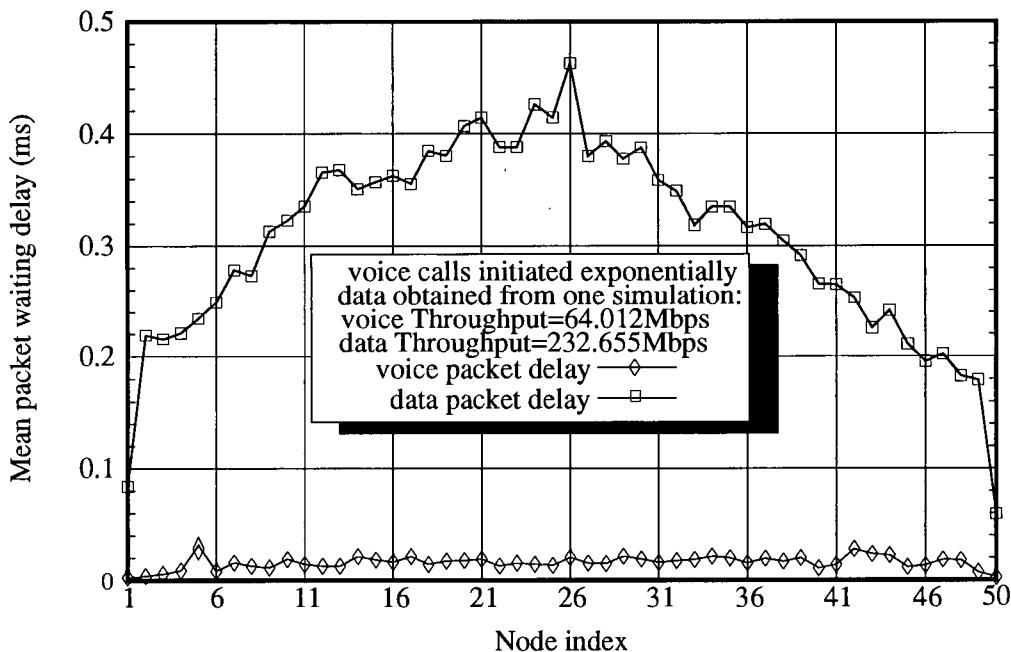


Figure 16. Mean packet waiting delay vs. node location

The results of the simultaneous and exponential voice call arrivals without silence deletion are summarized in Table 4 below.

Number of voice calls in the network	QA voice throughput achieved (Mbps)		Mean voice packet waiting delay (ms)		Percentage of discarded voice packets (%)	
	exponential voice calls arrival	simultaneous voice calls arrival	exponential voice calls arrival	simultaneous voice calls arrival	exponential voice calls arrival	simultaneous voice calls arrival
1000	64.012	64.012	0.018	0.711	0	0
2000	127.996	127.996	0.022	1.424	0	0
3000	192	192	0.041	2.102	0	0
3500	223.986	224.014	0.057	2.409	0	0
4000	255.992	243.527	0.163	2.694	0	4.86
4500	267.688	248.608	0.546	2.798	5.22	13.64
5000	276.375	254.087	1.992	2.886	13.55	20.55

Table 4. Results summary of simultaneous and exponential voice calls arrival without silence detection

The results of this section shows clearly the need for simulations which include the effects of random call arrivals.

### **3.3.3 Exponentially distributed voice call arrivals with silence deletion**

In this section, results using a simulation model with Poisson voice call arrivals with silence deletion are presented, to show how silence deletion affects networking performance. In the previous two models, all voice calls remained active and in talkspurt mode throughout the simulations. Now we assume that all voice calls are remain active throughout the simulation but that each voice call alternates independently between a talkspurt and a silence period, with exponentially distributed random durations of means 1.5 s and 2.25 s, respectively. Voice packets are generated every 6.0 ms ( voice coding rate=64 kbps ) from those voice calls in a talkspurt period. As shown previously in this chapter, the probability that an active voice call is in talkspurt mode at an arbitrary time instant equals 0.4; therefore, we assume that 40% of all voice calls are in talkspurt mode and that the remaining 60% are in silence mode at the beginning of each simulation. Each voice call alternates between silence and talkspurt according to time intervals randomly chosen from two independent exponentially distributed random distributions.

Silence deletion prevents a voice call from generating a packet for a specific period of time and will not have any effect on the mean packet delay for any given QA voice ( talkspurt ) throughput value. Consequently, the mean voice and data packet waiting delays are identical to those obtained earlier without silence deletion. Table 5 shows some examples of the number of voice calls needed to achieve listed QA voice throughputs with their respective mean voice packet waiting delays for exponentially distributed voice call arrivals with and without silence detection.

QA voice throughput (Mbps)	Number of voice calls needed to achieve the throughput		Mean voice packet waiting delay (ms)	
	with silence deletion	without silence deletion	with silence deletion	without silence deletion
127	5000	2000	0.022	0.022
192	7500	3000	0.039	0.041
256	10000	4000	0.159	0.163

Table 5. Comparison between models with and without silence deletion

It is almost impossible to achieve an exact match on the QA voice throughputs for a given number of voice calls for the two models used. The matching between the QA voice throughput and corresponding number of voice calls in Table 5 is a good approximation. The mean waiting delays for two cases are almost equal. Therefore, we can conclude that silence deletion will not affect the mean voice and data packet waiting delay with respect to QA voice throughput.

On the contrary, the number of voice calls that the network can accommodate with silence deletion must be larger than the one without silence deletion. The network can accommodate approximately 10500 voice calls with an acceptable 2% discarded voice packets with silence deletion, as shown in Figure 17. This is more than double in comparison to the case without silence deletion.

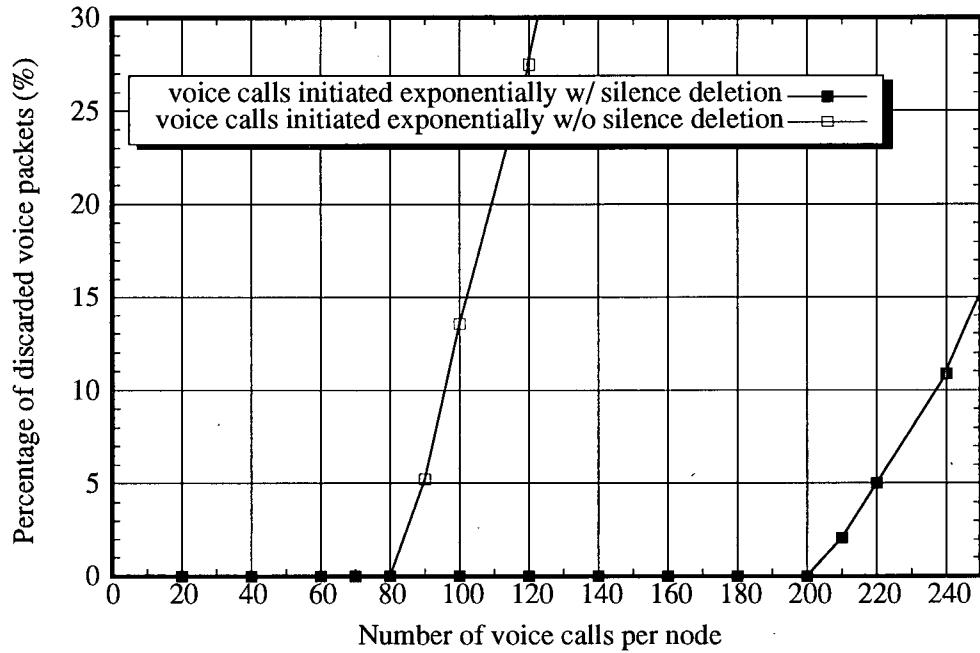


Figure 17. Percentage of discarded voice packets vs. number of voice calls per node

### 3.3.4 Exponentially distributed voice calls arrival for different voice coding rate with silence deletion

In this section, we consider the effect of voice coding rate on voice and data packet delay in the network. The simulation models are basically the same as the one in previous section. Each node has the same number of voice calls with Poisson arrivals. All voice calls are assumed to be active throughout the simulations and they alternate between talkspurt and silence period independently with exponentially distributed random durations of means 1.5 s and 2.25 s, respectively. As before, 40% of the randomly chosen voice calls are in talkspurt and the other 60% in silence at the beginning of each simulation.

However, the time between two successively generated voice packets at a voice station is not fixed at 6.0 ms ( 64 kbps ) because we are varying the voice coding rate of the voice station. Therefore, the time between two successive voice packets in a voice station is 6.0 ms for 64

kbps voice coding rate, 12.0 ms for 32 kbps voice coding rate, 24.0 ms for 16 kbps voice coding rate and so on. As before, voice packets not transmitted within 5.5 ms will be discarded.

Figures 18 to 21 show the mean data packet waiting delay versus the QA data throughput for different numbers of active voice calls and different voice coding rates. These curves show clearly that a reduction in the voice coding rate leaves more capacity for data packet transmission.

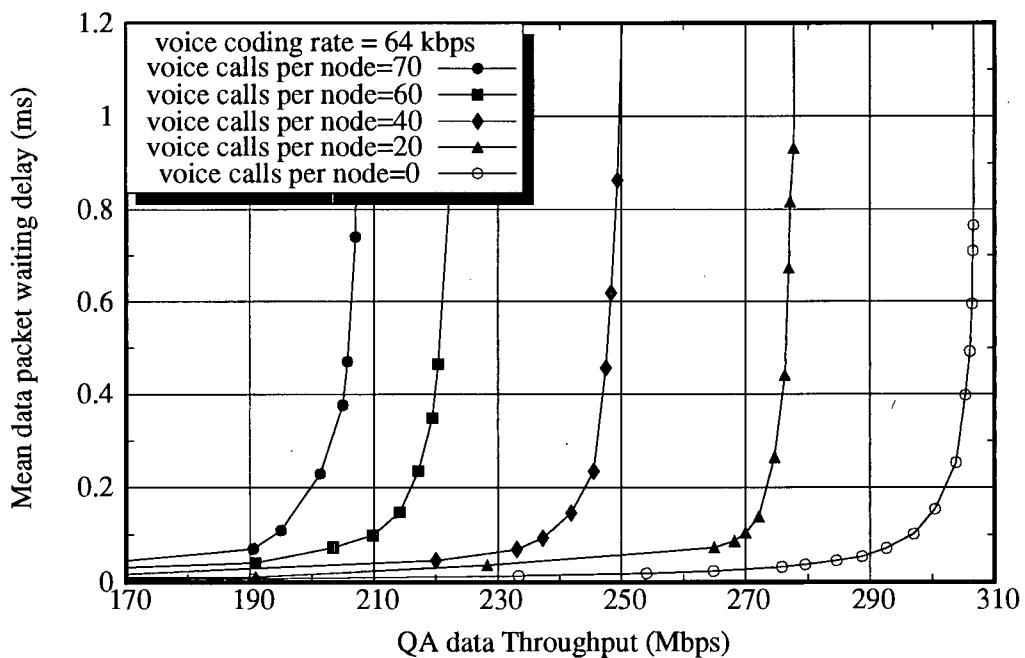


Figure 18. Mean data packet waiting delay vs. QA data Throughput for various number of active voice calls with voice coding rate of 64 kbps

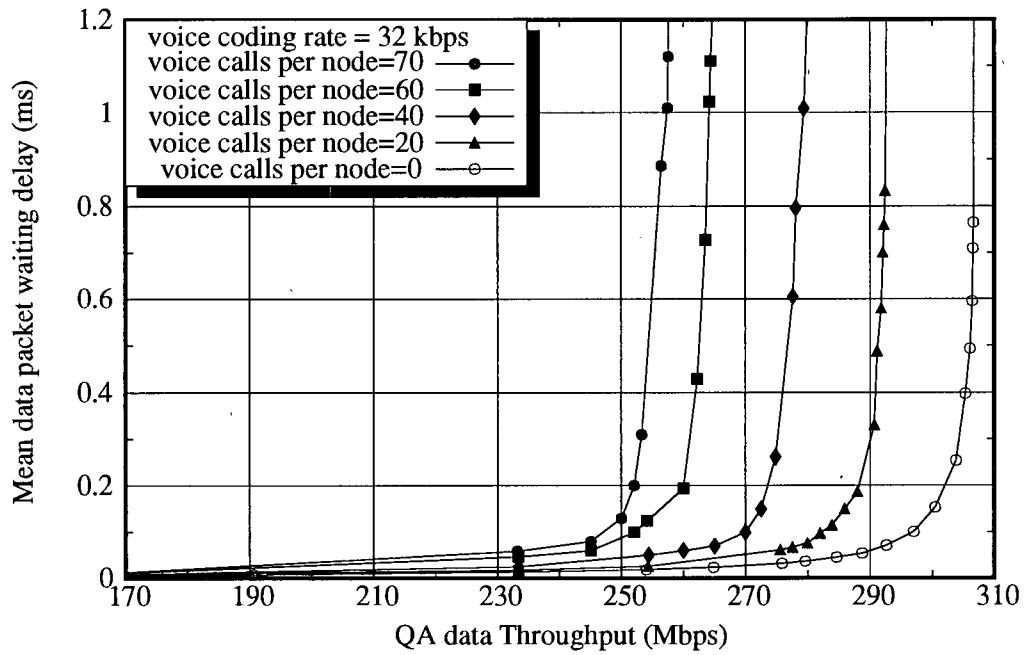


Figure 19. Mean data packet waiting delay vs. QA data Throughput for various number of active voice calls with voice coding rate of 32 kbps

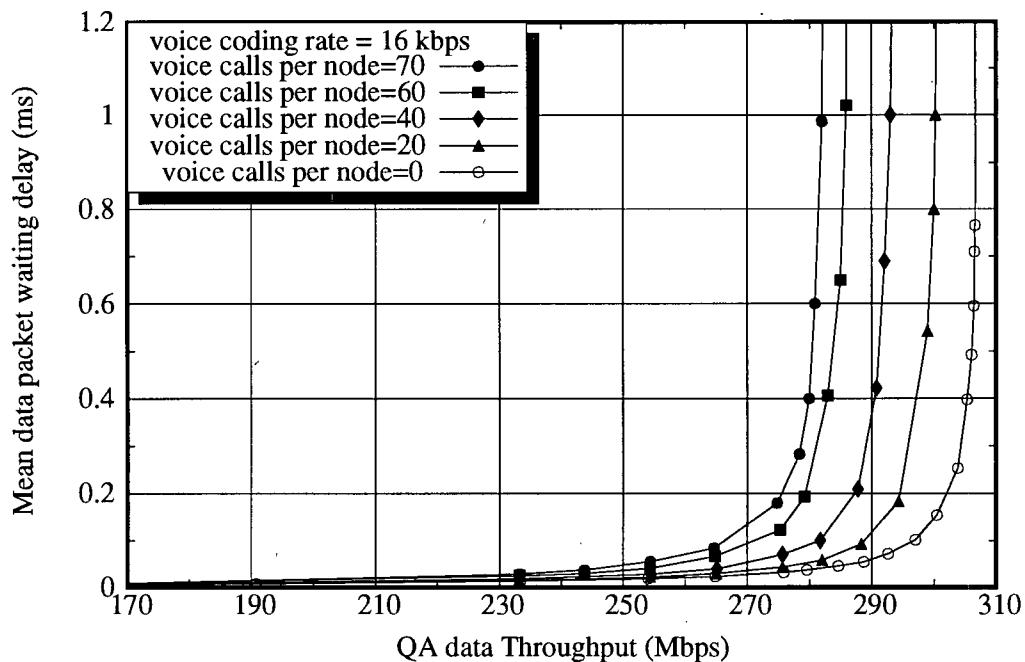


Figure 20. Mean data packet waiting delay vs. QA data Throughput for various number of active voice calls with voice coding rate of 16 kbps

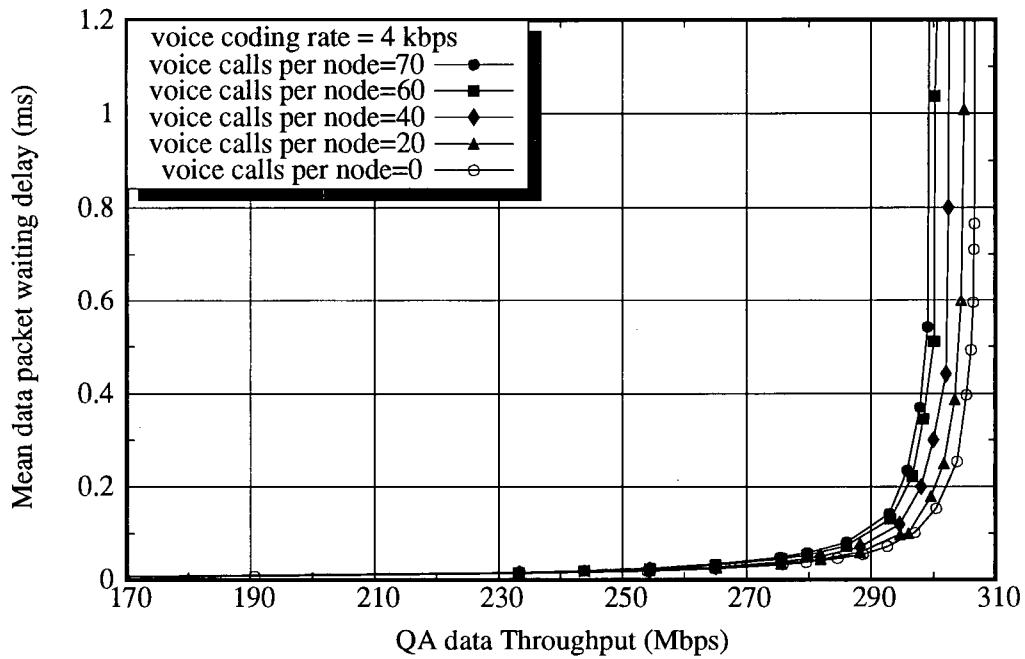
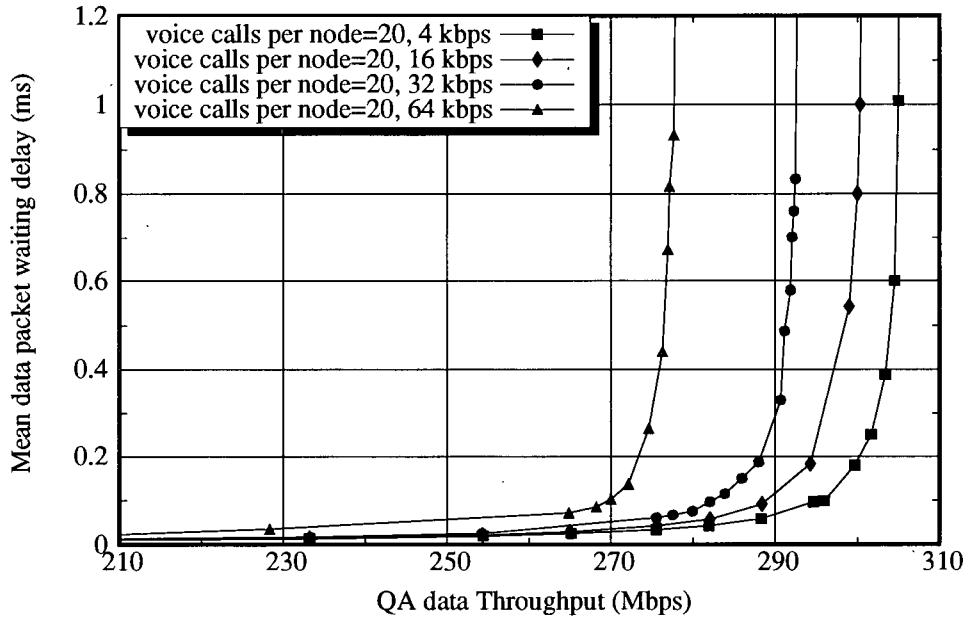


Figure 21. Mean data packet waiting delay vs. QA data Throughput for various number of active voice calls with voice coding rate of 4 kbps

Figure 22 shows the mean data packet waiting delay versus the QA data throughput for different voice coding rates with 20 and 70 active voice calls activated per node. Reduction of the voice coding rate from 64 to 4 kbps increases the maximum QA data throughput from 280 to 305 Mbps. With 70 active voice calls per node, the QA data throughput increases from 200 to 300 Mbps as the voice coding rate is reduced from 64 to 4 kbps.

(a) 20 voice calls per node



(b) 70 voice calls per node

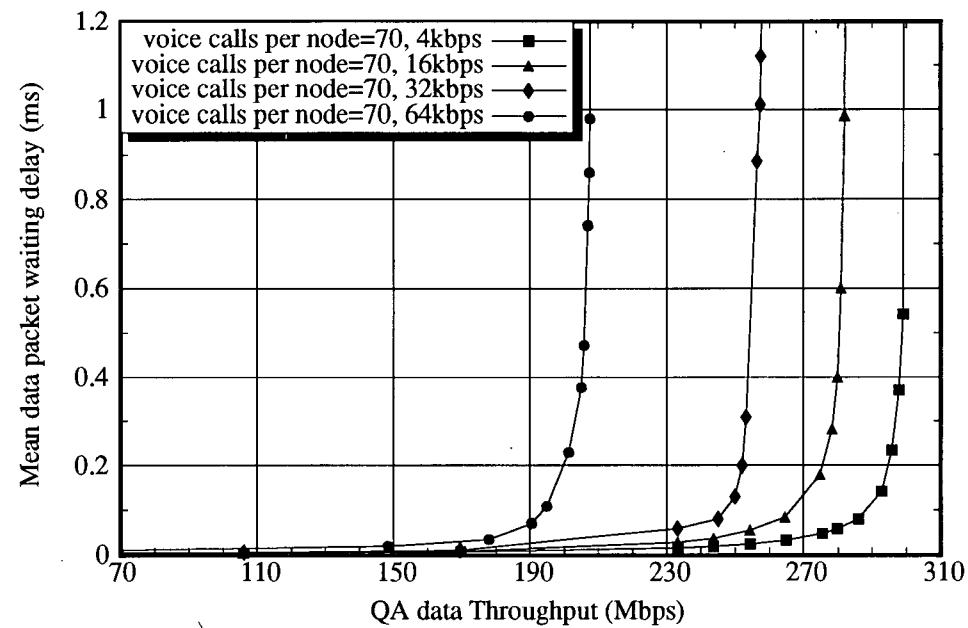


Figure 22. Mean data packet waiting delay vs. QA data Throughput with different numbers of active voice calls per node for various voice coding rates

There is a simple relationship between the voice coding rate and the number of active voice calls, as follows, where  $R_i$  denotes the voice coding rate,  $N_i$  denotes the number of active voice calls and  $C$  is a constant:

$$R_1 N_1 = R_2 N_2 = C$$

From this equation, one sees that the mean data packet waiting delay curve for 64 kbps voice and 20 voice calls per node is the same as the mean data packet waiting delay curve for 32 kbps voice and 40 voice calls per node. Both configurations give the same QA voice throughput in the network and the corresponding curves of mean data packet delay versus QA data throughput overlay each other, as shown in Figure 23. Similar comments apply for other sets of curves. Simulation results of the QA voice throughput achieved for different voice coding rates and number of voice calls activated are summarized in Table 6.

number of voice calls	QA voice Throughput achieved ( Mbps )			
	64 kbps	32 kbps	16 kbps	4 kbps
1000	25.970	12.771	6.345	1.657
2000	51.659	25.769	12.830	3.229
3000	78.246	38.524	19.625	4.913
3500	89.617	44.536	22.376	5.784

Table 6. QA voice Throughput achieved for various voice coding rates and number of active voice calls

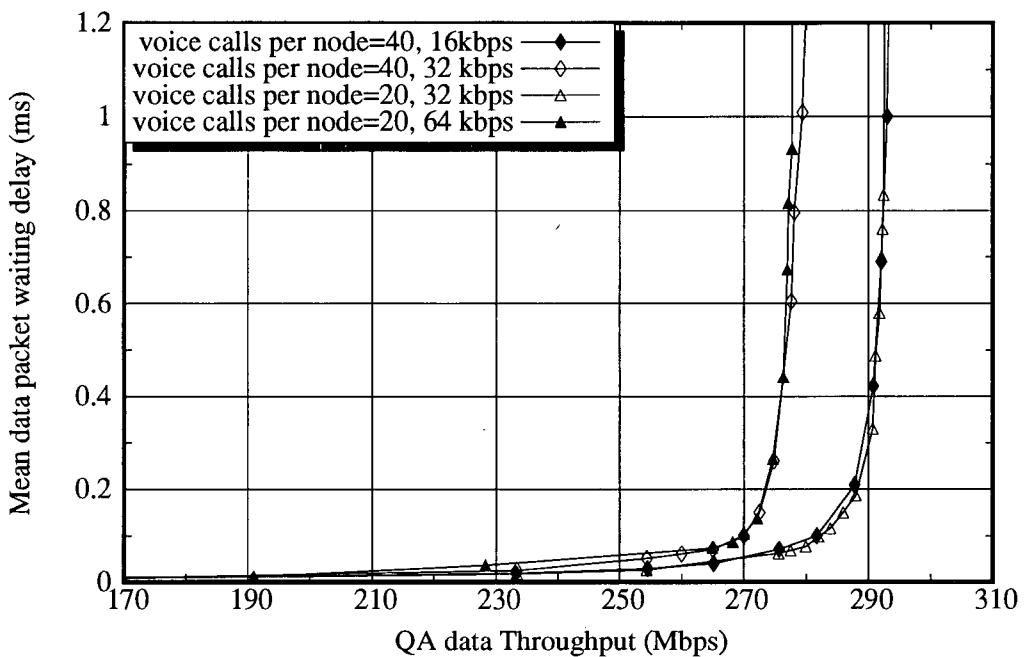


Figure 23. Relation between voice coding rate and number of active voice calls activated

Because the voice coding rate at the network stations primarily affects the QA voice throughput achieved in the network, the maximum number of voice calls a network can accommodate within an acceptable 2% of discarded voice packets range may vary with voice coding rate. The network can hold approximate 21000 voice calls for 32 kbps voice coding rate, 42000 voice calls for 16 kbps voice coding rate and 168000 voice calls for 4 kbps voice coding rate. These numbers are based on the assumptions of our simulation models.

Since the voice coding rate limits the QA voice throughput achieved in the network for a fixed number of active voice calls, the mean voice packet waiting delay versus the QA voice throughput remains the same as before and is as shown in Figure 13. However, for a fixed number of voice calls activated in the network, the time from when a voice station encodes a voice cell to the time of decoding at the destination may vary substantially with voice coding

rate. Table 7 shows two typical examples of 1000 active voice calls ( 20 voice calls per node for 50 nodes ) and 3500 active voice calls in the network.

number of voice calls in network	voice coding rate (kbps)	voice cell encoding time (ms)	mean voice packet waiting delay (ms)	propaga-tion delay (ms)	voice cell decoding time (ms)	total elapse time (ms)
1000	64	6.0	0.012	t	6.0	12.012+t
	32	12.0	0.010	t	12.0	24.010+t
	16	24.0	0.010	t	24.0	48.010+t
	4	96.0	0.009	t	96.0	192.009+t
3500	64	6.0	0.016	t	6.0	12.016+t
	32	12.0	0.013	t	12.0	24.013+t
	16	24.0	0.011	t	24.0	48.011+t
	4	96.0	0.009	t	96.0	192.009+t

Table 7. Delay comparison for different voice coding rate voice stations

As shown in Table 7, the voice coding rate of a voice station plays a very important role in the total elapse time for transmitting a voice cell since the propagation delay, t, is almost constant for all cases in uniform destination distribution and the variation of mean voice waiting delay is very small compared to the encoding/decoding time.

## **Chapter 4. Delay Analysis of QA Access for Inter- and Intra-MAN Voice and Data Integrated Traffic within a Backbone MAN Network**

In this chapter, voice and data packet transmission delays will be analyzed for both inter-and intra-MAN traffic. Inter-MAN traffic is that flowing on two or more DQDB MANs within a backbone MAN network; intra-MAN traffic is confined to one access DQDB MAN only. The number of homogeneous bridges connecting each access MAN depends on its physical location within the backbone MAN network, as illustrated in Figure 24.

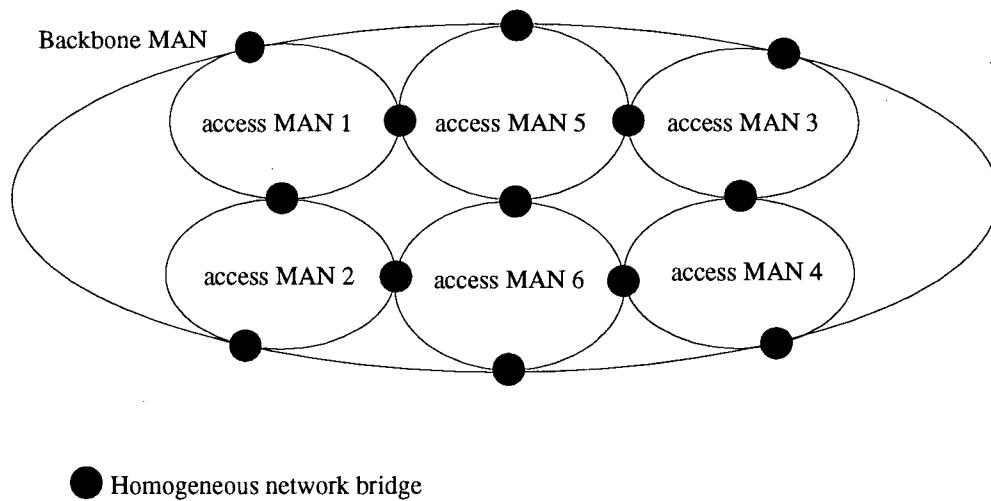


Figure 24. Geographical location of access MANs within a backbone MAN

In Figure 24, six access MAN are connected to adjacent MANs and to a backbone MAN. Four homogeneous bridges connect both MAN 5 and MAN 6 which are located in the middle of the MAN cluster. Each of the other four access MANs has three homogeneous connecting bridges. According to this connection scheme, any inter-MAN traffic need only go through two homogeneous bridges, at most, in order to reach the destinated MAN within a backbone MAN network. Details of the homogeneous bridges connection scheme are provided next.

## **Section 4.1**

### **Homogeneous Bridges Connection Scheme**

From the previous chapter we know that the upstream nodes have a transmission advantage over the downstream nodes; therefore, bridge connections are arranged according to this property in order to reduce the overall delay, including bridge delay.

For MANs 1 to 4, two of the three bridges connect to the two neighborhood MANs while the third bridge is used to direct packets to the remaining three MANs via the backbone MAN. The bridge connected to the backbone MAN will have a heavier traffic load than the other two bridges provided that the inter-MAN traffic is uniformly distributed among the six access MAN. Therefore, we connected the bridge with the heaviest traffic load as the head node of the access MAN, as shown in Figure 25.

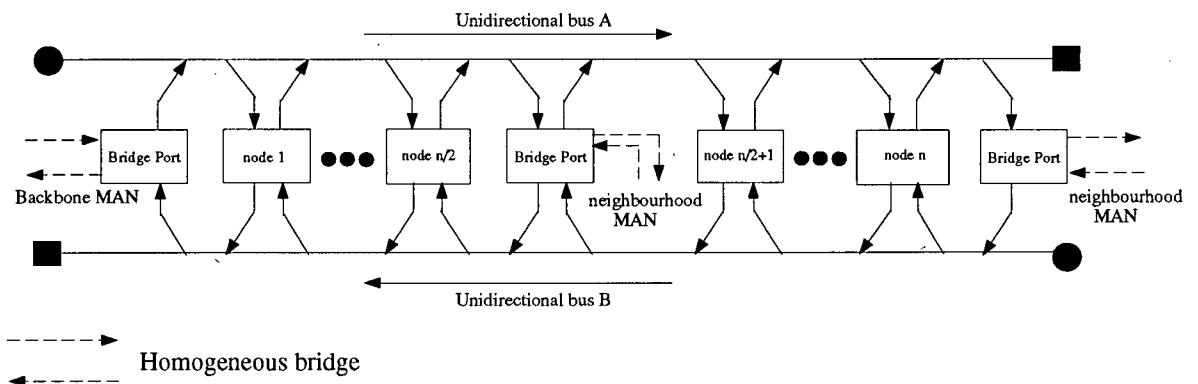


Figure 25. Illustration of three homogeneous bridges connection

For MAN 5 and MAN 6, three of the four bridges are connect to the three neighborhood MANs and the fourth bridge is used to deliver packets to the other two MANs via the backbone MAN. With four bridges, we can actually put one more bridge port into the access MAN or simply connect the fourth bridge to either one of the other three bridge ports. In our connection scheme, we implemented the latter scheme, for simplicity. Since upstream nodes provide a transmission

advantage, the fourth bridge is connected to the head bridge port for the neighboring MAN. As a result, two head bridge ports will be responsible for delivering inter-MAN traffic destined to four access MANs ( each bridge port for two access MANs ), as illustrated in Figure 26.

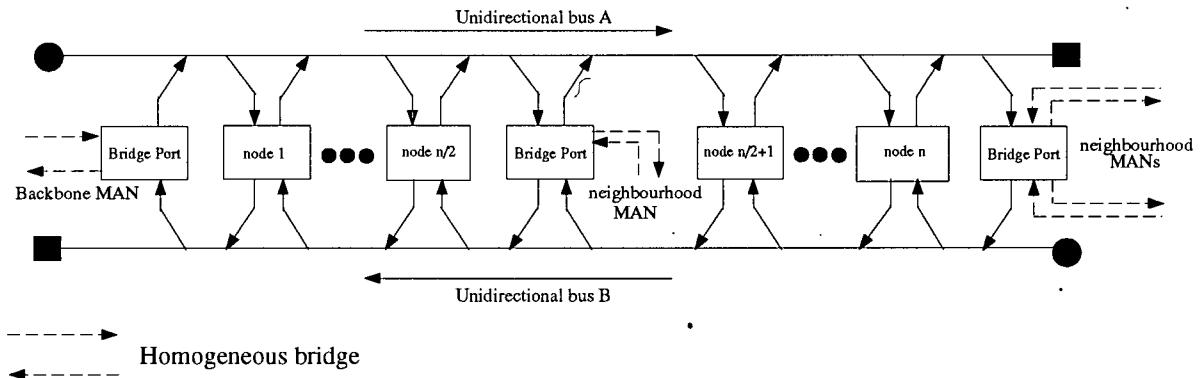


Figure 26. Illustration of four homogeneous bridges connection

Each bridge port functions as a node, with the same nodal media access mechanism. Twenty buffers were used for each priority level in each bridge port in our model. In our simulations, the processing delay and the propagation delay of the bridges are neglected. Each of these delays would typically be less than 0.25 ms. Figure 27 shows the physical arrangement and bridge connections of six access MANs within a backbone MAN network in our model.

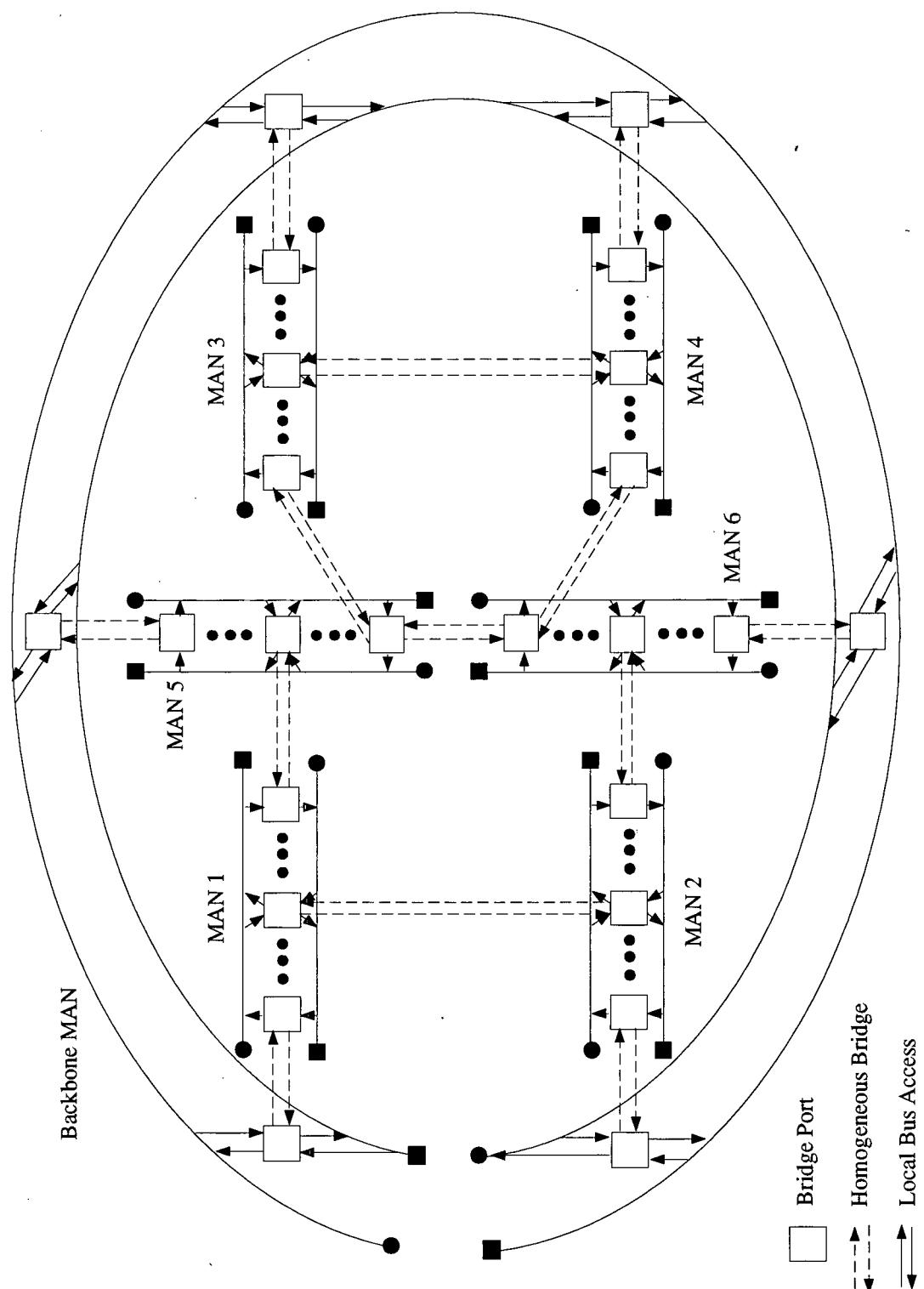


Figure 27. Physical representation of a backbone MAN network

## **Section 4.2 Simulation Model**

Each node of each access MAN within the backbone MAN network is basically the same as the previous models except that now a predetermined level of voice calls and data packets in each node will be inter-MAN traffic. Inter-MAN traffic occurs at random. The destination node and the destination MAN for all inter-MAN traffic of each access MAN is uniformly distributed among the two hundred and fifty nodes in the other five access MANs. Since the DQDB QA protocol allows only three priority access levels and since the highest priority level is reserved for control signalling, inter-MAN and intra-MAN traffic priorities are equal. All assumptions from the previous simulation models apply for each access MAN within the backbone MAN network. Inter-MAN voice call destination are paired. For example, a voice call initiated in MAN 1, node 26 with destination at MAN 5, node 4 must be paired with a voice call from MAN 5, node 4 to MAN 1, node 26. All voice calls are processed to have silence deletion with a steady-state talkspurt probability of 0.4.

Each of the six backbone MAN bridge port functions as an access MAN node. Sixty time slots separate successive bridge ports ( nodes ) in the backbone MAN. Twenty buffers serve each priority level at each backbone MAN node. An extremely high level of inter-MAN traffic will cause some lost packets, because of limited buffer queue capacity.

The source code of the backbone MAN network simulation model is provided in Appendix A. The model is written in **SIMSCRIPT II.5** [7].

## **Section 4.3 Preliminary Simulation Results**

In a preliminary analysis, the packet mean delay of inter-MAN and intra-MAN packets was

investigated for a fixed percentage of inter-MAN voice and data traffic. Voice coding rates of 32 and 4 kbps were used, with 50 voice calls activated in each node. A total of  $300 \times 50 = 15000$  voice calls were active at any time in the six-MAN network.

Figures 28 and 29 show the mean inter- and intra-MAN data packet transmission delay<sup>3</sup> versus the total QA data traffic<sup>4</sup> for 32 kbps voice coding. Figures 30 and 31 show the mean inter- and intra-MAN data packet transmission delay versus the total QA data traffic for 4 kbps voice coding.

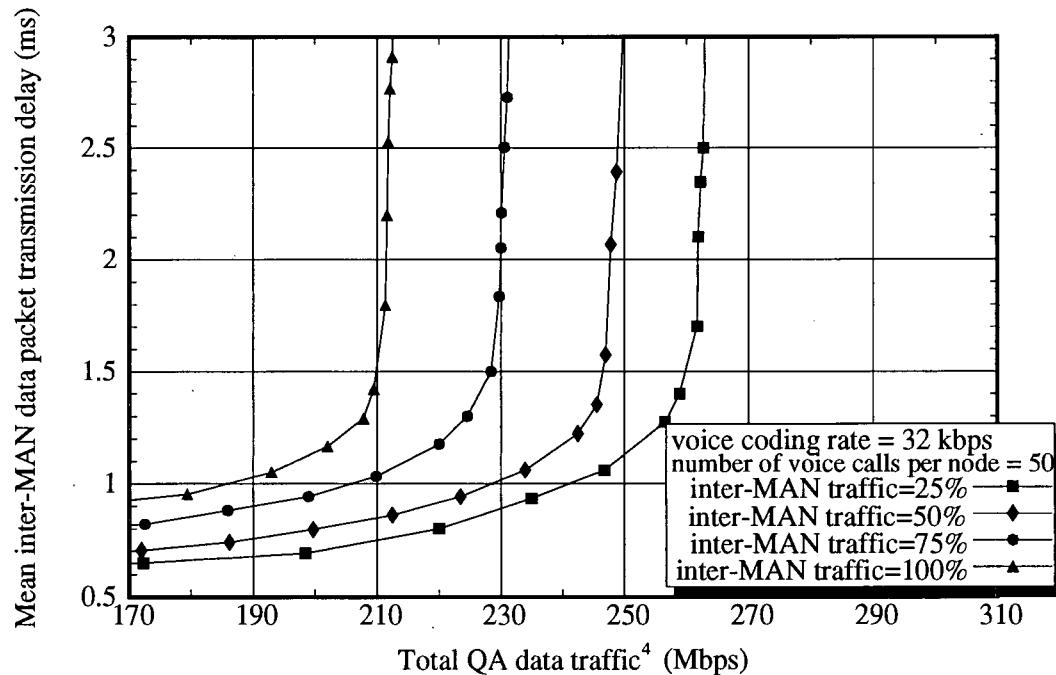


Figure 28. Mean inter-MAN data packet transmission delay vs. total QA data traffic for 32 kbps voice coding rate

<sup>3</sup> Mean transmission delay is calculated from the time a packet enters the buffer of a node until it reaches the destined node.

<sup>4</sup> Total QA data traffic is the sum of inter- and intra-MAN data traffic per single access MAN.

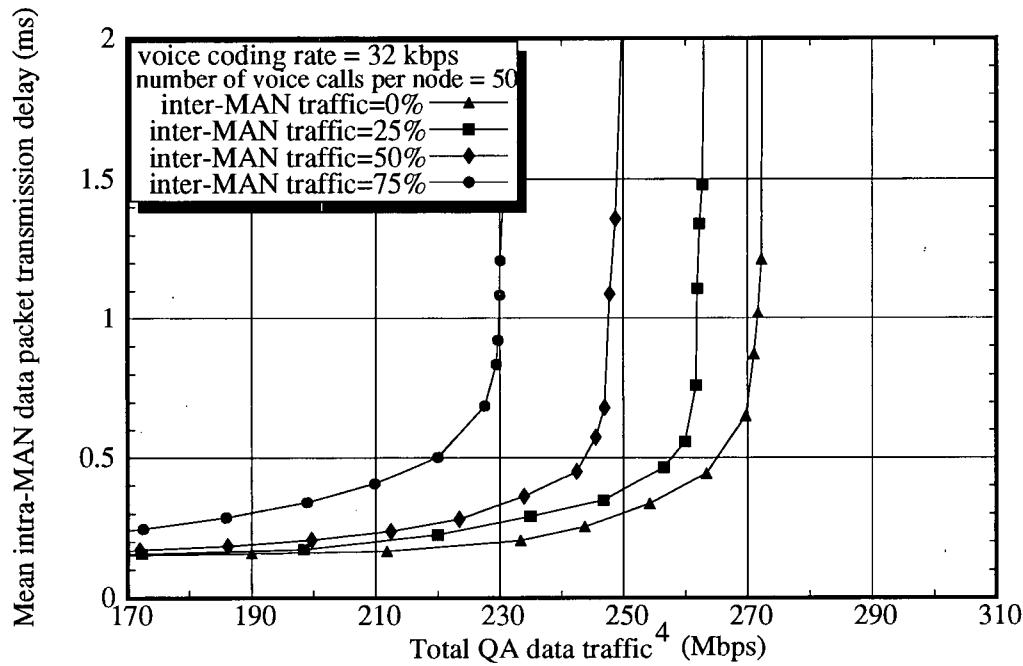


Figure 29. Mean intra-MAN data packet transmission delay vs. total QA data traffic for 32 kbps voice coding rate

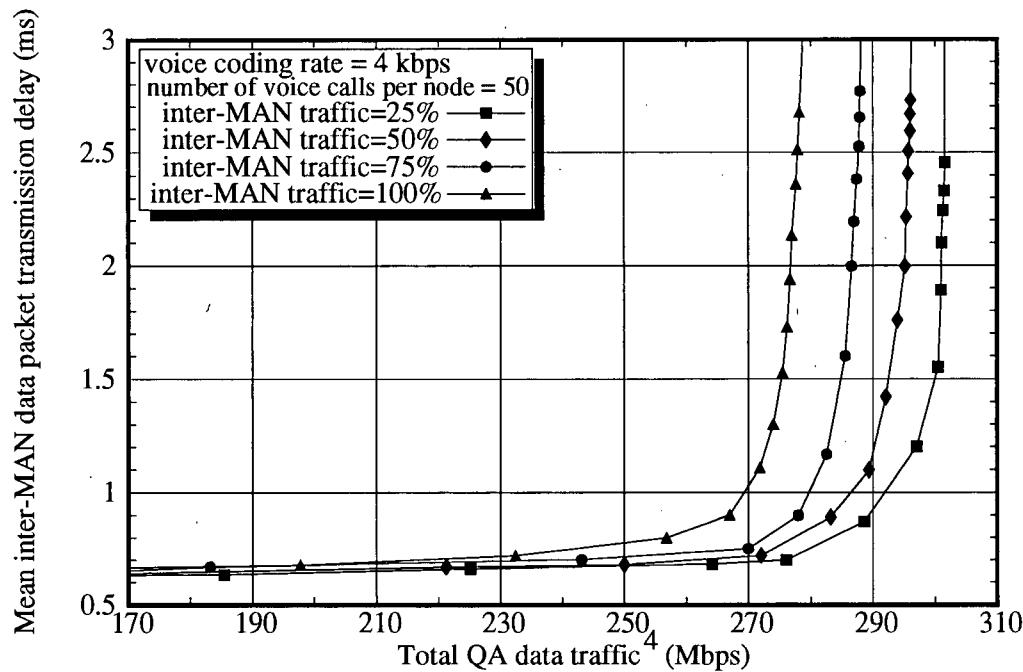


Figure 30. Mean inter-MAN data packet transmission delay vs. total QA data traffic for 4 kbps voice coding rate

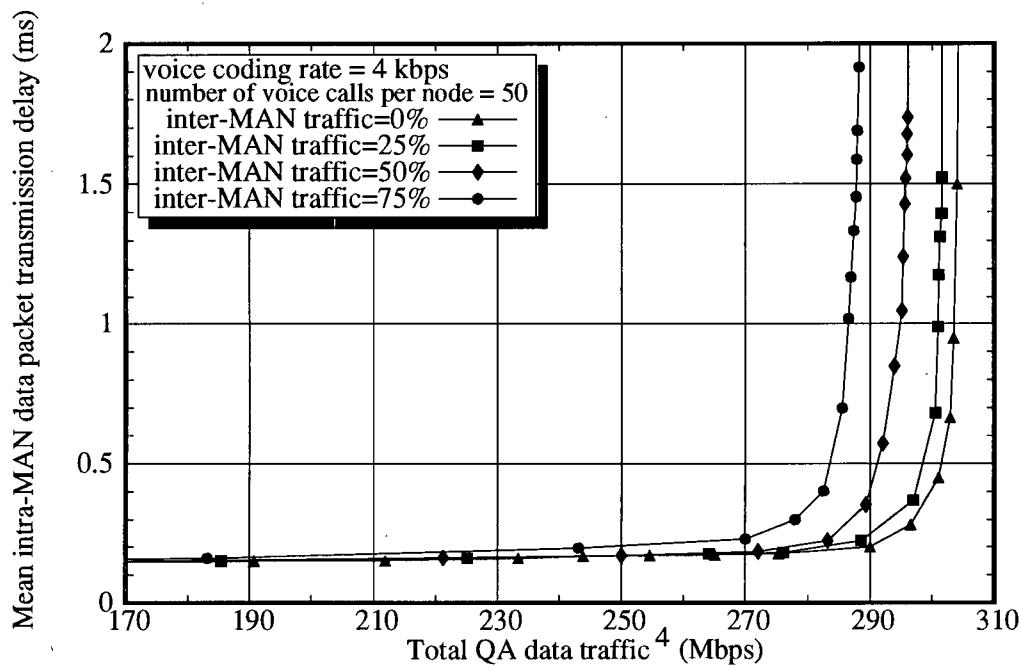


Figure 31. Mean intra-MAN data packet transmission delay vs. total QA data traffic for 4 kbps voice coding rate

The mean inter- and intra-MAN voice packet transmission delay<sup>5</sup> with different percentages of inter-MAN traffic and the two voice coding rates for 50 voice calls per node are summarized in Table 8. Since there is a limited buffer space in each bridge port, the voice throughput achieved may be less than the percent inter-MAN traffic might suggest. Moreover, the voice throughput achieved is calculated excluding any overheads.

As shown in Table 8, the inter- intra-MAN voice packet transmission delays depend not only on their respective traffics, but also on the total traffic achieved in the network ( i.e. the summation of the inter- and intra-MAN voice traffics ). For example, the fifth entry of the mean intra-MAN voice packet transmission delay in Table 8 is equal to 0.153 ms with a 32.209 Mbps intra-MAN voice traffic; on the other hand, the second entry of the mean intra-MAN voice

<sup>5</sup> The value of the voice packet transmission delay is obtained close to the network's capacity.

packet transmission delay is equal to 0.159 ms with a 7.983 Mbps intra-MAN voice traffic. The latter entry has a longer mean transmission delay than the fifth entry with its smaller intra-MAN voice traffic. This result can be explained by the fact that the total network voice traffic for the second entry is more than 1.5 times that for the fifth entry. Therefore, the mean intra-MAN voice packet transmission delay in the second entry is larger than for the fifth entry.

voice coding rate (kbps)	% inter-MAN traffic	inter-MAN voice traffic (Mbps)	intra-MAN voice traffic (Mbps)	total voice traffic (Mbps)	total source-destination voice throughput (Mbps)	mean inter-MAN voice packet tx. delay (ms)	mean intra-MAN voice packet tx. delay (ms)
32	100	62.025	0	62.025	29.841	1.224	0.0
	75	47.176	7.983	55.159	31.010	1.024	0.159
	50	31.047	16.167	47.214	31.047	0.735	0.151
	25	16.043	24.134	40.177	32.132	0.680	0.151
	0	0	32.209	32.209	32.209	0.0	0.153
4	100	7.843	0	7.843	3.772	0.990	0.0
	75	5.875	1.070	6.945	3.835	0.899	0.160
	50	4.163	1.986	6.149	4.028	0.733	0.152
	25	2.248	2.802	5.050	4.066	0.687	0.151
	0	0	4.071	4.071	4.071	0.0	0.152

Table 8. Mean inter- and intra-MAN voice packet transmission delay for 50 voice calls per node

## Section 4.4

### Statistical Results for The Six-MAN Network

The following statistics are obtained from one typical simulation result with 50 active voice calls per node for 32 kbps voice coding; The mean data packet Poisson arrival rate per node is

8500 packets/sec, with 25% inter-MAN traffic. Table 9 summarizes relevant delay and throughput results.

inter-MAN voice throughput (Mbps)	8.006
inter-MAN voice packet tx. delay (ms)	0.677
intra-MAN voice throughput (Mbps)	24.017
intra-MAN voice packet tx.delay (ms)	0.147
total voice throughput (Mbps)	39.711
inter-MAN data throughput (Mbps)	45.043
inter-MAN data packet tx. delay (ms)	0.681
intra-MAN data throughput (Mbps)	135.129
intra-MAN data packet tx. delay (ms)	0.171
total data throughput (Mbps)	225.091
total network throughput (Mbps)	264.802

Table 9. A typical simulation results for the six-MAN network

Figure 32 shows the mean inter-MAN voice and data packet waiting delay over the homogeneous bridges for each access MAN. Bridge 1 connects to the backbone MAN and is located at the head of one bus, bridge 2 is located in the middle of each access MAN, while bridge 3 is another bridge located at the head of another MAN bus.

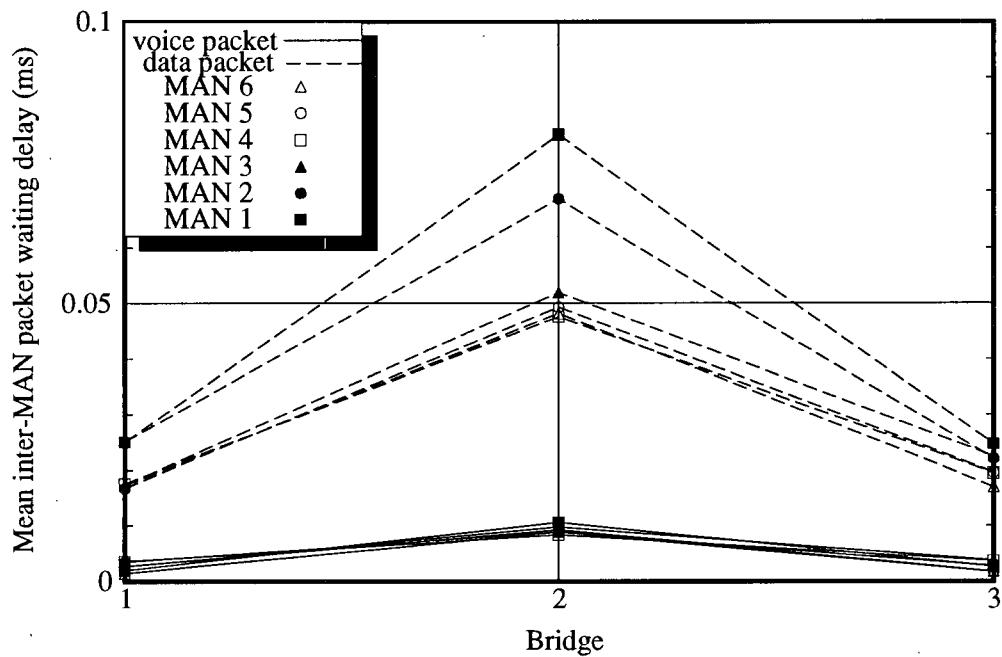


Figure 32. Mean inter-MAN packets waiting delay over homogeneous bridges

As shown in Figure 32, the mean packet waiting delay in bridge 2 (the bridge located in the middle of each access MAN) is higher than at the bridges located at the head of two buses (bridges 1 and 3). Therefore, it is useful to route much of the inter-MAN traffic through the homogeneous bridges located at the head of bus.

Figure 33 shows the mean inter-MAN voice and data packet transmission delay via different bridges. The locations of bridges 1, 2 and 3 are as before. Figure 32 shows the packet waiting delay of different bridges, while Figure 33 shows the packet transmission delay that via the same bridges as in Figure 32.

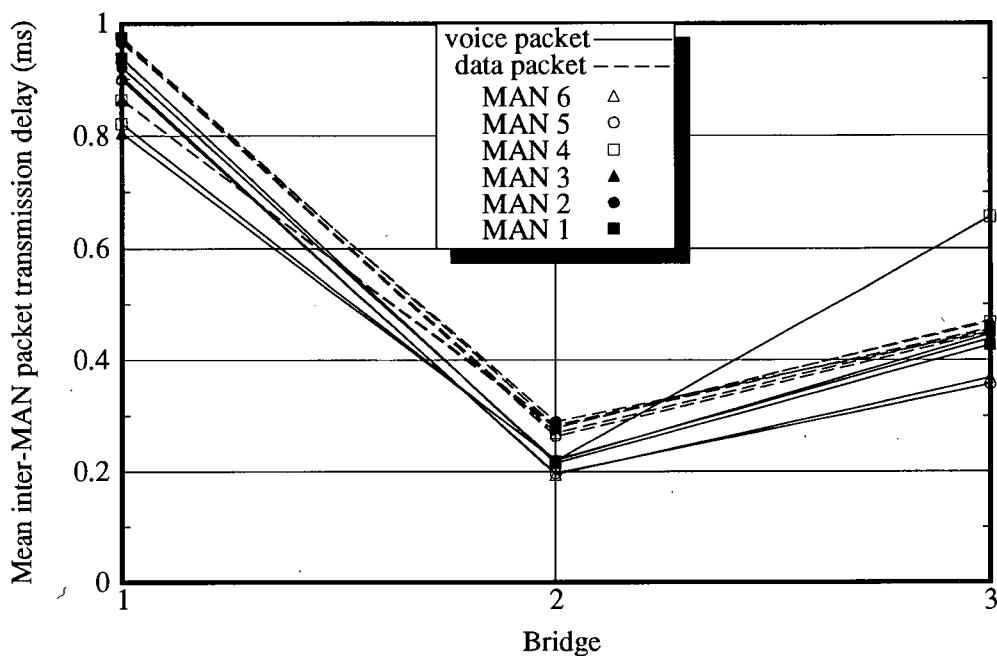


Figure 33. Mean inter-MAN packet transmission delay via different bridges

From figure 33, the mean inter-MAN voice and data packet transmission delay via bridge 2 is less than the packet transmission delay via the other two bridges because of the physical location of bridge 2, and because the packet waiting delay is negligible compared to the propagation delay. From the data obtained for Figure 33, the mean time a packet spends in the backbone MAN is approximately 0.5 ms.

It is shown in the previous chapter that the upstream nodes have a transmission advantage over the nodes located in the middle of two buses in that the waiting delay of a packet at upstream nodes is smaller than for the middle nodes. However, Figure 33 together with Figure 32 shows that the propagation delay is relatively large compared to the waiting delay of a packet; therefore, the transmission delay of the uniformly distributed destination intra-MAN voice and data packets with respect to the nodes' locations will have a shape as shown in Figure 34. All

six access MANs have the same shape of intra-MAN packets transmission delay, we only show the data for one access MAN, access MAN 1.

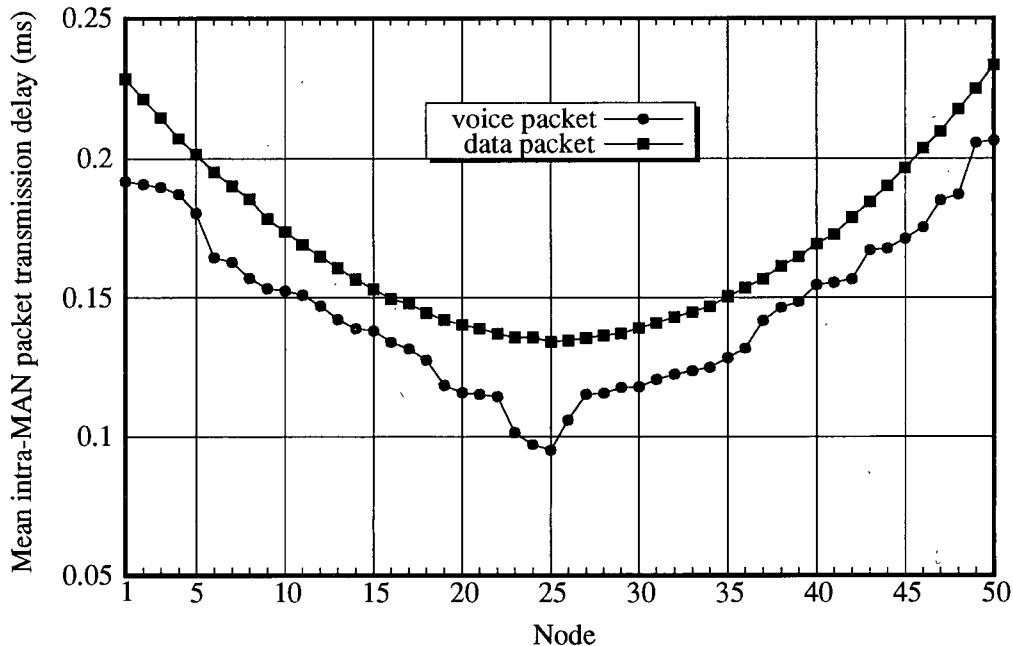


Figure 34. Mean intra-MAN packets transmission delay vs. nodes locations

As expected, the transmission delays for both intra-MAN voice and data packets are larger at the upstream nodes than for the nodes located in the middle of the buses. The reason why the voice packet transmission delay curve is not as smooth as the data packet transmission delay curve in Figure 34 is because the number of voice calls activated in each node is not large enough to uniformly spread the voice call destinations of a node among the remaining 49 MAN nodes.

The inter-MAN voice and data packet transmission delay will depend not only on the nodes' locations, but also on the connection scheme for the homogeneous bridges that connect to each access MAN. Figure 35 shows the inter-MAN data packet transmission delay with respect to the nodes locations. Access MANs 1 to 4 have exactly the same homogeneous bridges connection scheme, and access MANs 5 and 6 have connection schemes which are identical to each other.

Therefore, we plot only one curve for access MANs 1 to 4 and one curve for access MANs 5 and 6.

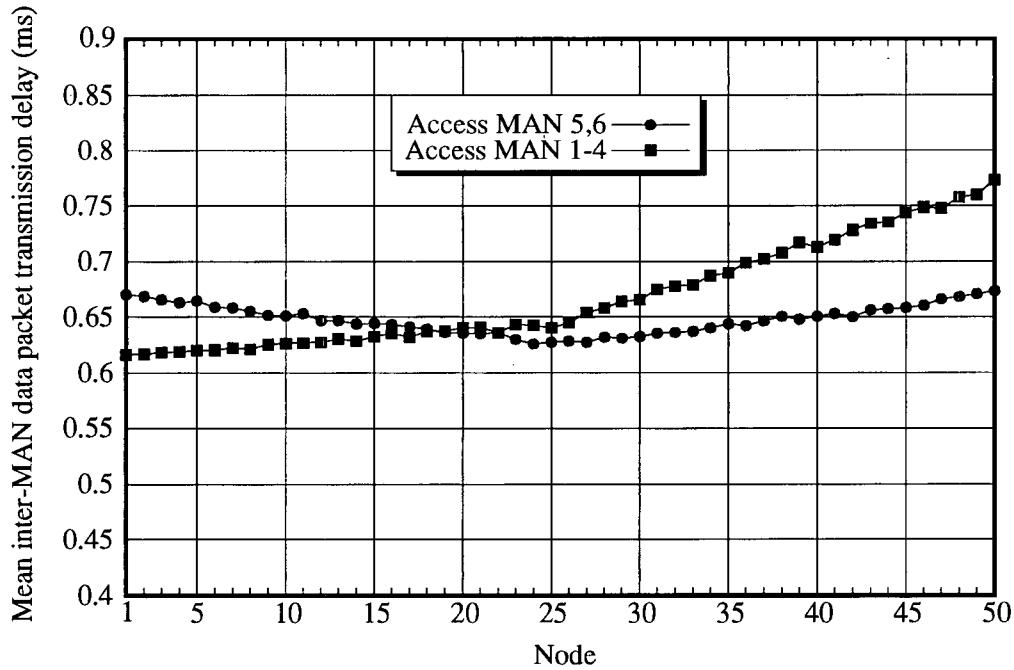


Figure 35. Mean inter-MAN data packet transmission delay vs. nodes locations

In Figure 35, node 1 is closest to the homogeneous bridge connected to the backbone MAN, while node 50 is the farthest node from the backbone MAN. For access MANs 1 to 4, three out of the five other MANs require inter-MAN traffic to be route through the backbone MAN. Since the destination MAN traffic is uniformly distributed, the nodes further away from the backbone MAN will have larger transmission delays than nodes closer to the backbone MAN, as shown in Figure 35. For access MANs 5 and 6, the homogeneous bridges connection scheme divides the uniformly distributed destination MAN traffic relatively even among the four bridges, and the fluctuation in nodal delay is not as large as access MAN 1 to 4. Moreover, the nodal inter-MAN packet transmission delay for access MANs 5 and 6 also preserved the shape as the intra-MAN packet transmission delay.

Figure 36 shows the inter-MAN voice packet transmission delay with respect to the node locations; this figure reflects the same general characteristics as in Figure 35.

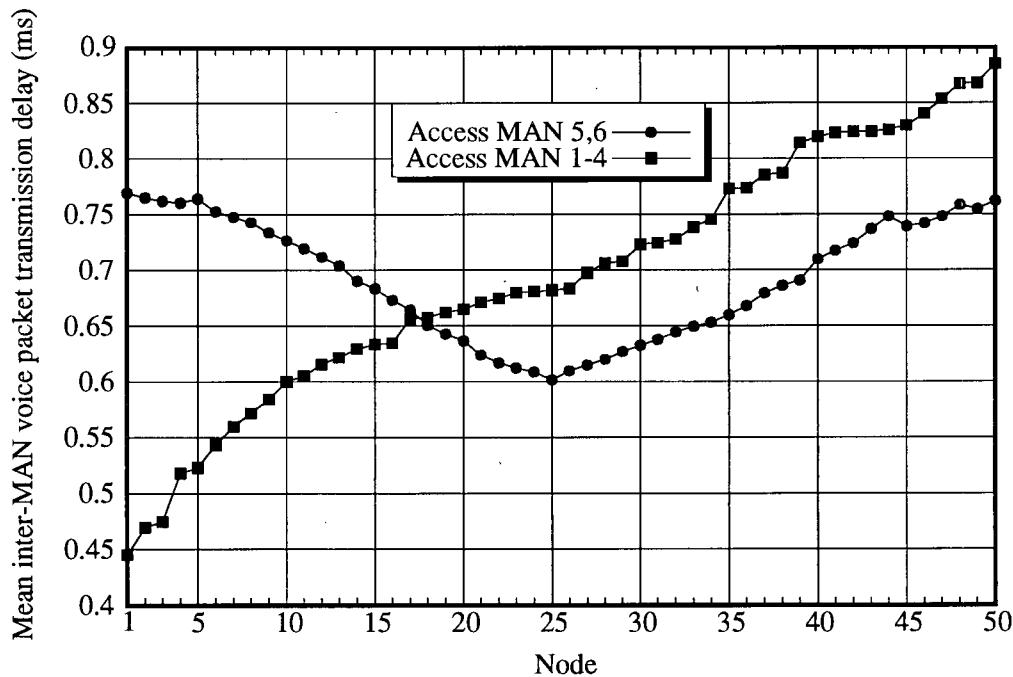


Figure 36. Mean inter-MAN voice packet transmission delay vs. nodes locations

For this typical simulation in this chapter, the percentage of discarded voice packets in the bridge ports is equal to 0.0014%, whereas the percentage of discarded data packets is 0.165%. Recall that voice packets are not retransmitted, whereas data packets are retransmitted.

## **Section 4.5** **General Results of The Inter- and Intra-MAN Traffic** **within a Backbone MAN Network**

In this section, the inter- and intra-MAN voice and data packet transmission delay for different throughput and inter-MAN traffic levels are presented. In order to achieve the desired voice throughput, the assumption of 50 voice calls per node in the previous section is waived.

Instead, additional active voice calls per node are simulated for higher voice throughput. Figure 37 to 39 show the mean inter- and intra-MAN data packet transmission delay versus the total QA data traffic for different voice throughputs and different inter-MAN traffic levels.

Figures 40 and 41 show the mean inter- and intra-MAN voice packet transmission delay versus the total QA voice traffic<sup>6</sup> for different inter-MAN traffic levels.

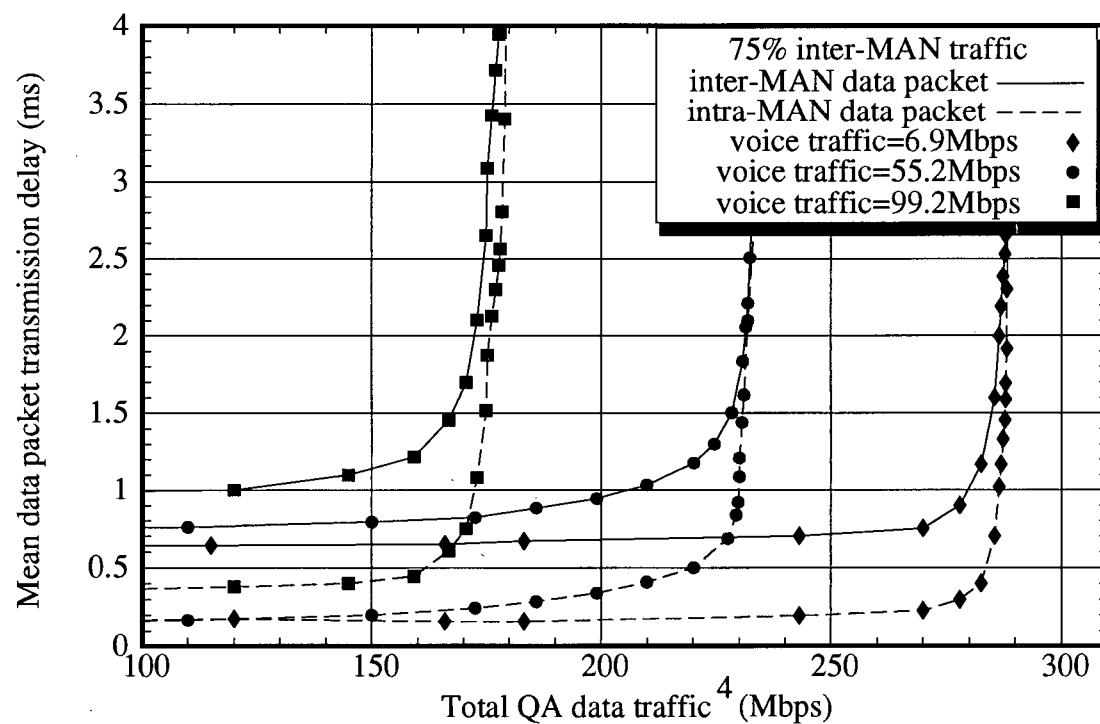


Figure 37. Mean inter- and intra-MAN data packet transmission delay vs. total data traffic for 75% inter-MAN traffic

<sup>6</sup> Total QA voice traffic is the sum of inter- and intra-MAN voice traffic per single access MAN.

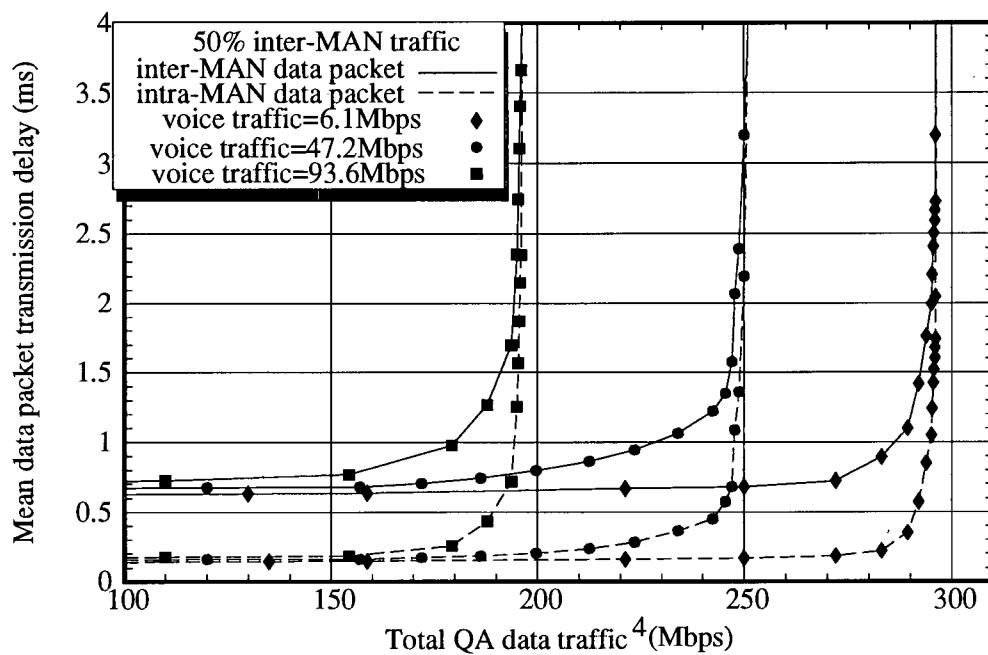


Figure 38. Mean inter- and intra-MAN data packet transmission delay vs. total data traffic for 50% inter-MAN traffic

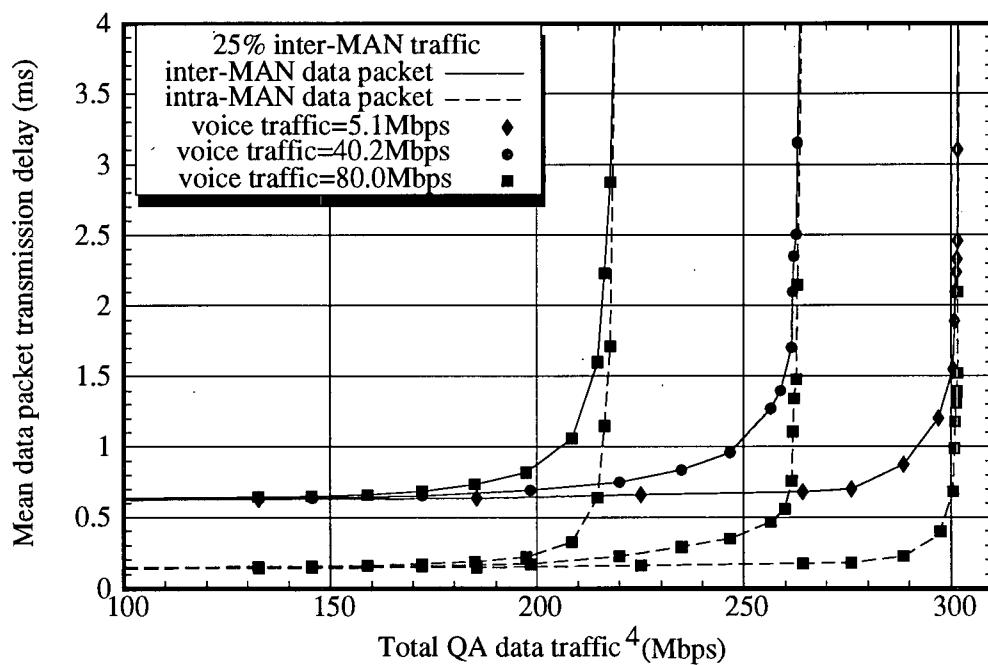


Figure 39. Mean inter- and intra-MAN data packet transmission delay vs. total data traffic for 25% inter-MAN traffic

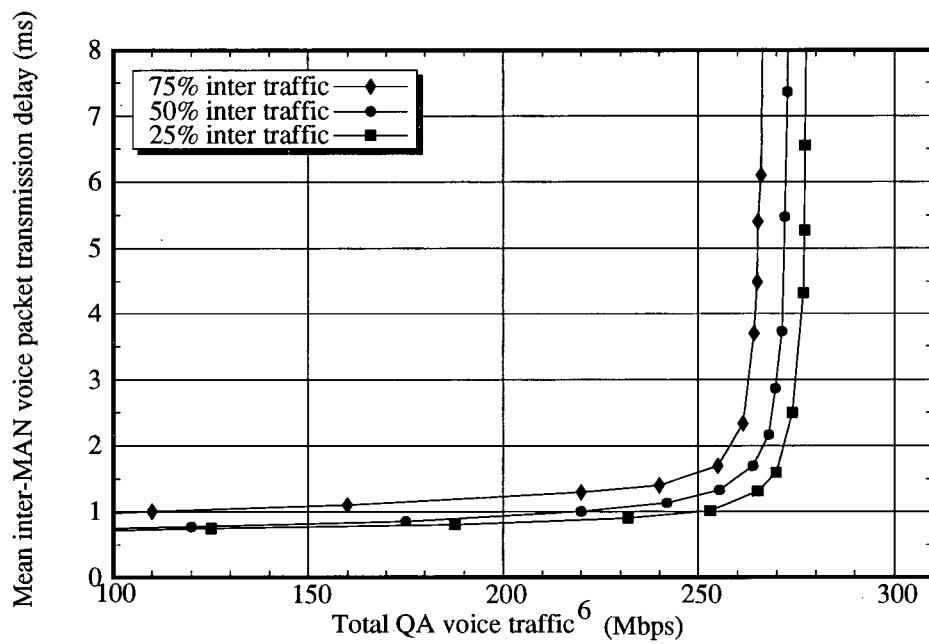


Figure 40. Mean inter-MAN voice packet transmission delay vs. total QA voice traffic for various levels of inter-MAN traffic

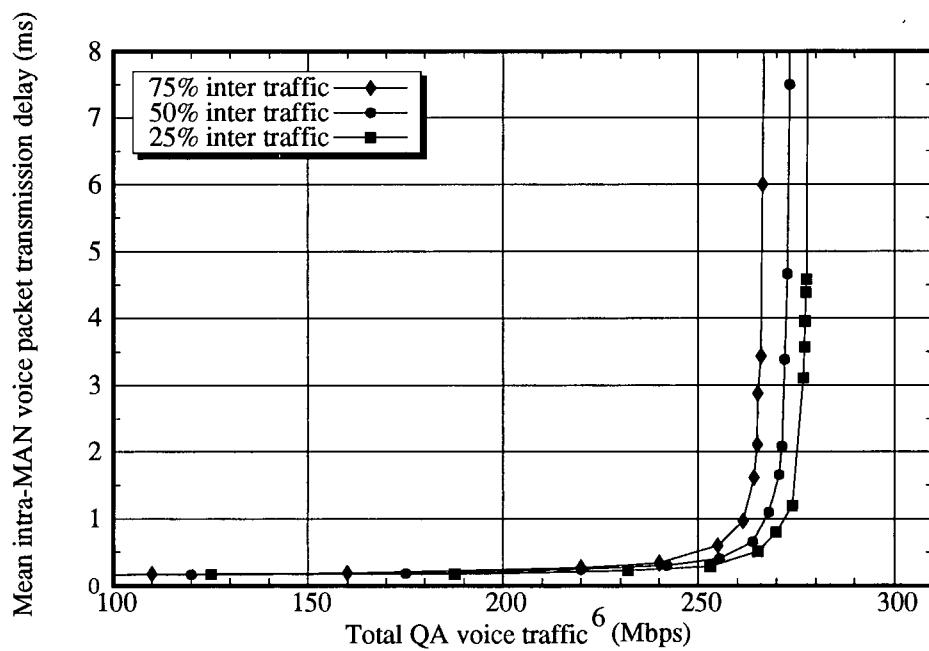


Figure 41. Mean intra-MAN voice packet transmission delay vs. total QA voice traffic for various levels of inter-MAN traffic

Table 10 summarizes the maximum data throughput achieved and its respective data packet transmission delay for various voice throughput levels. Note that the total throughput ( voice and data ) falls as the voice throughput increases since voice packet overheads throughput also increases as the voice throughput increases. Packet overheads are excluded for the entries of voice throughput in Table 10.

Voice throughput (Mbps)	Maximum data throughput achieved (Mbps)	Total throughput (Mbps)	Mean intra-MAN data packet tx. delay (ms)	Mean inter-MAN data packet tx. delay (ms)
102.87	195.08	297.95	2.27	5.04
127.32	167.77	295.09	2.55	5.60
153.04	139.84	292.88	3.23	6.96
178.64	112.01	290.65	4.14	8.78
202.96	85.47	288.43	5.70	11.89
230.16	55.65	285.81	8.90	18.30
255.54	27.71	283.25	19.81	40.11
270.03	9.25	279.28	36.94	74.38

Table 10. Data packet transmission delay and throughput achieved for different voice throughput levels

With the assumptions and restrictions of our simulation model, the backbone MAN network can accommodate approximately 185 active voice calls per node for 64 kbps voice coding rate with an acceptable 2% discarded voice packets for a 25% inter-MAN traffic level network. This loading represents 9250 active voice calls per MAN and 55500 active voice calls per six-MAN network. For a 50% inter-MAN traffic level network, the system can accommodate approximately 175 active voice calls per node, 8750 per MAN and 52500 per six-MAN network. For a 75% inter-MAN traffic, the system can accommodate approximately 160 active voice calls per node,

8000 per MAN and 48000 per six-MAN network. All delays show in Figures 37 to 41, inclusive, are less than 1.0 ms provided that the QA throughput is limited below some maximum value.

## **Section 4.6**

### **Call Setup Delay**

In this section, we compare the call setup delay of isochronous QA and PA traffic. The setup delay for QA access is expected to be smaller than for PA since QA avoids the complexity and the delay for the Bandwidth Manager ( BWM ) and VCI to allocate valid PA slots. Moreover, QA access for isochronous traffic eliminates the call clearing process required to release PA bandwidth reserved for voice calls.

#### **4.6.1 Signalling protocol**

Below, we explore the possibility of using CCITT Q.931 as the call setup protocol for the MAN-based PCN network. The choice of CCITT Q.931 as the signalling protocol is based on its adoption for the ISDN User-Network Interface ( UNI ). Because this signalling protocol and call control procedures are designed for wireline networks with centralized switching to enable exchange of control signals between fixed points, it is suitable for implementation in the connection-oriented MAN. In CCITT Q.931, call establishment is initiated when the caller sends a SETUP message across the UNI. Upon receipt of the SETUP message, the network returns a SETUP ACKNOWLEDGE message to the caller. The caller then sends any remaining call information in an INFORMATION message. After receiving this message, the network sends a CALL PROCEEDING message to the caller. Upon receiving ALERTING indication initiated at the called address, the network sends an ALERTING message to the caller. The network then sends a CONNECT message to the caller upon receiving a CONNECT indication from the

called party. Call setup is completed after the caller sends a CONNECT ACKNOWLEDGE message to the network [10].

If PA slots are used to transfer isochronous traffic, the network's Signalling Termination ( ST ), Bandwidth Manager & VCI Serve ( NM ) and HOB start to exchange information for allocating an isochronous channel for the caller after the ST receives the SETUP message from the caller. The ST will then send the SETUP ACKNOWLEDGE message back to the caller if and only if an isochronous channel is successfully allocated for that caller. The allocation of an isochronous channel for the called party is similar to that for the caller, except that the request for an isochronous bandwidth is initiated by the ST only after a channel has been allocated for the caller. The ST and NM could be located anywhere on the bus [10]. Figure 42 shows the call setup procedures between the caller and callee.

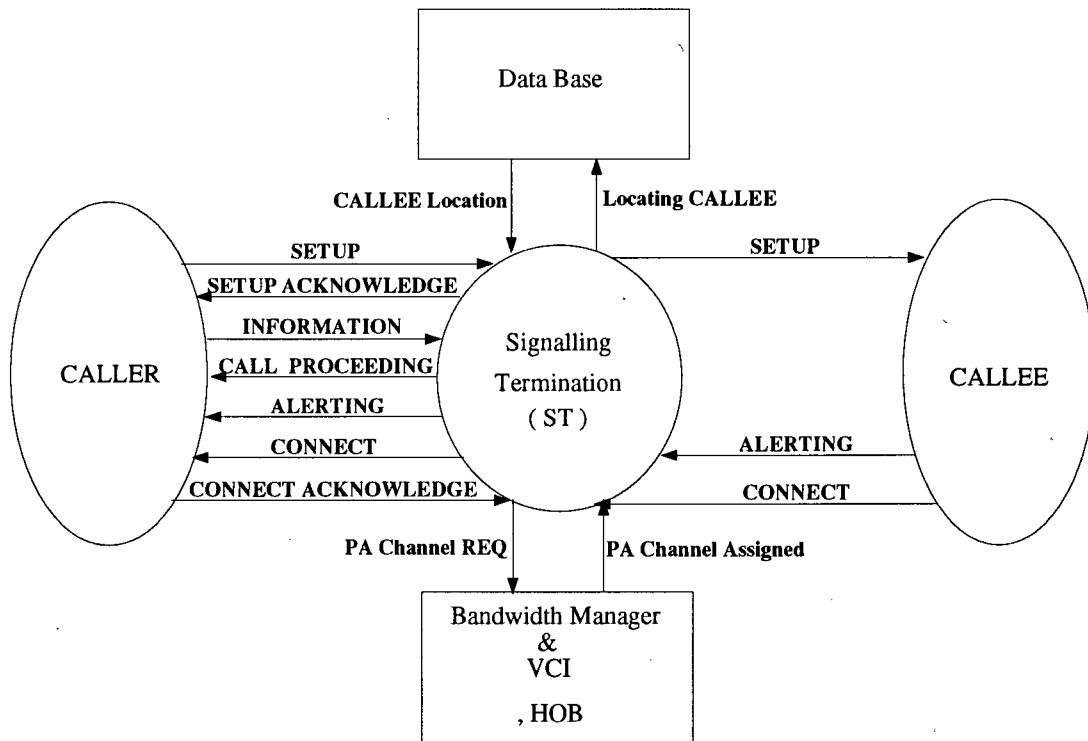


Figure 42. Call setup procedure of Q.931 protocol for PA access. For QA access, signalling involving the Bandwidth Manager and VCI is eliminated.

If QA slots are used to transfer isochronous traffic, the delay and complexity of PA channel allocations for the caller and callee during call setup can be eliminated.

#### **4.6.2 Call setup procedure**

From [2], we know that an overlapping call setup procedure provides a lower mean call setup delay than a sequential call setup procedure. The overlapping call setup procedure reduces the setup time by processing isochronous channel setup and the second information phase at the same time ( see [2] for more details ). Therefore, we expect that reduction in call setup delay for QA access for isochronous traffic over PA access will be greater with sequential call setup than in an overlapping setup procedure.

#### **4.6.3 Performance analysis**

As before, the network configuration consists of 50 nodes uniformly distributed in each access MAN, with two successive nodes separated by 3 slots transmission time. Mean nodal processing delay is  $100 \mu s$  for each signalling message. The length of the SETUP message is assumed to be 5 slots and that of other signalling messages is 1 slot. Signalling messages are transferred in the highest QA access priority ( i.e. priority 2 ) for the network.

Results in Table 11 and 12 are obtained from the simulation results of the best signalling architecture ( which collocates the HOB and NM ) in [2] for mean call setup delay<sup>7</sup> versus the background PA throughput. To show the effects of limited QA capacity for signalling traffic, we have assumed the significant bandwidth ( i.e. more than 200 Mbps ) is unavailable for QA access, but is instead used for PA access. Otherwise, all signalling traffic delay for QA access is constant and, in all cases, less than 0.2 ms.

---

<sup>7</sup> Call setup delay represents the duration from the time the caller initiates the call to the time the SETUP message reaches the callee. If PA access is used, allocation of PA channel for caller is included in the delay.

Background PA throughput (Mbps)	Mean call setup delay (ms)		
	PA access isochronous traffic		QA access isochronous traffic
	Sequential	Overlapping	
204.0	7.72	4.95	4.03
240.8	8.62	5.52	4.48
256.6	8.88	5.68	4.61
273.8	9.52	6.00	4.83

Table 11. Mean inter-MAN call setup delay vs. background PA throughput over 2 transitional MANs

Background PA throughput (Mbps)	Mean call setup delay (ms)		
	PA access isochronous traffic		QA access isochronous traffic
	Sequential	Overlapping	
204.0	3.19	2.88	2.27
240.8	3.53	3.18	2.49
256.6	3.62	3.27	2.56
273.8	3.84	3.45	2.67

Table 12. Mean intra-MAN call setup delay vs. background PA throughput

As expected, the mean QA access call setup delay for isochronous traffic is smaller than the mean PA access setup delay regardless of the call setup procedures used. Since the mean delay for QA access of isochronous packets increases as the traffic throughput increases, it is wise to implement QA access for isochronous traffic under light traffic load and to implement PA access for isochronous traffic under heavy traffic. PA access guarantees bounded delay by denying admission to any packets for which isochronous slots are not available.

# **Chapter 5. Summary and Conclusions**

To handle the increasing communication traffic loads which result from the rapid growth in demand for Personal Communications Services, a distributed switching network architecture based on the IEEE 802.6 MAN with a distributed queue dual bus ( DQDB ) protocol is proposed. This distributed network has several potential advantages relative to current centralized network architectures for supporting PCS. Performance analysis of the proposed network is needed. Accordingly, network delay versus throughput performance for mixed data and real-time voice traffic is determined using computer simulations.

## **Section 5.1 Results Summary**

The efficiency of the DQDB protocol is analyzed first by comparing its delay-vs-throughput performance with and without slot reuse. Two slot reuse schemes, *PSI* and *Destination Release*, are employed for the comparison. Simulation results indicate that performance with slot reuse is not much improved relative to non-reuse when implemented under light traffic load. Slot reuse schemes provide some reduction in packet delay under heavy traffic load, but do not provide any significant improvement in throughput. The added implementation complexity associated with slot reuse does not seem worthwhile when DQDB is employed.

Packet transmission delay for inter- and intra-MAN integrated voice and data traffic was evaluated in detail, using Queue-Arbitrated ( QA ) access only. QA slots with priority 1 access ( the middle level ) are used for transporting voice traffic in our simulation models. QA slots with the lowest priority access ( level zero ) are used to transport data traffic and QA slots with the highest priority ( level two ) would be used to transport control signalling. Simulation results

prove that the waiting delay of a packet is negligible compared to the propagation delay of the packet. Since QA access is used for packets transfer, packet transmission delay for each priority level will depend not only on its own priority level traffic, but also on higher priority level traffic.

The homogeneous bridge schemes employed for MAN interconnection affect inter-MAN packet transmission delay. The four-bridge connection scheme proposed preserves the same packet transmission delay characteristics for inter-MAN traffic as for intra-MAN traffic. As voice throughput increases from 103 to 270 Mbps, the total throughput ( voice plus data ) varies from 298 to 279 Mbps with a two percent bound on dropped voice packets. Mean inter- and intra-MAN delays vary from 5.0 and 2.3 ms, respectively, at 103 Mbps voice throughput to 74.4 and 36.9 ms at 270 Mbps voice throughput.

Call setup and clearing delays are compared for Queue-Arbitrated Access Isochronous Traffic ( QAAIT ) and Pre-Arbitrated Access Isochronous Traffic ( PAAIT ). QAAIT call setup excludes the complex and time consuming process of isochronous channel allocation, and is simpler and faster than for PAAIT. Call clearing involves release of the isochronous channel after call completion, and can be omitted for QAAIT. Although QAAIT has faster call setup and clearing than PAAIT, use of QA access is most appropriate under light traffic loading. Use of PA access under heavy traffic loading is recommended, to provide a bounded packet transmission delay by limiting channel access. QA access results in delays which may become unbounded as traffic levels increases.

## **Section 5.2 Suggestions for Future Work**

There is much future work which would be of interest. Of particular interest are the following issues:

- Performance Analysis for Different Homogeneous Bridges Connection Schemes:

In order to optimize the performance of the Backbone MAN network architecture, it may be useful to investigate further the Backbone MAN network performance by relocating the homogeneous bridges for the interconnected MANs or by altering the physical location of the interconnected MANs within the backbone MAN cluster.

- Performance Analysis for Interworking Between IEEE 802.6 MAN and Heterogeneous Networks:

Since a DQDB MAN network acts as an information exchange media for its subnetworks, analysis in interworking between a DQDB MAN network and its subnetworks is essential. For example, it would be useful to investigate the possibility of using ATM network as an interconnection network between interconnected DQDB MANs or using DQDB MAN network as an interconnection network for interconnected ATM networks. It would be useful as well to examine the performance of using heterogeneous bridges for connecting DQDB MANs with different subnetworks ( e.g. LAN, PBX ).

- Performance and Feasibility of Coexisting PA and QA Traffic:

It would be of interest to examine the feasibility and performance advantages of using QA access for all traffic types under light loading, while employing PA access for real-time traffic

*Chapter 5. Summary and Conclusions*

under heavy traffic load. One of the difficulties is the delay which occurs in packing QA or PA slots for low-bit rate voice coding; this delay is 96 ms for 48-byte PA slots at 4 kbps coding rate. One approach to overcoming the delay is to pack more than one call per cell, which could then involve different parts of a cell having different destinations. The complexity increases considerably, and requires thoughtful analysis of all of the relevant issues.

## Appendix A Source Code

Backbone MAN network Simulation Model Source Code is written in  
**SIMSCRIPT II.5**

## Appendix A

```

*****+
* THE PREAMBLE routine gives the description of each modelling element like
* variable type, temporary element, permanent element, process, etc.
* +
*****+
preamble
normally, mode is undefined

processes include STATION.VD.BUSA, STATION.VD.BUSB, PACKET.ARR
processes include VOICE.CALL.ARR, VOICE.PACKET.UPDATE, STOP.SIM
processes include BRIDGE.VD.BUSA, BRIDGE.VD.BUSB, BRIDGE.VD.BUSA
processes include BRIDGE.VD.BUSB, BRIDGE.VD.BUSA, BRIDGE.VD.BUSA
processes include BACKBONE.MAN, BACKBONE.HEAD, VD.BUSA, BACKBONE.HEADS.VD.BUSA,
BACKBONE.NODE.VD.BUSA, BACKBONE.NODE.VD.BUSB

permanent entities
every STRIN own a vbbufferq and a dbufferq
every BRIDGE own a vbbufferq and a dbufferq
every BBNODE own a vbbufferq and a dbbbufferq

temporary entities
every vbbuffer has a vtemp.add, a vdelay.temp, a vdest.man
and belong to the vbbufferq
every dbuffer has a temp.add, a delay.temp, a sdest.man
and belong to the dbufferq
every vbbuffer has a Vadd.bridge, a Vdest.MAN,
every vbbuffer has a Vadd.bridge, a Vdelay.bridge, a dest.MAN,
every vbbuffer has a add.bridge, a delay.bridge, a dest.MAN,
ini.man, a bdelay, a ini.add and belong to the dbbufferq
every vbbuffer has a add.bridge, a delay.bridge, a dest.MAN,
ini.man, a bdelay, a ini.add and belong to the vbbufferq
every vbbuffer has a vadd.bb, a vdelay.bb, a vdest.MAN.bb,
a Vini.MAN.bb, a bdelay.bb, a vbbufferq
every dbbuffer has a add.bb, a delay.bb, a dest.MAN.bb,
a ini.MAN.bb, a bdelay.bb and belong to the dbbufferq
a ini.add.bb and belong to the dbbufferq

define II, LL, I, spf, man.noAA, Vman as integer variables // global index
define man,noAA, man,noBB, man,noCB, man,noCA, man,noBC as integer variables
define man,noAA, man,noSB, man,noPS, man,noVS as integer variables
define N.stn as integer variable // number of stations
define SD as real variable // distance between two successive stations, slots
define num.slot as integer variable // number of slots/nodes for simulation
define Poisson.arr as real variable // poisson arrival rate of stations
define slot.add.A, slot.add.B as 2-dimensional, integer array
define B.A, B.B as 2-dimensional, integer array // busy bits of slots
define tx.counter.A, vx.counter.A as 2-dimensional, integer array
define tx.counter.B, vx.counter.B as 2-dimensional, integer array
define req.counter.A, vreq.counter.B as 2-dimensional, integer array
define tx.counter.BA, vtx.counter.BA as 2-dimensional, integer array
define req.counter.BA, vreq.counter.BA as 2-dimensional, integer array
define num.b, drnum.b, drnum.bb, vnum.b as integer variables
define DR.A, DR.B as 2-dimensional, integer array // request bits of voice slots
define MAN.IA, MAN.IB, MAN2.A, MAN2.B as 2-dimensional, integer array
define inter.A, inter.B as 2-dimensional, integer array
define voice.A, voice.B as 2-dimensional, integer array
define delay.A, delay.B as 2-dimensional, real array
define delay, delay as 4-dimensional, real array
define delay, vdelay as real variables
define temp.add, vtemp.add as integer variables
define dest.man,Vdest.man as integer variables
define Vdelay.bridge, delay.bridge as real variables
define ini.man, Vini.man as integer variables
define dest.man, Vdest.man as integer variables

define Vini.add, ini.add, Vini.add.bb, ini.add.bb as integer variables
define dbufferlimit, vbufferlimit, vbufferlimit as integer variables
define dbufferlimit, vbufferlimit as integer variables
define dbufsize, dbufsize as 2-dimensional, integer array
define Vrx.add, Vrx.man.tx.add, tx.man as 2-dimensional, integer array
define Vtx.add, Vtx.man, btx.add, btx.man as 2-dimensional, integer array
define END.TIME as real variable
define node, voice.call.num, voice.tx, data.tx, pac.arr as integer variables
define slot.dur as real variable
define d.inter, man.count as 1-dimensional, integer array
define bpac.enter, bpac.total as integer variables
define BBMAN as integer variable
define BB.slot.add.A, BB.slot.add.B as 1-dimensional, integer array
define BB.ini.node.A, BB.ini.node.B as 1-dimensional, integer array
define BB.BA, BB.BB as 1-dimensional, integer array
define BB.DR.A, BB.DR.B as 1-dimensional, integer array
define BB.VN.RA, BB.VN.RB as 1-dimensional, integer array
define BB.VN1.A, BB.VN1.B as 1-dimensional, integer array
define BB.MAN1.A, BB.MAN1.B as 1-dimensional, integer array
define BB.MAN2.A, BB.MAN2.B as 1-dimensional, integer array
define BB.inter.A, BB.inter.B as 1-dimensional, integer array
define BB.voice.A, BB.voice.B as 1-dimensional, integer array
define BB.delay.A, BB.delay.B as 1-dimensional, real array
define BB.tx.counter.A, BB.tx.counter.B as 1-dimensional, integer array
define BB.tx.counter.BA, BB.tx.counter.BB as 1-dimensional, integer array
define BB.req.counter.A, BB.req.counter.B as 1-dimensional, integer array
define BB.vreq.counter.A, BB.vreq.counter.B as 1-dimensional, integer array
define BB.vdelay, bbdelay, add as 1-dimensional, integer array
define Vnum.bb, drnum.bb as 1-dimensional, integer array
define Vdelay.tx, BB.data.tx as integer variables
define Vadd.bb, Vdest.MAN.bb, Vini.MAN.bb as integer variables
define add.bb, dest.MAN.bb, ini.MAN.bb as integer variables
define bbdelay, bbdelay, vdelay, vdelay.bb as real variables
define vbbufsize, ddbbufsize as 1-dimensional, integer array
define BB.vdelay, BB.ddelay as 1-dimensional, real array

define bridge.A, bridge.B as 2-dimensional, integer array
define ini.node.A, ini.node.B as 2-dimensional, real array
define via.bridge as 3-dimensional, real array
define num.via.bridge as 3-dimensional, integer array

define Bv to mean 384 // num of bits in voice packet, 48 octets
define Rv to mean 32000 // voice digitized rate, bps
define BP to mean 424 // num of bits in a slot
define RN to mean 155000000 // network tx rate/bus, 155Mbps
define .secs to mean DAYS
define .us to mean MINUTES

define percent.traffic, percent.inter.traffic.v as integer variables
define sn.act as integer variable
define vtn.act as integer variable
define vtn.act to mean HOURS
define .us to mean MINUTES

end // preamble
*****+
* The function of MAIN routine is to input variables parameters into the
* system, initialize, set up the configuration of the system and schedule
* the beginning time of each process.
*****+
main
define m, n, k, s as integer variables // indexs
let HOURS.V=1000
let MINUTES.V=1000

```

## *Appendix A*

```

let I=1
let II=1
let Lij=1
let N=1
let noAA=1
let man.noAB=1
let man.noB=1
let man.noBA=1
let man.noCA=1
let man.noSA=1
let man.noBS=1
let man.noPS=1
let man.noTS=1
let scn.act=0
let TIME.V=0.0
let voice.tx=0
let data.tx=0
let pac.arr=0

print 1 line thus
Enter the number of nodes (exclude head & end stations):
read N,scn

print 1 line thus
Enter the separation distance 2 successive nodes in slots:
read SD

print 1 line thus
Enter the Poisson arrival rate of node, packets/sec:
read Poisson.arr

print 1 line thus
Enter the number of buffers per node:
read num.buffer

print 1 line thus
Enter the number of voice stations activated per node:
read Vstn.act

print 1 line thus
Enter the time for simulations, in secs:
read END.TIME

print 1 line thus
Enter the percentage of inter-MAN traffic for data traffic, in %:
read data.percentage.inter.traffic

print 1 line thus
Enter the percentage of inter-MAN traffic for voice traffic, in %:
read voice.percentage.inter.traffic

let slot.dur=BP/RN
let num.slot=ttrunc.F(END.TIME/slot.dur)
let sp=(N.scn+)*SD+5
let NSTATNE=N.scn+2
let N_BRIDGE=6^3
let N_BBNODE=6
let dbufferlimit=num.buffer
let vbufferlimit=vstn.act
let dbufferlimit=num.buffer
let vbufferlimit=num.buffer

create slot.add.A(*,*), slot.add.B(*,*), B.A(*,*), B.B(*,*), MAN1.A(*,*),
MAN2.A(*,*), DR.B(*,*), DR.B(*,*), VR.B(*,*), MAN2.A(*,*),
MAN2.B(*,*), inter.A(*,*), inter.B(*,*), voice.A(*,*), voice.B(*,*),
delay.B(*,*), delay.B(*,*), bridge.A(*,*), bridge.B(*,*),
ini.node.B(*,*), ini.node.B(*,*), in.node.B(*,*), as 6 by spf
reserve tx.counter.(*,*,*) req.counter.(*,*,*) Vtx.counter.(*,*,*)
Vtx.counter.B(*,*), tx.counter.B(*,*), reg.counter.B(*,*),
reg.counter.B(*,*), tx.counter.B(*,*), reg.counter.B(*,*),
Vtx.counter.B(*,*), reg.counter.B(*,*), as 6 by N.scn+2 by Vstn.act
create every STATION
create every BRIDGE
create every BNODE

reserve slot.add.A(*,*), slot.add.B(*,*), B.A(*,*), B.B(*,*), MAN1.A(*,*),
MAN2.A(*,*), DR.B(*,*), DR.B(*,*), VR.B(*,*), MAN2.A(*,*),
MAN2.B(*,*), inter.A(*,*), inter.B(*,*), voice.A(*,*), voice.B(*,*),
delay.B(*,*), delay.B(*,*), bridge.A(*,*), bridge.B(*,*),
ini.node.B(*,*), ini.node.B(*,*), in.node.B(*,*), as 6 by spf
reserve tx.counter.(*,*,*) req.counter.(*,*,*) Vtx.counter.(*,*,*)
Vtx.counter.B(*,*), tx.counter.B(*,*), reg.counter.B(*,*),
reg.counter.B(*,*), tx.counter.B(*,*), reg.counter.B(*,*),
Vtx.counter.B(*,*), reg.counter.B(*,*), as 6 by N.scn+2 by 2
```

## *Appendix A*

## Appendix A

```

loop
end '' VOICE.CALL.ARR
*****+
*   * The VOICE.PACKET.UPDATE routine is use to generate voice packet according
*   * to the voice coding rate of the system during a call which is in talkspurt.
*   * This routine also alternate the voice calls between talkspurt and silence
*   * for an exponentially period of time.
*   *
process VOICE.PACKET.UPDATE
define stn,no, MAN, num.call as integer variables
define J as integer variables
define P, END.TIME.VPU as real variable
let stn,no=none
let MAN=Vman
let num.call=voice.call.num
let END.TIME.VPU=END.TIME*stn,no*SD*slot.dur
do J=1 to Vstn.act*0.4
  with num.call=J
  find the first case
  if none
    go to 'SILENT'
  always
    while TIME.V < END.TIME.VPU
      do P=TIME.V+exponential.f(1,5,4)
        while (TIME.V>P) and TIME.V<END.TIME.VPU
          do
            if vsbufsize(MAN,stn,no) < vsbufferlimit
              add 1 to vsbufsize(MAN,stn,no)
            create a vsbuffer
            let Vtemp.add(vsbuffer)=voice.add.dest(MAN,stn,no,num.call)
            let Vdelay,temp(vsbuffer)=TIME.V
            let Vdest.man(vsbuffer)=voice.man.dest(MAN,stn,no,num.call)
            if vsbufsize(MAN,stn,no)=1
              let Vtx.add(MAN,stn,no)=Vtemp.add(vsbuffer)
            let Vtx.man(MAN,stn,no)=Vdest.man(vsbuffer)
            always
              file this vsbuffer in vsbuffer erg((MAN-1)*50+stn,no)
              add 1 to Vnum(MAN,stn,no)
            always
              if (TIME.V+BV/Rv)>END.TIME.VPU
                wait BV/Rv .secs
              else
                go to 'FINISH'
            always
            loop
          let P=TIME.V+exponential.f(2,25,4)
          if P < END.TIME.VPU
            wait P-TIME.V .secs
          else
            go to 'FINISH'
        always
      loop
    'FINISH' end '' packet.arr
  always
  * The following routines are used to update the transmission address of each
  * packet in each node. One of the routines is used to predetermine the voice
  * calls addresses.
  *
  * The PACKET.ARR routine is responsible to generate a data packet enters to
  * each node according to a Poisson arrival rate. This routine also responds to
  * for determining the destination MAN and node of the packet according to an
  define MAN, stn,no as integer variables

```

```

* uniform distribution and inserting the generated data packet to the nodal
* queue.
*
process PACKET.ARR
define stn,no, s, MAN as integer variables
define P, END.TIME.PACKET as real variable
let stn,no=11
let MAN=man.noPS
add 1 to man.noPS
let END.TIME.PACKET=END.TIME+II*SD*slot.dur
while TIME.V<END.TIME.PACKET
do
  if dsbufsize(MAN,stn,no) < dsbufferlimit
    add 1 to dsbufsize(MAN,stn,no)
    create a dsbuf for
    let delay temp(dsbuffer)=TIME.V
    if percent.inter.traffic<5
      go to 'INTRAS'
    always
      s=randi.f(1,100,5)
      if s<percent.inter.traffic
        let s=randi.f(1,6,7)
        let temp.add(dsbuffer)=randi.f(1,N,stn+2,4)
      do
        let s=randi.f(1,6,4)
      loop
      let sdest.man(dsbuffer)=MAN
      let s=randi.f(1,N,stn+2,4)
      while stn,no>8
        do
          let s=randi.f(1,N,stn+2,4)
        loop
        let sdest.man(dsbuffer)=MAN
        let s=randi.f(1,N,stn+2,4)
        always
          add 1 to man.count(sdest.man(dsbuffer))
          if dsbufsize(MAN,stn,no)=1
            let tx.add(MAN,stn,no)=temp.add(dsbuffer)
            let tx.man(MAN,stn,no)=sdest.man(dsbuffer)
            always
              file this dsbuffer in dsbuffer erg((MAN-1)*50+stn,no)
              add 1 to pac.arr
            always
              let p-TIME.V+exponential.f(1/Poisson.arr,4)
              if p < END.TIME.PACKET
                wait P-TIME.V .secs
              else
                go to 'FINISH'
            always
          loop
        'FINISH' end '' packet.arr
      *
      * The following routines are used to update the transmission address of each
      * packet in each node. One of the routines is used to predetermine the voice
      * calls addresses.
      *
      routine VSTX.ADD.UPDATE Given MAN, stn,no
      define MAN, stn,no as integer variables

```

## Appendix A

```

remove first vbuffer from vbufferq((MAN-1)*50+stn.no)
let Vxx.add(MAN,stn.no)=Vtemp.add(vbuffer)
let Vxx.man(MAN,stn.no)=vdest.man(vbuffer)
file vbuffer First in vbufferq((MAN-1)*50+stn.no)
end '' vtx.add.update
'',,
routine stx.add.update given MAN, stn.no
define MAN, stn.no as integer variables
remove first dsbuffer from dsbufferq((MAN-1)*50+stn.no)
let tx.add(MAN,stn.no)=temp.add(dsbuffer)
let tx.man(MAN,stn.no)=dest.man(dsbuffer)
file dsbuffer First in dsbufferq((MAN-1)*50+stn.no)
end '' stx.add.update
'',,
routine vtx.add.update given MAN, bridge.no
define MAN, bridge.no as integer variables
remove first vbuffer from vbufferq((MAN-1)*3+bridge.no)
let Vbtx.add(MAN,bridge.no)=yadd.bridge(vbuffer)
let Vbtx.man(MAN,bridge.no)=vdest.bridge(vbuffer)
file vbuffer First in vbufferq((MAN-1)*3+bridge.no)
end '' vtx.add.update
'',,
routine btx.add.update given MAN, bridge.no
define MAN, bridge.no as integer variables
remove first dbbuffer from dbbufferq((MAN-1)*3+bridge.no)
let btx.add(MAN,bridge.no)=add.bridge(dbbuffer)
let btx.man(MAN,bridge.no)=dest.man(dbbuffer)
file dbbuffer First in dbbufferq((MAN-1)*3+bridge.no)
end '' btx.add.update
'',,
routine vbbtx.add.update given stn.no
define stn.no as integer variable
remove first vbbbuffer from vbbbufferq(stn.no)
if stn.no=1
    if vdest.MAN.bb(vbbbuffer)=3
        let vbbtx.add(stn.no)=3
    else
        if vdest.MAN.bb(vbbbuffer)=4
            let vbbtx.add(stn.no)=4
        else
            if vdest.MAN.bb(vbbbuffer)=5
                let vbbtx.add(stn.no)=5
            always
                always
            else
                if stn.no=2
                    if vdest.MAN.bb(vbbbuffer)=2
                        let vbbtx.add(stn.no)=2
                    else
                        if vdest.MAN.bb(vbbbuffer)=4
                            let vbbtx.add(stn.no)=4
                        else
                            let vbbtx.add(stn.no)=4
                always
                always
            end '' vbbtx.add.update
        end '' vbbtx.add.update
    end '' vbbtx.add.update
else
    if stn.no=3
        if vdest.MAN.bb(vbbbuffer)=3
            let vbbtx.add(stn.no)=3
        else
            if vdest.MAN.bb(vbbbuffer)=4
                let vbbtx.add(stn.no)=4
            else
                let vbbtx.add(stn.no)=4
        always
        always
    end '' vbbtx.add.update
end '' vbbtx.add.update
file vbbbuffer first in vbbbufferq(stn.no)
routine vbbtx.add.update given stn.no
define stn.no as integer variable
remove first vbbbuffer from vbbbufferq(stn.no)
if stn.no=1
    if vdest.MAN.bb(vbbbuffer)=3
        let vbbtx.add(stn.no)=3
    else
        if vdest.MAN.bb(vbbbuffer)=4
            let vbbtx.add(stn.no)=4
        else
            if vdest.MAN.bb(vbbbuffer)=5
                let vbbtx.add(stn.no)=5
            always
            always
        end '' vbbtx.add.update
    end '' vbbtx.add.update
else
    if stn.no=2
        if vdest.MAN.bb(vbbbuffer)=2
            let vbbtx.add(stn.no)=2
        else
            if vdest.MAN.bb(vbbbuffer)=4
                let vbbtx.add(stn.no)=4
            else
                let vbbtx.add(stn.no)=4
        always
        always
    end '' vbbtx.add.update
end '' vbbtx.add.update
routine vbbtx.add.update given stn.no
define stn.no as integer variable
remove first vbbbuffer from vbbbufferq(stn.no)
if stn.no=1
    if vdest.MAN.bb(vbbbuffer)=3
        let vbbtx.add(stn.no)=3
    else
        if vdest.MAN.bb(vbbbuffer)=4
            let vbbtx.add(stn.no)=4
        else
            if vdest.MAN.bb(vbbbuffer)=5
                let vbbtx.add(stn.no)=5
            always
            always
        end '' vbbtx.add.update
    end '' vbbtx.add.update
else
    if stn.no=2
        if vdest.MAN.bb(vbbbuffer)=2
            let vbbtx.add(stn.no)=2
        else
            if vdest.MAN.bb(vbbbuffer)=4
                let vbbtx.add(stn.no)=4
            else
                let vbbtx.add(stn.no)=4
        always
        always
    end '' vbbtx.add.update
end '' vbbtx.add.update

```

## *Appendix A*

```

    ''
    routine voice.call.predet
    define k, n, m, s, c, JJ, trial, a, b, x as integer variables
    '',
    ' uniformly chosen the voice calls for inter Man
    ,

    if percent.inter.traffic.v>5
    while inter(traffic.f(vstn.act*percent.inter.traffic.v/100))=0
    do
        let k=randi .f(1,vstn.act/4)
        for m=1 to trunc.f(vstn.act*percent.inter.traffic.v/100)
        with k=inter(n)
        find the first case
        if none
        for n=1 to trunc.f(vstn.act*percent.inter.traffic.v/100)
        with inter(n)=0
        find the first case
        if found
            let inter(n)=k
            always
        always
        loop
        always

        '',
        ' pre-determine the inter v.c. destinations
        ,

        if percent.inter.traffic.v>5
        for k=1 to 6 '' by percent.inter.traffic
        do
            for n=1 to N str+n2
            do
                for m= 1 to vstn.act
                do
                    let trial=1
                    if voice.add.dest(k,n,m)=0
                    for x=1 to trunc.f(vstn.act*percent.inter.traffic.v/100)
                    with voice.add.dest(m,int(x))
                    find the first case
                    if found
                        let s=randi .f(1,6,8)
                        let c=randi .f(1,N,str+n2,4)
                        REPEAT
                        let voice.add.dest(s,b,int(c))
                        while k=s
                        do
                            let s=randi .f(1,6,8)
                        loop
                        for JJ= 1 to trunc.f(vstn.act*percent.inter.traffic.v/100)
                        with voice.add.dest(s,c,int(er(JJ)))=0
                        find the first case
                        if found
                            let voice.add.dest(k,n,m)=c
                            let voice.add.dest(s,b,int(er(JJ)))=n
                            let voice.man.dest(k,n,m=s
                            let voice.man.dest(s,b,int(er(JJ)) )=k
                            let voice.add.dest(k,n,m)=b
                            if trial>20
                                for s back from 6 to 1
                                with s>x to trunc.f(vstn.act*percent.inter.traffic.v/100)
                                for a1 to trunc.f(vstn.act*percent.inter.traffic.v/100)
                                for b back from N.str+n2 to 1
                                with voice.add.dest(s,b,int(a))=0
                                find the first case
                                if found
                                    let voice.man.dest(k,n,m)=s
                                    let voice.man.dest(s,b,int(a))=k
                                    let trial=1
                                else
                                    always
                                    add 1 to trial
                            else
                                always

```

## Appendix A

```

        go to 'REPEAT'
    always
    always
    always
    loop
    loop
    always
    for m= 1 to vstn.act          '' pre-determine the intra v.c. destinations
    do
        for k1 to 6 '' by percent. inter.traffic
        do
            if voice.add.dest(k,n,m)=0
            do
                'REPEAT' let trial=1
                let s=randi.f(1,N,stn+2,4)
                let d=vstn.act
                if voice.add.dest(k,n,m)=0
                do
                    if voice.add.dest(k,s,JJ)=0
                    find the first case
                    if found
                    let voice.add.dest(k,n,m)=s
                    let voice.add.dest(k,s,JJ)=n
                    let voice.man.dest(k,n,m)=k
                    let voice.man.dest(k,s,JJ)=k
                    let trial=1
                else
                    for JJ= 1 to vstn.act
                    with voice.add.dest(k,s,JJ)=0
                    find the first case
                    if found
                    for a=1 to vstn.act
                    for b back from N.stn+2 to 1
                    with voice.add.dest(k,b,a)=0 and b>n
                    find the first case
                    if found
                    let voice.add.dest(k,n,m)=b
                    let voice.add.dest(k,b,a)=n
                    let voice.man.dest(k,n,m)=k
                    let voice.man.dest(k,b,a)=k
                    let trial=1
                else
                    add 1 to trial
                    go to 'REPEAT'
                always
            always
            loop
            loop
            for m=1 to trunc.f(vstn.act*0.4*percent.inter.traffic.v/100)
            do
                let d(m)=inter(m)
            loop
            while d(vstn.act*0.4)=0          '' uniformly chosen the v.c. in talk period.
            do
                let x=randi.f(1,vstn.act,4)
                for n=1 to trunc.f(vstn.act*percent.inter.traffic.v/100)
                with x=inter(n)
                find the first case
                if none
                    for n= 1 to vstn.act*0.4
                    with d(n)
                    find the first case
                    if none
                        for n=1 to vstn.act*0.4
                        with d(n)=0
                        if found
                            let d(n)=x
                always
            always
        always
    always

```

```

        always
        always
        loop
        end '' voice.call.predict
        ****
        *
        *   The following routines combine to form the backbone MAN.
        *
        process BACKBONE.MAN
        ****
        define m as integer variable
        define star as real variable
        let BBman=1
        let BB.voice.tx=0
        let BB.dctc.tx=0
        let star=(N.stn+3) * SD+0.5+6/15
        for m1 to 6
        do
            let BB.tx.counter.A(m1)=1
            let BB.VTX.counter.A(m1)=1
            let BB.tx.counter.B(m1)=1
            let BB.VTX.counter.B(m1)=1
        loop
        activate a BACKBONE.HEADA.VD.BUSA in star*slot.dur .secs
        activate a BACKBONE.HEADB.VD.BUSB in (star+0.5)*slot.dur .secs
        for m1 to 4
        do
            activate a BACKBONE.NODE.VD.BUSA in (star+4*60)*slot.dur .secs
            activate a BACKBONE.NODE.VD.BUSB in (star+0.5*m+60)*slot.dur .secs
        loop
        activate a BACKBONE.HEADB.VD.BUSA in (star+5*60)*slot.dur .secs
        activate a BACKBONE.HEADA.VD.BUSA in (star+0.5+m+60)*slot.dur .secs
        end '' BACKBONE.MAN
        /*
        /*
        /*
        process BACKBONE.HEADA.VD.BUSA
        ****
        *
        * this process describes the homogeneous bridge at the head station,
        * refer to bus A, towards bus B in the BACKBONE MAN.
        * Slots travel along bus A are started From this bridge.
        */
        */
        */
        define stn.no, slotA, frameLoop as integer variables
        define END.TIME,BBHA as real variables
        let stn.no=1
        let slotA=1
        let frameLoop=1
        let END.TIME,BBHA=END.TIME+(N.stn+3)*SD+0.5+6/15)*slot.dur
        while TIME.V < END.TIME,BBHA
        do
            if slotA>spf/2
                let slotA=1
                add 1 to frameLoop
            always
            let BB.slot.add.A(slotA)=0
            let BB.B.A(slotA)=0
            let BB.DR.A(slotA)=0

```

## *Appendix A*

## *Appendix A*

## *Appendix A*

```

*****+
* this process describe the action of the homogeneous bridges located at the
* backbone MAN rather than two heads towards bus A.
*****+
define stn.no, slotA, Tx, frame.loop as integer variables
let stn.no=BBman+1
let slotA=1
let frame.loop=1

while TIME.V < END.TIME+ ( N. stn+3 ) *SD-0 .5+6/15*BBman*60 ) *slot.dur
do if slotA>spf/2
    let slotA=1
    add 1 to frame.loop
    always
        if BB.VTX.counter.A(slotA).no=0           ' TRANSMISSION FOR VOICE PACKET
            remove first vbbuffer from vbbufferfqr(stn.no)
            let BB.slot.add(A(slotA))=add.bb(vbbuffer)
            let BB.ini.node(A(slotA))=Vini.add.bb(vbbuffer)
            let BB.MAN1.A(slotA)=Vini.MAN.bb(vbbuffer)
            let BB.MAN2.A(slotA)=Vdest.MAN.bb(vbbuffer)
            let BB.inter.A(slotA)=2                   ' reset tx.counter
            let BB.voice.A(slotA)=2                  ' set busy bit
            let BB.delay.A(slotA)=delay.bb(vbbuffer)
            let BB.Vtx.counter(A(stn.no))=1
            let BB.B(A(slotA))=1
        if frame.loop>5
            let BB.vdelay(stn.no)BB.vdelay(stn.no)+TIME.V-bbvdelay(vbbuffer)
            add 1 to vnum.bb.vdelay(stn.no)
            add 1 to BB.voice.tx
        always
            destroy this vbbuffer
            subtract 1 from vbbuffersize(stn.no)
            if vbbufferfqr(stn.no) is not empty
                call vbbtx.add.update giving stn.no
            else
                let vbbtx.add(stn.no)=0
            always
                let Tx=
            else if BB.Vtx.counter.A(stn.no)>0
                subtract 1 from BB.Vtx.counter.A(stn.no)
            else
                if BB.Vreq.counter.A(stn.no)>0 and BB.Vtx.counter.A(stn.no)<0
                    subtract 1 from BB.Vreq.counter.A(stn.no)
                else
                    if BB.tx.counter.A(stn.no)=0           ' TX FOR DATA PACKET
                        remove first dbbbuffer from dbbbufferfqr(stn.no)
                        let BB.slot.add(A(slotA))=add.bb(dbbbuffer)
                        let BB.ini.node(A(slotA))=ini.add.bb(dbbbuffer)
                        let BB.MAN1.A(slotA)=ini.MAN.bb(dbbbuffer)
                        let BB.MAN2.A(slotA)=dest.MAN.bb(dbbbuffer)
                        let BB.inter.A(slotA)=2
                        let BB.voice.A(slotA)=1
                        let BB.delay.A(slotA)=delay.bb(dbbbuffer)
                        let BB.tx.counter(A(stn.no))-1      ' reset tx.counter
                        let BB.B(A(slotA))=1
                    if frame.loop>5
                        let BB.ddelay(stn.no)=BB.ddelay(stn.no)+TIME.V-bbvdelay(dbbbuffer)
                    always
                        destroy this dbbbuffer
                        subtract 1 from dbbbufferfqr(stn.no)
                        if dbbbufferfqr(stn.no) is not empty
                            call bbttx.add.update giving stn.no
                        else
                            let BBTX.add(stn.no)=0
                            let Tx=

```

```

let BB.MAN1.B(slotB).Vini.MAN.bbb(vbbuffer)
let BB.MAN2.B(slotB).Vdest.MAN.bb(vbbuffer)
let BB.inter.B(slotB)=2
let BB.voice.B(slotB)=2
let BB.delay.B(slotB)=Vdelay.bb(vbbuffer)      // reset tx counter
let BB.BB.delay.B(slotB)=1                         // set busy bit
let BB.Vtx.Counter.B(stn.no)=1
let BB.BB.Vtx.Counter.B(stn.no)=1
add 1 to vnum.bb(stn.no)
add 1 to BB.voice.CX

always
destroy this vbbuffer
subtract 1 from vbbuffersize(stn.no)
if vbbuffer(stn.no) is not empty
call vbbtx.add.update giving stn.no
else
    let Vbbtx.add(stn.no)=0
always

else
if BB.Vreq.counter.B(stn.no)>0
    subtract 1 from BB.Vreq.counter.B(stn.no)
else
if BB.Vreq.counter.B(stn.no)>0 and BB.Vtx.counter.B(stn.no)<0
    subtract 1 from BB.Vreq.counter.B(stn.no)
else
    if BB.tx.counter.B(stn.no)=0           //TX FOR DATA PACKET
        remove first dbbbuffer from dbbbufferf(stn.no)
        let BB.slot.add.B(slotB)=add.bb(dbbbuffer)
        let BB.ini.node.B(slotB)=ini.add.bb(dbbbuffer)
        let BB.MAN.B(slotB)=ini.MAN.bb(dbbbuffer)
        let BB.MAN.B(slotB)=dest.MAN.bb(dbbbuffer)
        let BB.inter.B(slotB)=2
        let BB.voice.B(slotB)=2
        let BB.BB.inter.B(slotB)=1
        let BB.BB.delay.B(slotB)=delay.bb(dbbbuffer)
        let BB.tx.counter.B(stn.no)=1      // reset tx counter
        let BB.B.B(slotB)=1
        if frame.loop>5
            BB.ddelay(stn.no)=BB.ddelay(stn.no)+TIME.V-bbdddelay(dbbbuffer)
            add 1 to dnum.bb(stn.no)
            add 1 to BB.data.CX
        always
            destroy this dbbbuffer
            subtract 1 from dbbbuffersize(stn.no)
            if dbbbuffer(stn.no) is not empty
                call bbtx.add.update giving stn.no
            else
                let bbtx.add(stn.no)=0
        always
            if BB.tx.counter.B(stn.no)>0
                subtract 1 from BB.tx.counter.B(stn.no)
            else
                if BB.req.counter.B(stn.no)>0 and BB.tx.counter.B(stn.no)<0
                    subtract 1 from BB.req.counter.B(stn.no)
                always
                    always
                    always
                    always
                    add 1 to slotB
                    wait slot.dur .secs
                loop
            end // BACKBONE.HEADB.VD.BUSA
        '''
        '''
        '''
process BACKBONE.NODE.VD.BUSA

```

## *Appendix A*

## *Appendix A*

## *Appendix A*

```

add 1 to bvpac.enter
always
` add 1 to bvpac.total
else
  if BB.voice.B(slotB)=1
    if dbbsize(BB.MAN2.B(slotB),1) < vbbufferlimit
      add 1 to dbbsize(BB.MAN2.B(slotB),1)
      create a dbbuffer
      let add_bridge(dbbuffer)=BB.slot.add.B(slotB)
      let ini_bridge(dbbuffer)=BB.ini.node.B(slotB)
      let delay_bridge(dbbuffer)=BB.delay.B(slotB)
      let ini_MAN(dbbuffer)=BB.MAN1.B(slotB)
      let bddelay(dbbuffer)=TIME.V
      if dbbsize(BB.MAN2.B(slotB),1)=1
        let btx.add(BB.MAN2.B(slotB),1)=add_bridge(dbbuffer)
        let btx.man(BB.MAN2.B(slotB),1)=dest.MAN(dbbuffer)
        always
          file this dbbuffer in dbbufferfq((BB.MAN2.B(slotB)-1)*3+1)
          add 1 to bdpac.enter
        always
          add 1 to bdpac.total
        always
        else
          if str.no=4 and BB.MAN2.B(slotB)=4
            if BB.voice.B(slotB)=2
              if dbbsize(BB.MAN2.B(slotB),1) < vbbufferlimit
                add 1 to vbbusise(BB.MAN2.B(slotB),1)
                create a vbbusise
                let add_bridge(vbbuffer)=BB.slot.add.B(slotB)
                let ini_VIN.add(vbbuffer)=BB.ini.node.B(slotB)
                let VDECT.MAN(vbbuffer)=BB.MAN1.B(slotB)
                let VINI.MAN(vbbuffer)=BB.MAN1.B(slotB)
                let VDECT.MAN(vbbuffer)=BB.MAN2.B(slotB)
                let VINI.MAN(vbbuffer)=BB.MAN2.B(slotB)
                let VDECT.MAN(vbbuffer)=TIME.V
                if vbbusise(BB.MAN2.B(slotB),1)=1
                  let bddelay(vbbuffer)=TIME.V
                  if vbbusise(BB.MAN2.B(slotB),1)=1
                    let btx.add(BB.MAN2.B(slotB),1)=Vadd.bridge(vbbuffer)
                    let btx.man(BB.MAN2.B(slotB),1)=Vdeset.MAN(vbbuffer)
                    always
                      file this vbbusise in vbbusisefq((BB.MAN2.B(slotB)-1)*3+1)
                      add 1 to bvpac.enter
                    always
                      add 1 to bvpac.total
                    always
                    else
                      if BB.voice.B(slotB)=1
                        if dbbsize(BB.MAN2.B(slotB),1) < vbbufferlimit
                          add 1 to dbbsize(BB.MAN2.B(slotB),1)
                          create a dbbuffer
                          let add_bridge(dbbuffer)=BB.slot.add.B(slotB)
                          let ini_bridge(dbbuffer)=BB.ini.node.B(slotB)
                          let delay_bridge(dbbuffer)=BB.delay.B(slotB)
                          let ini_MAN(dbbuffer)=BB.MAN1.B(slotB)
                          let dest_MAN(dbbuffer)=BB.MAN2.B(slotB)
                          let bddelay(dbbuffer)=TIME.V
                          if dbbsize(BB.MAN2.B(slotB),1)=1
                            let btx.add(BB.MAN2.B(slotB),1)=add_bridge(dbbuffer)
                            let btx.man(BB.MAN2.B(slotB),1)=dest.MAN(dbbuffer)
                            always
                              file this dbbuffer in dbbufferfq((BB.MAN2.B(slotB)-1)*3+1)
                              add 1 to bdpac.enter
                            always
                              add 1 to bdpac.total
                            always
                            else
                              if str.no=5 and BB.MAN2.B(slotB)=6
                                if BB.voice.B(slotB)=2
                                  if dbbsize(BB.MAN2.B(slotB),1) < vbbufferlimit
                                    add 1 to vbbusise(BB.MAN2.B(slotB),1)
                                    create a vbbusise
                                    let add_bridge(vbbuffer)=BB.slot.add.B(slotB)
                                    let ini_VIN.add(vbbuffer)=BB.ini.node.B(slotB)
                                    let VDECT.MAN(vbbuffer)=BB.MAN1.B(slotB)
                                    let VINI.MAN(vbbuffer)=BB.MAN2.B(slotB)
                                    let VDECT.MAN(vbbuffer)=TIME.V
                                    if vbbusise(BB.MAN2.B(slotB),1)=1
                                      let btx.add(BB.MAN2.B(slotB),1)=Vadd.bridge(vbbuffer)
                                      let btx.man(BB.MAN2.B(slotB),1)=Vdeset.MAN(vbbuffer)
                                      always
                                        file this vbbusise in vbbusisefq((BB.MAN2.B(slotB)-1)*3+1)

```

## *Appendix A*

## *Appendix A*

## *Appendix A*

## *Appendix A*

```

let dest.MAN.bb(vbbuffer) = MAN2.B(MAN, slotB)
let bdelay.v(vbbuffer) = TIME.V
file this vbbuffer in dbbuffer(frq1b)
if vbbufsize(bb)=1
  call vbbxx.add.update giving b
    always
      add 1 to bdpac.enter
    always
      add 1 to bdpac.total
    always
      add 1 to slotB
    always
      wait slot.dur . secs
    loop
      add 1 to slotB
      wait slot.dur . secs
      go to 'finish'
do
  go to 'req'
  if B(MAN,slotB)=1 and (MAN2.B(MAN,slotB)=2 or
    'regones' if B(MAN,slotB)=4)
    MAN2.B(MAN,slotB)=1
    let b=2
    if voice.B(MAN, slotB)=2
      if vbbufsize(bb) < vbbufferrlimit
        add 1 to vbbuffer.size (b)
        create a vbbuffer
        let Vadd.bb(vbbuffer) = slot.add.B(MAN, slotB)
        let Vini.add.bb(vbbuffer) = slot.ini.node.B(MAN, slotB)
        let Vdelay.bb(vbbuffer) = slot.delay.B(MAN, slotB)
        let Vini.MAN.bb(vbbuffer) = slot.ini.MAN1.B(MAN, slotB)
        let Vdest.MAN.bb(vbbuffer) = slot.dest.MAN2.B(MAN, slotB)
        let bdelay.bb(vbbuffer) = TIME.V
        file this vbbuffer in vbbuffer(frq1b)
        if vbbufsize(bb)=1
          call vbbxx.add.update giving b
            always
              add 1 to bvpac.enter
            always
              add 1 to bvpac.total
            always
              add 1 to slotB
            always
              wait slot.dur . secs
            loop
              go to 'finish'
do
  go to 'req'
  if voice.B(MAN,slotB)=1
    if vbbufsize(bb) < vbbufferrlimit
      add 1 to vbbuffer.size (b)
      create a vbbuffer
      let add.bb(vbbuffer) = slot.add.B(MAN, slotB)
      let ini.add.bb(vbbuffer) = slot.ini.node.B(MAN, slotB)
      let delay.bb(vbbuffer) = slot.delay.B(MAN, slotB)
      let ini.MAN.bb(vbbuffer) = slot.ini.MAN1.B(MAN, slotB)
      let dest.MAN.bb(vbbuffer) = slot.dest.MAN2.B(MAN, slotB)
      let bdelay(bb(vbbuffer)) = TIME.V
      file this vbbuffer in dbbuffer(frq1b)
      if vbbufsize(bb)=1
        call vbbxx.add.update giving b
          always
            add 1 to bdpac.enter
          always
            add 1 to bdpac.total
          always
            add 1 to slotB
          always
            wait slot.dur . secs
          loop
            go to 'finish'
do
  go to 'req'
  if B(MAN,slotB)=1 and (MAN2.B(MAN,slotB)=1 or
    'regones' if B(MAN,slotB)=4)
      add 1 to slotB
      wait slot.dur . secs
      go to 'finish'

```

```

MAN2.B(MAN,slotB)=3
let b=5
if voice.B(MAN,slotB)=2
  if vbbbuffer.size(b) < vbbbuffer.limit
    add 1 to vbbbuffer.size(b)
    create a vbbbuffer
    let vbb.(b)(vbbbuffer)=slot.add.B(MAN,slotB)
    let Vani.ddd.bb(vbbbuffer)=ini.node.B(MAN,slotB)
    let Vaneiy.bb.(vbbbuffer)=delay.B(MAN,slotB)
    let Vani.MAN.bb.(vbbbuffer)=MAN1.B(MAN,slotB)
    let Vdest.MAN.bb.(vbbbuffer)=MAN2.B(MAN,slotB)
    let bdelay.(vbbbuffer)=TIME.V
    file this vbbbuffer in vbbbufferq.(b)
    if vbbbuffer.size(b)=1
      call vbbcx.add.update giving b
    always
      add 1 to bvpac.enter
    always
      add 1 to bvpac.total
  else
    if voice.B(MAN,slotB)=1
      if dbbbuffer.size(b) < dbbbuffer.limit
        add 1 to dbbbuffer.size(b)
        create a dbbbuffer
        let ini.add.bb.(dbbbuffer)=slot.add.B(MAN,slotB)
        let ini.add.bb.(dbbbuffer)=ini.node.B(MAN,slotB)
        let delay.bb.(dbbbuffer)=delay.B(MAN,slotB)
        let dest.MAN.bb.(dbbbuffer)=MAN1.B(MAN,slotB)
        let dest.MAN.bb.(dbbbuffer)=MAN2.B(MAN,slotB)
        let bdelay.(dbbbuffer)=TIME.V
        file this dbbbuffer in dbbbufferq.(b)
        if dbbbuffer.size(b)=1
          call bbtx.add.update giving b
        always
          add 1 to bdpac.enter
        always
          add 1 to bdpac.total
        always
          add 1 to slotB
          wait slot.dur .secs
        always
          add 1 to slotB
          wait slot.dur .secs
        loop
      'finish' end // BRIDGEA.VD.BUSB
      /*
      ****
      * The BRIDGEB routines represent the homogeneous bridges that located at the
      * head of another bus for connecting two access MANS together.
      */
      ****
process BRIDGEB.VD.BUSA

define bridge.no. slotA, MAN as integer variables
define END.TIME.BRIDGEB as real variable

let MAN=man,noBA
let bridge.no-3
let slotA=1
add 1 to man,noBA

let END.TIME.BRIDGEB=END.TIME+( (N..strn+3) *SD+(MAN-1)/15 ) *slot.dur

```

## *Appendix A*

```

if MAN=1
  go to 'first'
else
  if MAN=2
    go to 'second'
  else
    if MAN=3
      go to 'third'
    else
      if MAN=4
        go to 'forth'
      else
        if MAN=5
          go to 'fifth'
        else
          go to 'sixth'
    always
    'reg'
    let slotA=1
  always
  if vbtX.add(MAN,bridge.no)>0 and vrx.counter.BB(MAN,bridge.no)<0
    let vrx.counter.BB(MAN,bridge.no)=vreq.counter.BB(MAN,bridge.no)
  always
  if vrx.add(MAN,bridge.no)=0 and tx.counter.BB(MAN,bridge.no)<0
    let tx.counter.BB(MAN,bridge.no)=req.counter.BB(MAN,bridge.no)
  always
  if vr.A(MAN,slotA)=1
    add 1 to vreq.counter.BB(MAN,bridge.no)
  always
  if DR.A(MAN,slotA)=1
    add 1 to reg.counter.BB(MAN,bridge.no)
  always
  if VR.A(MAN,slotA)=1
    add 1 to req.counter.BB(MAN,bridge.no)
  always
  if MAN=1
    go to 'reqdone1'
  else
    if MAN=2
      go to 'reqdone2'
    else
      if MAN=3
        go to 'reqdone3'
      else
        if MAN=4
          go to 'reqdone4'
        else
          if MAN=5
            go to 'reqdone5'
          else
            go to 'reqdone6'
  always
  'first' while TIME.V < END.TIME.BRIDGEBA
  do
    go to 'req'
    'reqdone1' if B.A(MAN,slotA)=1 and MAN2.A(MAN,slotA)=5
    if vbufsize(MAN2.A(MAN,slotA),2)< vbufufferlimit
      add 1 to vbufsize(MAN2.A(MAN,slotA),2)
    create a vbufuffer
    let vadd.bridge(vbufuffer)=slot.add.A(MAN,slotA)
    let ini.add(vbufuffer)=ini.node.A(MAN,slotA)
    let delay.bridge(vbufuffer)=delay.A(MAN,slotA)
    let ini.MAN(vbufuffer)=MAN1.A(MAN,slotA)
    let dest.MAN(vbufuffer)=MAN2.A(MAN,slotA)
    let delay.V(vbufuffer)=TIME.V
    if vbufsize(MAN2.A(MAN,slotA),2)=1
      let vbx.add(MAN2.A(MAN,slotA),2)=add.bridge(dbbuffer)
    let vbx.man(MAN2.A(MAN,slotA),2)=dest.MAN(dbbuffer)
    always
    file this vbufuffer in vbufifferg((MAN2.A(MAN,slotA)-1)*3+2)
    add 1 to bvpac.enter
    always
    add 1 to bvpac.total
  always
  'reqdone2' if B.A(MAN,slotA)=1 and MAN2.A(MAN,slotA)=6
  if voice.A(MAN,slotA)=2
    if vbufsize(MAN2.A(MAN,slotA),2)< vbufufferlimit
      add 1 to vbufsize(MAN2.A(MAN,slotA),2)
    create a vbufuffer
    let vadd.bridge(vbufuffer)=slot.add.A(MAN,slotA)
    let ini.add(vbufuffer)=ini.node.A(MAN,slotA)
    let delay.bridge(vbufuffer)=delay.A(MAN,slotA)
    let ini.MAN(vbufuffer)=MAN1.A(MAN,slotA)
    let dest.MAN(vbufuffer)=MAN2.A(MAN,slotA)
    let delay.V(vbufuffer)=TIME.V
    if vbufsize(MAN2.A(MAN,slotA),2)=1
      let vbx.add(MAN2.A(MAN,slotA),2)=add.bridge(dbbuffer)
    let vbx.man(MAN2.A(MAN,slotA),2)=dest.MAN(dbbuffer)
    always
    file this vbufuffer in vbufifferg((MAN2.A(MAN,slotA)-1)*3+2)
    add 1 to bvpac.enter
    always
    add 1 to bvpac.total
  always
  'reqdone3' if voice.A(MAN,slotA)=1
  if vbufsize(MAN2.A(MAN,slotA),2)< vbufufferlimit
    add 1 to vbufsize(MAN2.A(MAN,slotA),2)
  create a vbufuffer
  let vadd.bridge(vbufuffer)=slot.add.A(MAN,slotA)
  let ini.add(vbufuffer)=ini.node.A(MAN,slotA)
  let delay.bridge(vbufuffer)=delay.A(MAN,slotA)
  let ini.MAN(vbufuffer)=MAN1.A(MAN,slotA)
  let dest.MAN(vbufuffer)=MAN2.A(MAN,slotA)
  let delay.V(vbufuffer)=TIME.V
  if vbufsize(MAN2.A(MAN,slotA),2)=1
    let vbx.add(MAN2.A(MAN,slotA),2)=add.bridge(dbbuffer)
  let vbx.man(MAN2.A(MAN,slotA),2)=dest.MAN(dbbuffer)
  always
  file this vbufuffer in vbufifferg((MAN2.A(MAN,slotA)-1)*3+2)
  add 1 to bvpac.enter
  always
  add 1 to bvpac.total
  always
  'reqdone4' if voice.A(MAN,slotA)=1
  if vbufsize(MAN2.A(MAN,slotA),2)< vbufufferlimit
    add 1 to vbufsize(MAN2.A(MAN,slotA),2)
  create a vbufuffer
  let vadd.bridge(vbufuffer)=slot.add.A(MAN,slotA)
  let ini.add(vbufuffer)=ini.node.A(MAN,slotA)
  let delay.bridge(vbufuffer)=delay.A(MAN,slotA)
  let ini.V(vbufuffer)=MAN1.A(MAN,slotA)
  let dest.V(vbufuffer)=TIME.V
  if vbufsize(MAN2.A(MAN,slotA),2)=1
    let vbx.add(MAN2.A(MAN,slotA),2)=add.bridge(dbbuffer)
  let vbx.man(MAN2.A(MAN,slotA),2)=dest.MAN(dbbuffer)
  always
  file this vbufuffer in vbufifferg((MAN2.A(MAN,slotA)-1)*3+2)
  add 1 to bvpac.enter
  always
  add 1 to bvpac.total
  always
  'reqdone5' if voice.A(MAN,slotA)=1
  if vbufsize(MAN2.A(MAN,slotA),2)< vbufufferlimit
    add 1 to vbufsize(MAN2.A(MAN,slotA),2)
  create a vbufuffer
  let vadd.bridge(vbufuffer)=slot.add.A(MAN,slotA)
  let ini.add(vbufuffer)=ini.node.A(MAN,slotA)
  let delay.bridge(vbufuffer)=delay.A(MAN,slotA)
  let ini.V(vbufuffer)=MAN1.A(MAN,slotA)
  let dest.V(vbufuffer)=TIME.V
  if vbufsize(MAN2.A(MAN,slotA),2)=1
    let vbx.add(MAN2.A(MAN,slotA),2)=add.bridge(dbbuffer)
  let vbx.man(MAN2.A(MAN,slotA),2)=dest.MAN(dbbuffer)
  always
  file this vbufuffer in vbufifferg((MAN2.A(MAN,slotA)-1)*3+2)
  add 1 to bvpac.enter
  always
  add 1 to bvpac.total
  always
  'reqdone6' if voice.A(MAN,slotA)=1
  if vbufsize(MAN2.A(MAN,slotA),2)< vbufufferlimit
    add 1 to vbufsize(MAN2.A(MAN,slotA),2)
  create a vbufuffer
  let vadd.bridge(vbufuffer)=slot.add.A(MAN,slotA)
  let ini.add(vbufuffer)=ini.node.A(MAN,slotA)
  let delay.bridge(vbufuffer)=delay.A(MAN,slotA)
  let ini.V(vbufuffer)=MAN1.A(MAN,slotA)
  let dest.V(vbufuffer)=TIME.V
  if vbufsize(MAN2.A(MAN,slotA),2)=1
    let vbx.add(MAN2.A(MAN,slotA),2)=add.bridge(dbbuffer)
  let vbx.man(MAN2.A(MAN,slotA),2)=dest.MAN(dbbuffer)
  always
  file this vbufuffer in vbufifferg((MAN2.A(MAN,slotA)-1)*3+2)
  add 1 to bvpac.enter
  always
  add 1 to bvpac.total
  always

```

## *Appendix A*

## Appendix A

```

always
    add 1 to bdpac.total
    always
        always
            go to 'sixth' while TIME.V < END.TIME.BRIDGEBA
do
    go to 'req'
    if B.A(MAN,slotA)=1 and (MAN2.A(MAN,slotA)=4 or
        MAN2.A(MAN,slotA)=5)
        if voice.A(MAN,slotA)=2
            if vbufsize(MAN2.A(MAN,slotA),3) < vbufbuffer.limit
                add 1 to vbufx
                vbufx.add(MAN2.A(MAN,slotA),3)=vdest.MAN(vbufbuffer)
                let Vadd.bridge(vbufbuffer)=slot.add.A(MAN,slotA)
                let Vini.add(vbufbuffer)=ini.node.(MAN,slotA)
                let Vdelay.bridge(vbufbuffer)=delay.A(MAN,slotA)
                let Vini.MAN(vbufbuffer)=MAN1.A(MAN,slotA)
                let Vdest.MAN(vbufbuffer)=MAN2.A(MAN,slotA)
                let bdelay.(vbufbuffer)=TIME.V
                if vbufsize(MAN2.A(MAN,slotA),3)=1
                    let vbufx.add(MAN2.A(MAN,slotA),3)=vadd.bridge(vbufbuffer)
                    let vbufx.man(MAN2.A(MAN,slotA),3)=vdest.MAN(vbufbuffer)
                    always
                        file this vbufbuffer in vbufbufferq!(MAN2.A(MAN,slotA)-1)*3+3)
                    add 1 to bvpac.enter
                    always
                        add 1 to bvpac.total
                    else
                        if voice.A(MAN,slotA)=1
                            if dbbufsize(MAN2.A(MAN,slotA),3) < dbbufbuffer.limit
                                add 1 to dbbufx
                                dbbufx.add(MAN2.A(MAN,slotA),3)=slot.add.A(MAN,slotA)
                                create a dbbuf for
                                let add.bridge(dbbuf)=slot.add.A(MAN,slotA)
                                let ini.add(dbbuf)=ini.node.A(MAN,slotA)
                                let delay.bridge(dbbuf)=delay.A(MAN,slotA)
                                let ini.MAN(dbbuf)=MAN1.A(MAN,slotA)
                                let dest.MAN(dbbuf)=MAN2.A(MAN,slotA)
                                let bdelay.(dbbuf)=TIME.V
                                if dbbufsize(MAN2.A(MAN,slotA),3)=1
                                    let bx.x.add(MAN2.A(MAN,slotA),3)=add.bridge(dbbuf)
                                    let bx.x.man(MAN2.A(MAN,slotA),3)=dest.MAN(dbbuf)
                                    always
                                        file this dbbuf in dbbufbufferq!(MAN2.A(MAN,slotA)-1)*3+3)
                                    add 1 to bdpac.enter
                                    always
                                        add 1 to bdpac.total
                                    always
                                        always
                                            add 1 to slotA
                                            wait slot.dur .secs
                                        loop
                                        'finish' end // BRIDGE.BD.BUSA
                                        //'
                                        //'
                                        //'
                                        process BRIDGE.BD.BUSB
                                        ****
                                        * this process describes the homogeneous bridge as the head station,
                                        * refer to bus B, towards bus B. Slots travel along bus B are started
                                        * from this bridge.
                                        *

```

```

*****+
define bridge.no, slotB, MAN, frame.loop as integer variables
let MAN=man.nobb
let bridge.no=3
let slots=1
let frame.loop=1
add 1 to man.nobb

loop
while TIME.V < END.TIME+(0.5+(MAN-1)/15)*slot.dur
do
    if slotB>spf
        let slotB=1 // initialize frames & slots
        add 1 to frame.loop
    always
        let slot.add.B(MAN,slotB)=0
        let B.B(MAN,slotB)=0
        let DR.B(MAN,slotB)=0
        let VR.B(MAN,slotB)=0
        let MAN1.B(MAN,slotB)=0
        let MAN2.B(MAN,slotB)=0
        let inter.B(MAN,slotB)=0.
        let voice.B(MAN,slotB)=0
        let delay.B(MAN,slotB)=0
        let bridge.B(MAN,slotB)=0
        let ini.node.B(MAN,slotB)=ini.MAN(vbufbuffer)
        if Vtx.counter.BB(MAN,bridge.no)=0 // TRANSMISSION FOR VOICE PACKET
            remove first vbufbuffer from vbufbufferq!(MAN-1)*3+bridge.no)
            let slot.add.B(MAN,slotB)=vadd.bridge(vbufbuffer)
            let ini.node.B(MAN,slotB)=vini.MAN(vbufbuffer)
            let MAN1.B(MAN,slotB)=dest.MAN(vbufbuffer)
            let MAN2.B(MAN,slotB)=vdest.MAN(vbufbuffer)
            let inter.B(MAN,slotB)=2
            let bridge.B(MAN,slotB)=bridge.no
            let voice.B(MAN,slotB)=0
            let delay.B(MAN,slotB)=Vdelay(MAN,bridge.no)+TIME.V-bdelay(vbufbuffer)
            if frame.loop>5
                let Vtx.counter.BB(MAN,bridge.no)=1
                let B.B(MAN,slotB)=1 // set busy bit
                if vdelay(MAN,bridge.no)=vdelay(MAN,bridge.no)+TIME.V-bdelay(vbufbuffer)
                    add 1 to vnum.b(MAN,bridge.no)
                    add 1 to voice.tx
                    always
                        destroy this vbufbuffer
                        subtract 1 from vbufsize(MAN,bridge.no)
                        if vbufbufferq!(MAN-1)*3+bridge.no) is not empty
                            call Vtx.counter.BB(MAN,bridge.no).update giving MAN, bridge.no
                        else
                            let Vtx.add(MAN,bridge.no)=0
                            let Vtx.man(MAN,bridge.no)=0
                            always
                                if Vtx.counter.BB(MAN,bridge.no)>0
                                    subtract 1 from Vtx.counter.BB(MAN,bridge.no)
                                else
                                    if Vreq.counter.BB(MAN,bridge.no)>0 and Vtx.counter.BB(MAN,bridge.no)<0
                                        subtract 1 from Vreq.counter.BB(MAN,bridge.no)
                                    else
                                        if tx.counter.BB(MAN,bridge.no)=0
                                            remove first dbbuf from dbbufq!(MAN-1)*3+bridge.no)
                                            let slot.add.B(MAN,slotB)=add.bridge(dbbuf)
                                            let ini.node.B(MAN,slotB)=ini.add(dbbuf)
                                            let MAN1.B(MAN,slotB)=ini.MAN(dbbuf)
                                            let MAN2.B(MAN,slotB)=dest.MAN(dbbuf)
                                            let inter.B(MAN,slotB)=2
                                            let voice.B(MAN,slotB)=bridge.no
                                            let bridge.B(MAN,slotB)=bridge.no
                                            let delay.B(MAN,slotB)=delay.bridge(dbbuf)
                                            let tx.counter.BB(MAN,bridge.no)=-1 // reset tx.counter
                                            let B.B(MAN,slotB)=1 // set busy bit
                                            if frame.loop>5
                                                let delay.(MAN,bridge.no)=drum.b(MAN,bridge.no)+TIME.V-bdelay(dbbuf)
                                                add 1 to data.tx
                                            
```

## *Appendix A*

## Appendix A

```

always
  add 1 to bopac.total
  if voice.A(MAN, slotA)=1
    if dbbufsize(MAN2.A(MAN, slotA), 2) < dbbufferlimit
      add 1 to dbbufsize(MAN2.A(MAN, slotA), 2)
    let add_a_dbbuffer
    create_a_dbbuffer
    let add_bridge(dbbuffer)=slot, add_A(MAN, slotA)
    let ini, add(dbbuffer)=ini, node, A(MAN, slotA)
    let delay_bridge(dbbuffer)=delay, A(MAN, slotA)
    let ini, MAN(dbbuffer)=MAN1, A(MAN, slotA)
    let dest, MAN(dbbuffer)=MAN2, A(MAN, slotA)
    let bopac.total
    always
      add 1 to bopac.total
    else
      if voice.A(MAN, slotA)=1
        if dbbufsize(MAN2.A(MAN, slotA), 2)=1
          let btx.add(MAN2.A(MAN, slotA), 2)=add_bridge(dbbuffer)
          let btx.man(MAN2.A(MAN, slotA), 2)=dest.MAN(dbbuffer)
        always
          file this dbbuffer in dbbufferq((MAN2.A(MAN, slotA)-1)*3+2)
          add 1 to bopac.enter
        always
          add 1 to bopac.total
        else
          if voice.A(MAN, slotA)=1
            if dbbufsize(MAN2.A(MAN, slotA), 2)=1
              let btx.add(MAN2.A(MAN, slotA), 2)=add_bridge(dbbuffer)
              let btx.man(MAN2.A(MAN, slotA), 2)=dest.MAN(dbbuffer)
            always
              file this dbbuffer in dbbufferq((MAN2.A(MAN, slotA)-1)*3+2)
              add 1 to bopac.enter
            always
              add 1 to bopac.total
            else
              if voice.A(MAN, slotA)=3
                if dbbufsize(MAN2.A(MAN, slotA), 2) < dbbufferlimit
                  add 1 to dbbufsize(MAN2.A(MAN, slotA), 2)
                let add_a_dbbuffer
                create_a_dbbuffer
                let add_bridge(dbbuffer)=slot, add_A(MAN, slotA)
                let ini, add(dbbuffer)=ini, node, A(MAN, slotA)
                let delay_bridge(dbbuffer)=delay, A(MAN, slotA)
                let ini, MAN(dbbuffer)=MAN1, A(MAN, slotA)
                let dest, MAN(dbbuffer)=MAN2, A(MAN, slotA)
                let bopac.total
                always
                  add 1 to bopac.total
                else
                  if voice.A(MAN, slotA)=3
                    if dbbufsize(MAN2.A(MAN, slotA), 2) < dbbufferlimit
                      add 1 to dbbufsize(MAN2.A(MAN, slotA), 2)
                    let add_a_dbbuffer
                    create_a_dbbuffer
                    let add_bridge(dbbuffer)=slot, add_A(MAN, slotA)
                    let ini, add(dbbuffer)=ini, node, A(MAN, slotA)
                    let delay_bridge(dbbuffer)=delay, A(MAN, slotA)
                    let ini, MAN(dbbuffer)=MAN1, A(MAN, slotA)
                    let dest, MAN(dbbuffer)=MAN2, A(MAN, slotA)
                    let bopac.total
                    always
                      add 1 to bopac.total
                    else
                      if voice.A(MAN, slotA)=2
                        if dbbufsize(MAN2.A(MAN, slotA), 2)=1
                          let btx.add(MAN2.A(MAN, slotA), 2)=add_bridge(dbbuffer)
                          let btx.man(MAN2.A(MAN, slotA), 2)=dest.MAN(dbbuffer)
                        always
                          file this dbbuffer in dbbufferq((MAN2.A(MAN, slotA)-1)*3+2)
                          add 1 to bopac.enter
                        always
                          add 1 to bopac.total
                        else
                          if voice.A(MAN, slotA)=2
                            if dbbufsize(MAN2.A(MAN, slotA), 2) < dbbufferlimit
                              add 1 to dbbufsize(MAN2.A(MAN, slotA), 2)
                            let add_a_dbbuffer
                            create_a_dbbuffer
                            let add_bridge(dbbuffer)=slot, add_A(MAN, slotA)
                            let ini, add(dbbuffer)=ini, node, A(MAN, slotA)
                            let delay_bridge(dbbuffer)=delay, A(MAN, slotA)
                            let ini, MAN(dbbuffer)=MAN1, A(MAN, slotA)
                            let dest, MAN(dbbuffer)=MAN2, A(MAN, slotA)
                            let bopac.total
                            always
                              add 1 to bopac.total
                            else
                              if voice.A(MAN, slotA)=2
                                if dbbufsize(MAN2.A(MAN, slotA), 2)=1
                                  let btx.add(MAN2.A(MAN, slotA), 2)=add_bridge(dbbuffer)
                                  let btx.man(MAN2.A(MAN, slotA), 2)=dest.MAN(dbbuffer)
                                always
                                  file this dbbuffer in dbbufferq((MAN2.A(MAN, slotA)-1)*3+2)
                                  add 1 to bopac.enter
                                always
                                  add 1 to bopac.total
                                else
                                  always

```

## *Appendix A*

```

always
  if B.B(MAN,slotB)=0
    if Vx.counter.BB(MAN,bridge.no)=0      // TRANSMISSION FOR VOICE PACKET
      remove first vbbuffer from vbbufferq (MAN-1)*3+bridge.no
      let slot.add.B(MAN,slotB)=add.bridge(vbbuffer)
      let ini.node.B(MAN,slotB)=vini.MAN(vbbuffer)
      let MAN1.B(MAN,slotB)=vdest.MAN(vbbuffer)'
      let MAN2.B(MAN,slotB)=vdest.MAN(vbbuffer)
      let voice.B(MAN,slotB)=2
      let inter.B(MAN,slotB)=2
      let bridge.B(MAN,slotB)=bridge.no
      let delay.B(MAN,slotB)=delay.B(MAN,bridge.no)
      let Vtx.counter.BB(MAN,bridge.no)=-1
      let B.B(MAN,slotB)=1      // set busy bit
      if frame.loop>5
        if vdelay(MAN,bridge.no)=delay(MAN,bridge.no)+TIME.V-bdelay(vbbuffer)
          add 1 to vnum.b(MAN,bridge.no)
          add 1 to voice.tx
        always
          destroy this vbbuffer
          subtract 1 from vbbufsize(MAN,bridge.no)
          if vbbufferq(MAN-1)*3+bridge.no) is not empty
            call vtx.add.update giving MAN,bridge.no
          else
            let vbtx.add(MAN,bridge.no)=0
            always
              let Tx=1
            else
              if Vtx.counter.BB(MAN,bridge.no)>0
                subtract 1 from Vreq.counter.BB(MAN,bridge.no)
              else
                if tx.counter.BB(MAN,bridge,no)>0 and vtx.counter.BB(MAN,bridge.no)<0
                  subtract 1 from Vreq.counter.BB(MAN,bridge.no)
                else
                  if tx.counter.BB(MAN,bridge,no)=0
                    ''TX FOR DATA PACKET
                    remove first dbbuffer from dbufferq( (MAN-1)*3+bridge.no)
                    let slot.add.B(MAN,slotB)=add.bridge(dbbuffer)
                    let ini.node.B(MAN,slotB)=ini.add(dbbuffer)
                    let MAN2.B(MAN,slotB)=dest.MAN(dbbuffer)
                    let inter.B(MAN,slotB)=2
                    let voice.B(MAN,slotB)=1
                    let bridge.B(MAN,slotB)=bridge.no
                    let delay.B(MAN,slotB)=delay.bridge(dbbuffer)
                    let tx.counter.BB(MAN,bridge.no)=-1      // reset tx.counter
                    let B.B(MAN,slotB)=1
                    // set busy bit
                    if frame.loop>5
                      if delay(MAN,bridge.no)=ddelay(MAN,bridge,no)+TIME.V-bdelay(dbbuffer)
                        add 1 to data.tx
                      always
                        let btx.add(MAN,bridge,no)=0
                        let btx.man(MAN,bridge.no)=0
                        always
                          let Tx=1
                        else
                          if tx.counter.BB(MAN,bridge,no)>0
                            subtract 1 from tx.counter.BB(MAN,bridge,no)
                          else
                            if req.counter.BB(MAN,bridge,no)>0 and tx.counter.BB(MAN,bridge,no)<0
                              subtract 1 from req.counter.BB(MAN,bridge,no)
                            always
                              always
                                always
                                  always
                                    always
                                      always
                                        always
                                          always
                                            always
                                              always
                                                always
                                                  always
                                                    always
                                                      always
                                                        always
                                                          always
                                                            always
                                                              always
                                                                always
                                                                  always
                                                                    always
                                                                      always
                                                                        always
                                                                          always
                                                                            always
                                                                              always
                                                                                always
                                                                                  always
                                                                                    always
                                                                

```

## Appendix A

```

if dbbufsize(MAN2.B(MAN,slotB),2)=1
  let bxx.add(MAN2.B(MAN,slotB),2)=add_bridge(dbbuffer)
  let bxx.man(MAN2.B(MAN,slotB),2)=dest.MAN(dbbuffer)
  always
    file this dbbuffer in dbbufferq((MAN2.B(MAN,slotB)-1)*3+2)
    always
      add 1 to bdpac.enter
      always
        add 1 to bdpac.total
        always
        else
      always
        if MAN=3 and MAN2.B(MAN,slotB)=4
          if voice.B(MAN,slotB)=2
            if vbbufsize(MAN2.B(MAN,slotB),2)<vbbufferlimit
              add 1 to vbbufsize(MAN2.B(MAN,slotB),2)
              create a vbbuffer
              always
                file this vbbuffer in vbbufferq((MAN2.B(MAN,slotB)-1)*3+2)
                always
                  let Vadd.bridge(vbbuffer)=slot.add.B(MAN,slotB)
                  let Vdelay.bridge(vbbuffer)=ini.node.B(MAN,slotB)
                  let Vini.Man(vbbuffer)=MAN1.B(MAN,slotB)
                  let Vwest.Man(vbbuffer)=MAN2.B(MAN,slotB)
                  let Vpdeelay(vbbuffer)=TIME.V
                  let Vtx.add(MAN2.B(MAN,slotB),2)=add_bridge(vbbuffer)
                  let Vtx.man(MAN2.B(MAN,slotB),2)=dest.MAN(vbbuffer)
                  always
                    file this vbbuffer in vbbufferq((MAN2.B(MAN,slotB)-1)*3+2)
                    add 1 to bdpac.enter
                    always
                      add 1 to bdpac.total
                      always
                      else
                    always
                    if voice.B(MAN,slotB)=1
                      if dbbufsize(MAN2.B(MAN,slotB),2)<dbbufferlimit
                        add 1 to dbbufsize(MAN2.B(MAN,slotB),2)
                        create a dbbuffer
                        let add_bridge(dbbuffer)=slot.add.B(MAN,slotB)
                        let ini.add(dbbuffer)=ini.node.B(MAN,slotB)
                        let dest.MAN(dbbuffer)=MAN1.B(MAN,slotB)
                        let bdelay(dbbuffer)=MAN2.B(MAN,slotB)
                        let bdeelay(dbbuffer)=TIME.V
                        let bbufsize(MAN2.B(MAN,slotB),2)=1
                        let bxx.add(MAN2.B(MAN,slotB),2)=add_bridge(dbbuffer)
                        always
                          file this dbbuffer in dbbufferq((MAN2.B(MAN,slotB)-1)*3+2)
                          always
                            add 1 to bdpac.enter
                            always
                              add 1 to bdpac.total
                              always
                              else
                            always
                            if MAN=4 and MAN2.B(MAN,slotB)=3
                              if voice.B(MAN,slotB)=2
                                if vbbufsize(MAN2.B(MAN,slotB),2)<vbbufferlimit
                                  add 1 to vbbufsize(MAN2.B(MAN,slotB),2)
                                  create a vbbuffer
                                  always
                                    file this vbbuffer in vbbufferq((MAN2.B(MAN,slotB)-1)*3+2)
                                    always
                                      let Vadd.bridge(vbbuffer)=slot.add.B(MAN,slotB)
                                      let Vdelay.bridge(vbbuffer)=ini.node.B(MAN,slotB)
                                      let Vini.Man(vbbuffer)=MAN1.B(MAN,slotB)
                                      let Vwest.Man(vbbuffer)=MAN2.B(MAN,slotB)
                                      let Vpdeelay(vbbuffer)=TIME.V
                                      let Vtx.add(MAN2.B(MAN,slotB),2)=add_bridge(vbbuffer)
                                      let Vtx.man(MAN2.B(MAN,slotB),2)=dest.MAN(vbbuffer)
                                      always
                                        file this vbbuffer in vbbufferq((MAN2.B(MAN,slotB)-1)*3+2)
                                        always
                                          let bxx.add(MAN2.B(MAN,slotB),2)=add_bridge(dbbuffer)
                                          always
                                            file this dbbuffer in dbbufferq((MAN2.B(MAN,slotB)-1)*3+2)
                                            always
                                              add 1 to bdpac.enter
                                              always
                                              else
                                            always
                                            if MAN=5 and MAN2.B(MAN,slotB)=2
                                              if voice.B(MAN,slotB)=2
                                                if vbbufsize(MAN2.B(MAN,slotB),b)<vbbufferlimit
                                                  add 1 to vbbufsize(MAN2.B(MAN,slotB),b)
                                                  create a vbbuffer
                                                  let add_bridge(vbbuffer)=slot.add.B(MAN,slotB)
                                                  let ini.add(dbbuffer)=ini.node.B(MAN,slotB)
                                                  let delay.bridge(vbbuffer)=ini.node.B(MAN,slotB)
                                                  let dest.MAN(vbbuffer)=MAN1.B(MAN,slotB)
                                                  let Vadd.bridge(vbbuffer)=MAN2.B(MAN,slotB)
                                                  let Vdelay.bridge(vbbuffer)=TIME.V
                                                  let Vtx.add(MAN2.B(MAN,slotB),b)=add_bridge(vbbuffer)
                                                  let Vtx.man(MAN2.B(MAN,slotB),b)=dest.MAN(vbbuffer)
                                                  always
                                                    file this vbbuffer in vbbufferq((MAN2.B(MAN,slotB)-1)*3+b)
                                                    add 1 to bdpac.enter
                                                    always
                                                      add 1 to bdpac.total
                                                      always
                                                      else
                                                    always
                                                    if voice.B(MAN,slotB)=1
                                                      if vbbufsize(MAN2.B(MAN,slotB),b)<vbbufferlimit
                                                        add 1 to dbbufsize(MAN2.B(MAN,slotB),b)
                                                        create a dbbuffer
                                                        let add_bridge(dbbuffer)=slot.add.B(MAN,slotB)
                                                        let ini.add(dbbuffer)=ini.node.B(MAN,slotB)
                                                        let delay.bridge(dbbuffer)=ini.node.B(MAN,slotB)
                                                        let dest.MAN(dbbuffer)=MAN1.B(MAN,slotB)
                                                        let Vadd.bridge(vbbuffer)=MAN2.B(MAN,slotB)
                                                        let Vdelay.bridge(vbbuffer)=TIME.V
                                                        let bxx.add(MAN2.B(MAN,slotB),b)=add_bridge(dbbuffer)
                                                        always
                                                          file this dbbuffer in dbbufferq((MAN2.B(MAN,slotB)-1)*3+b)
                                                          always
                                                            add 1 to bdpac.enter
                                                            always
                                                            else
                                                          always
                                                          if MAN=6 and MAN2.B(MAN,slotB)=2
                                                            if vbbufsize(MAN2.B(MAN,slotB),2)<vbbufferlimit
                                                              add 1 to vbbufsize(MAN2.B(MAN,slotB),2)
                                                              create a vbbuffer
                                                              let add_bridge(vbbuffer)=slot.add.B(MAN,slotB)
                                                              let Vadd.bridge(vbbuffer)=ini.node.B(MAN,slotB)
                                                              let Vdelay.bridge(vbbuffer)=ini.node.B(MAN,slotB)
                                                              let Vini.Man(vbbuffer)=MAN1.B(MAN,slotB)
                                                              let Vwest.Man(vbbuffer)=MAN2.B(MAN,slotB)
                                                              let Vpdeelay(vbbuffer)=TIME.V
                                                              let Vtx.add(MAN2.B(MAN,slotB),2)=add_bridge(vbbuffer)
                                                              let Vtx.man(MAN2.B(MAN,slotB),2)=dest.MAN(vbbuffer)
                                                              always
                                                                file this vbbuffer in vbbufferq((MAN2.B(MAN,slotB)-1)*3+2)
                                                                always
                                                                  add 1 to bdpac.total
                                                                  always
                                                                  else
                                                                always
                                                                if voice.B(MAN,slotB)=1

```

## *Appendix A*

## *Appendix A*

## Appendix A

```

go to 'data'
'localD' if stn.no>tx.add(MAN,stn.no)=0
    let DR.A(MAN,stn,no)=req.counter.B(MAN,stn,no)
    let rx.counter.B(MAN,stn,no)=0
    always
        go to 'streqdone'

'interd' if MAN=1
    if stn.no<25 and (tx.man(MAN,stn.no)=3 or
        tx.man(MAN,stn.no)=4 or tx.man(MAN,stn.no)=6)
        let DR.A(MAN,stl)=1
        let rx.counter.B(MAN,stn,no)=req.counter.B(MAN,stn,no)
        let req.counter.B(MAN,stn,no)=0
    else
        if stn.no>25 and (tx.man(MAN,stn.no)=2 or
            tx.man(MAN,stn.no)=3 or tx.man(MAN,stn.no)=4 or
            tx.man(MAN,stn.no)=5)
            let DR.A(MAN,stl)=1
            let rx.counter.B(MAN,stn,no)=req.counter.B(MAN,stn,no)
            let req.counter.B(MAN,stn,no)=0
        else
            if stn.no>25 and (tx.man(MAN,stn.no)=2 or
                tx.man(MAN,stn.no)=3 or tx.man(MAN,stn.no)=4 or
                tx.man(MAN,stn.no)=6)
                let DR.A(MAN,stl)=1
                let rx.counter.B(MAN,stn,no)=req.counter.B(MAN,stn,no)
                let req.counter.B(MAN,stn,no)=0
            always
            else
    else
        if MAN=2
            if stn.no<25 and (tx.man(MAN,stn.no)=3 or
                tx.man(MAN,stn.no)=4 or tx.man(MAN,stn.no)=5)
                let DR.A(MAN,stl)=1
                let rx.counter.B(MAN,stn,no)=req.counter.B(MAN,stn,no)
                let req.counter.B(MAN,stn,no)=0
            else
                if stn.no>25 and (tx.man(MAN,stn.no)=1 or
                    tx.man(MAN,stn.no)=3 or tx.man(MAN,stn.no)=4 or
                    tx.man(MAN,stn.no)=5)
                    let DR.A(MAN,stl)=1
                    let rx.counter.B(MAN,stn,no)=req.counter.B(MAN,stn,no)
                    let req.counter.B(MAN,stn,no)=0
                always
                else
        else
            if MAN=3
                if stn.no<25 and (tx.man(MAN,stn.no)=1 or
                    tx.man(MAN,stn,no)=2 or tx.man(MAN,stn,no)=6)
                    let DR.A(MAN,stl)=1
                    let rx.counter.B(MAN,stn,no)=req.counter.B(MAN,stn,no)
                    let req.counter.B(MAN,stn,no)=0
                else
                    if stn.no>25 and (tx.man(MAN,stn.no)=1 or
                        tx.man(MAN,stn.no)=2 or tx.man(MAN,stn.no)=4 or
                        tx.man(MAN,stn.no)=6)
                        let DR.A(MAN,stl)=1
                        let rx.counter.B(MAN,stn,no)=req.counter.B(MAN,stn,no)
                        let req.counter.B(MAN,stn,no)=0
                    always
                    else
            else
                if MAN=4
                    if stn.no<25 and (tx.man(MAN,stn.no)=1 or
                        tx.man(MAN,stn,no)=2 or tx.man(MAN,stn,no)=6)
                        let DR.A(MAN,stl)=1
                        let rx.counter.B(MAN,stn,no)=req.counter.B(MAN,stn,no)
                        let req.counter.B(MAN,stn,no)=0
                    else
                        if stn.no>25 and (tx.man(MAN,stn.no)=1 or
                            tx.man(MAN,stn.no)=2 or tx.man(MAN,stn.no)=4 or
                            tx.man(MAN,stn.no)=6)
                            let DR.A(MAN,stl)=1
                            let rx.counter.B(MAN,stn,no)=req.counter.B(MAN,stn,no)
                            let req.counter.B(MAN,stn,no)=0
                        always
                        else
                else
                    if MAN=5
                        if stn.no<25 and (tx.man(MAN,stn.no)=1 or
                            tx.man(MAN,stn,no)=2 or tx.man(MAN,stn,no)=6)
                            let DR.A(MAN,stl)=1
                            let rx.counter.B(MAN,stn,no)=req.counter.B(MAN,stn,no)
                            let req.counter.B(MAN,stn,no)=0
                        else
                            if stn.no>25 and (tx.man(MAN,stn.no)=1 or
                                tx.man(MAN,stn.no)=2 or tx.man(MAN,stn.no)=4)
                                let DR.A(MAN,stl)=1
                                let rx.counter.B(MAN,stn,no)=req.counter.B(MAN,stn,no)
                                let req.counter.B(MAN,stn,no)=0
                            always
                            else
                    else
                        if MAN=6
                            if stn.no<25 and (tx.man(MAN,stn,no)=1 or tx.man(MAN,stn,no)=3)
                                let DR.A(MAN,stl)=1
                                let rx.counter.B(MAN,stn,no)=req.counter.B(MAN,stn,no)
                                let req.counter.B(MAN,stn,no)=0
                            else
                                if stn.no>25 and (tx.man(MAN,stn.no)=1 or
                                    tx.man(MAN,stn,no)=2 or tx.man(MAN,stn,no)=5)
                                    let DR.A(MAN,stl)=1
                                    let rx.counter.B(MAN,stn,no)=req.counter.B(MAN,stn,no)
                                    let req.counter.B(MAN,stn,no)=0
                                always
                                always
                                always
                                always
                                always
                                always
                            always
                        always
                    always
                always
            always
        always
    always

```

## *Appendix A*

## Appendix A

```

slot.add.B(MAN,slotB)=stn.no and frame.loop>5
let asMAN1.B(MAN,slotB)
let bInit.node.B(MAN,slotB)
let cInit.B(MAN,slotB)
let dDevice.B(MAN,slotB)
let eBridge.B(MAN,slotB)
let delay.a,b,c,d=delay(a,b,c,d)+TIME.V-delay.B(MAN,slotB)
add 1 to num.re(a,b,c,d)
if c>2 and e>0
let delay.via.bridge(a,e,d)=delay.via.bridge(a,e,d)+TIME.V-delay.B(MAN,slotB)
add 1 to num.via.bridge(a,e,d)
always
always
let throw=0
while vsbufferG(MAN-1)*50+stn.no) is not empty
do
remove first vsbuffer from vsbufferG((MAN-1)*50+stn.no)
let vtime.wait=TIME.V-delay.temp(v$buffer)
if vtime.wait > 0.055
destroy this vsbuffer
subtract 1 from vsbufsize(MAN,stn.no)
add 1 to discard.vpacket(MAN,stn.no)
let throw=1
if vrx.counter.A(MAN,stn.no)>-0
vreq.counter.A(MAN,stn.no)=Vreq.counter.A(MAN,stn.no)
vtx.counter.A(MAN,stn.no)=1
else
if vrx.counter.B(MAN,stn.no)>=0
if vrx.counter.B(MAN,stn.no)+vtx.counter.B(MAN,stn.no)
vreq.counter.B(MAN,stn.no)=Vreq.counter.B(MAN,stn.no)+vtx.counter.B(MAN,stn.no)
vtx.counter.B(MAN,stn.no)=1
always
always
let vrx.add(MAN,stn.no)=0
let vrx.man(MAN,stn.no)=0
else
file vsbuffer first in vsbufferG((MAN-1)*50+stn.no)
if throw=1
call vstx.add.update giving MAN, stn.no
always
always
go to :OUT'
loop
if VR.B(MAN,slotB)=1
add 1 to Vreq.counter.A(MAN,stn.no)
always
if DR.B(MAN,slotB)=1
add 1 to req.counter.A(MAN,stn.no)
always
if VR.B(MAN,slotB)=0 and Vrx.counter.A(MAN,stn.no)<0 and
Vrx.add(MAN,stn.no)>0
if MAN>vtx.man(MAN,stn.no)
go to :localV'
else
go to :interv'
always
if TX=0
if DR.B(MAN,slotB)=0 and tx.counter.A(MAN,stn.no)<0 and
tx.add(MAN,stn.no)>0
if MAN>tx.man(MAN,stn.no)
if MAN>tx.man(MAN,stn.no)
else
go to :localV'
else
go to :interv'
always
if sreqdone' let Tx=0
add 1 to slotB
wait slot.dur .secs
loop

```

## *Appendix A*

## *Appendix A*

## *Appendix A*

```

    // Average Tx. delay of packets via which bridge
    for K=1 to 6
        do for L=1 to 3
            with num.via.bridge(K,L,2)>0
            do for L=1 to 3
                with num.via.bridge(K,L,1)>0
                do let mean.Vdelay.via.bridge[K,L]=delay.via.bridge[K,L,2]/num.via.bridge[K,L,2]
                do let mean.ddelay.via.bridge[K,L]=delay.via.bridge[K,L,1]/num.via.bridge[K,L,1]
                loop
            loop
        loop
    loop

    for K=1 to 6
        do for L=1 to 3
            for K=1 to 6
                with num.Vdelay.via.bridge(K,L)>0.0
                compute avg.Vdelay.via.bridge[L] as the mean of mean.Vdelay.via.bridge[K,L]
            loop
        loop
    loop

    for K=1 to 3
        do for K=1 to 6
            with mean.ddelay.via.bridge(K,L)>0.0
            compute avg.ddelay.via.bridge[L] as the mean of mean.ddelay.via.bridge[K,L]
        loop
    loop

    for L=1 to 3
        do for K=1 to 6
            with mean.vnum.via.bridge(K,L)>0.0
            compute sum.vnum[K] as the sum of vnum(K,L)
        loop
    loop

    for K = 1 to 6
        do for L = 1 to N.srn+2
            compute sum.discard.vpacket[K] as the sum of discard.vpacket[K,L]
        loop

    for K = 1 to 6
        do for L = 1 to N.srn+2
            compute sum.vpacket[K] as the sum of vnum(K,L)
        loop

    for L = 1 to 6
        compute total.discard.vpacket as the sum of sum.discard.vpacket[L]
    loop

    for L = 1 to 6
        compute total.vpacket as the sum of sum.vpacket[L]
    loop

let sum.tx=voice.tx+data.tx
let U = sum.tx/num.slot+spf)/2/6
let U.d = data.tx/(num.slot-5*spf)/2/6
let U.v = voice.tx/(num.slot-5*spf)/2/6*48/53
let BU.v=bridge.vpacket/(num.slot-5*spf)/2/6*48/53
let BU.d=bridge.dpacket/(num.slot-5*spf)/2/6
let pac.arr.rate.Poisson.arr*stn.act*slot.dur
let percent.vpacket.discard=total.discard.vpacket/total.vpacket*100

print 24 lines with (N.strn+2)*6, SD, Poisson.arr, pac.arr.rate, num.slot

```

## Appendix A

```

thus
total # of voice packet enter bridges: *; # of voice packet discarded: *
# voice packet enter buffer: *; # of data packet discarded: *
total # of data packet enter bridges: *; # of data packet discarded: *
# data packet enter buffer: *; # of data packet discarded: *

'
skip 1 output lines
for K=1 to 6
do
print 1 line with K, man.count(K) thus
MAN: *, # of data packets destined:
loop

skip 1 output lines
for L=1 to 3
do
print 1 line with K, L, vbufsize(K,L), dbbufsize(K,L) thus
MAN: *, Bridge: *, v.packet in buffer: *
loop

,
'
for L=1 to 6
do
skip 1 output lines
print 2 lines with L thus
MAN: *
inter
for K=1 to N.stn+2
do
print 1 line with K, meandelay.d.intra(L,K)*1000,
meandelay.v.intra(L,K)*1000, meandelay.d.inter(L,K)*1000,
meandelay.v.inter(L,K)*1000 thus
node: *; d: *.***** ms; v: *.***** ms; d: *.***** ms; v: *.***** ms;
loop

skip 1 output lines
for L=1 to 3
do
print 1 line with L, avg.vdelay.via.bridge(L)*1000,
avg.adelay.via.bridge(L)*1000 thus
Bridge: *, avg.voice.tx.delay: *.***** ms; avg.data.tx.delay: *.***** ms
for K=1 to 6
do
print 1 line with K, mean.vdelay.via.bridge(K,L)*1000,
mean.ddelay.via.bridge(K,L)*1000 thus
MAN: *, avg.voice.tx.delay: *.***** ms; avg.data.tx.delay: *.***** ms
loop

skip 1 output lines
,
'
for K=1 to 6*(N.stn+2)
do
while vbufsize(K) is not empty
do
remove first vbuffer from vbufsize(K)
destroy this vbuffer
loop
while dbbufsize(K) is not empty
do
remove first dbuffer from dbbufsize(K)
destroy this dbuffer
loop
loop
for K=1 to 6*3
do
while vbufsize(K) is not empty
do
remove first vbuffer from vbufsize(K)
destroy this vbuffer

```

## References

- [1] Andrew D. Malyan, Leslie J. Ng, Victor C. M. Leung and Robert W. Donaldson, "Network Architecture and Signalling for Wireless Personal Communications," *IEEE Journal on Selected Areas in Communications*, vol. 11, no. 6, pp. 830–841, August 1993.
- [2] Nanjian Qian, "Call Control Signalling for Personal Communications Over Interconnected Metropolitan Area Network," Master's thesis, UBC Department of Electrical Engineering M.A.Sc. Thesis, April 1993.
- [3] IEEE 802.6 Committee, *Distributed Queue Dual Bus (DQDB) Subnetwork of a Metropolitan Area Network (MAN)*. Institute of Electrical and Electronics Engineers, Inc., New York, NY, 1991.
- [4] I. S. Venieris, J. D. Angelopoulos and G. I. Stassinopoulos, "DQDB MAN as a Transit Network for ATM CPNs," *Proc. INFOCOM, Chicago*, pp. 2351–2357, 1992.
- [5] William Stallings, *Data and Computer Communication*. New York: Macmillan Publishing Company, 4th ed., 1994.
- [6] Oran Sharon and Adrian Segall, "A simple Scheme for Slot Reuse Without Latency for a Dual Bus Configuration," *IEEE/ACM Transactions on Networking*, vol. 1, no. 1, pp. 96–104, February 1993.
- [7] P. J. Kiviat, H. Markowitz ,and R. Villanueva, *Simscrip II.5 Programming Language*. California: CACI, 1987.
- [8] William H. Beyer, *CRC Standard Mathematical Tables*. Florida: CRC Press Inc., 28th ed., 1987.

## Bibliography

- [9] Kiriakos Apostolidis, "A Reservation Protocol for Packet Voice and Data Integration in Unidirectional Bus Networks," *IEEE Transactions on Communications*, vol. 41, no. 3, pp. 478–485, March 1993.
- [10] Xiaomei Qian, Sharad Kumar, Dhadesugor Vaman, Shukri Wakid and David Cypher, "Provision of Isochronous Service on IEEE 802.6," *IEEE Transactions on Communications*, vol. 42, no. 2/3/4, pp. 1778–1784, February/March/April 1994.
- [11] L. J. Ng, R. W. Donaldson and A. D. Malayan, "Distributed Architectures and Databases for Intelligent Personal Communication Networks," *ICWC*, pp. 300–304, 1992.
- [12] James Lane, "ATM kints voice, data on any net," *IEEE SPECTRUM*, pp. 42–45, February 1994.
- [13] Cem Ersoy and Shivendra S. Panwar, "Topological Design of Interconnected LAN/MAN Networks," *IEEE Journal on Selected Areas in Communications*, vol. 11, no. 8, pp. 1172–1182, October 1993.
- [14] Oran Sharon and Adrian Segall, "On the Efficiency of Slot Reuse in the Dual Bus Configuration," *IEEE/ACM Transactions on Networking*, vol. 2, no. 1, pp. 89–100, February 1994.
- [15] Kathleen S. Meier-Hellstern and David J. Goodman, "Network Protocols for the Cellular Packet Switch," *IEEE Transactions on Communications*, vol. 42, no. 2/3/4, pp. 89–100, February/March/April 1994.
- [16] Kin K. Leung, Bhaskar Sengupta and Raymond W. Yeung, "A Credit Manager for Traffic Regulation in High-Speed Networks: A Queueing Analysis," *IEEE Transactions on Networking*, vol. 1, no. 2, pp. 236–245, April 1993.

*Bibliography*

- [17] Riaz Ahmad and Fred Halsall, "Interconnecting High-Speed LANs and Backbones," *IEEE Network Magazine*, pp. 36–43, September 1993.
- [18] B. Grela-M'Poko, M. Mehmet-Ali and J. F. Hayes, "An Approximate Performance Analysis of Interconnected Token-Passing Rings," *IEEE Transactions on Communications*, vol. 41, no. 8, pp. 1142–1146, August 1993.
- [19] Joseph C. S. Cheung, Mark A. Beach and Joseph P. McGeehan, "Network Planning for Third-Generation Mobil Radio Systems," *IEEE Communications Magazine*, vol. 32, no. 11, pp. 54–59, November 1994.
- [20] Wolfgang Fischer, Eugen Wallmeier, Thomas Worster, Simon P. Davis and Andrew Hayter, "Data Communications Using ATM: Architecture, Protocols, and Resource Management," *IEEE Communications Magazine*, vol. 32, no. 8, pp. 24–33, August 1994.
- [21] Brett J. Vickers and Tatsuya Suda, "Connectionless Service for Public ATM Networks," *IEEE Communications Magazine*, vol. 32, no. 8, pp. 34–42, August 1994.
- [22] Toshikazu Suzuki, "ATM Adaptation Layer Protocol," *IEEE Communications Magazine*, vol. 32, no. 4, pp. 80–83, April 1994.
- [23] Brian W. Unger, Douglas J. Goetz and Stephen W. Maryka, "Simulation of SS7 Common Channel Signaling," *IEEE Communications Magazine*, vol. 32, no. 3, pp. 52–62, March 1994.
- [24] William E. Burr, Shukri Wakid, Xiaomei Qian and Dhadesugoor Vaman, "A Comparison of FDDI Asynchronous Mode and DQDB Queue Arbitrated Mode Data Transmission for Metropolitan Area Network Applications," *IEEE Transactions on Communications*, vol. 42, no. 2/3/4, pp. 1758–1768, February/March/April 1994.

*Bibliography*

- [25] Douglas W. Browning, "Flow Control in High-Speed Communication Networks," *IEEE Transactions on Communications*, vol. 42, no. 7, pp. 2480–2489, July 1994.
- [26] Dimitri Bertsekas and Robert Gallager, *Data Networks*. New Jersey: Prentice-Hall Inc, 1992.
- [27] Leslie J. Ng, "Distributed Architectures and Database for Personal Communications Networks and Services," Master's thesis, UBC Department of Electrical Engineering M.A.Sc. Thesis, March 1993.